

# Random Delay in Network Coding for Bidirectional Relaying

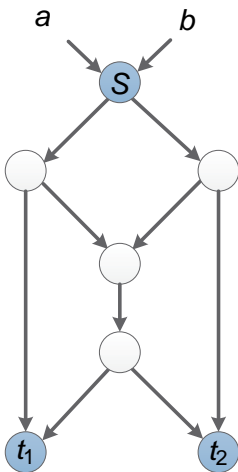
Niv Voskoboynik, Haim Permuter and Asaf Cohen

Ben Gurion University

June, 2014

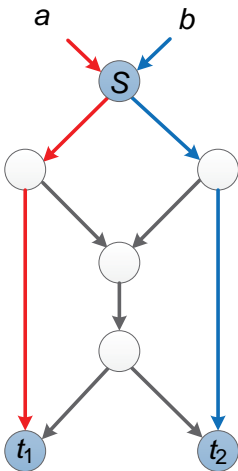
# What is Network Coding?

## Butterfly network



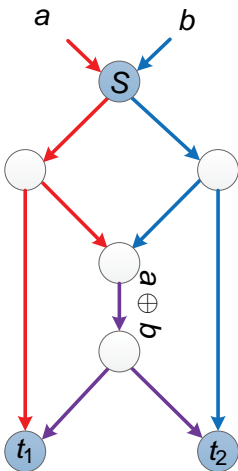
# What is Network Coding?

Using simple routing:



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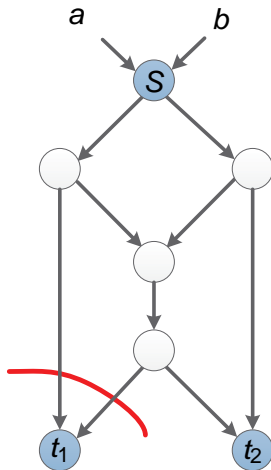
Using Network Coding:



# Max-Flow Min-Cut Theorem

The maximum value of a flow is equal to the minimum cut

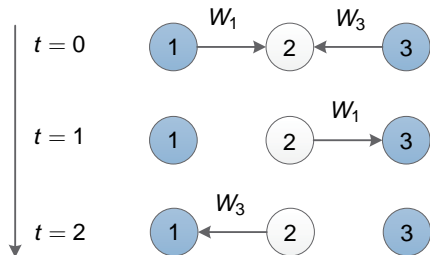
(R.Ahlswede *et al.* 2000)



How to minimize the number of transmissions?

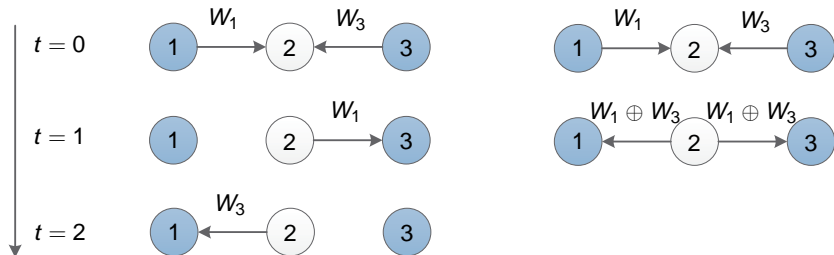
# Motivation

How to minimize the number of transmissions?



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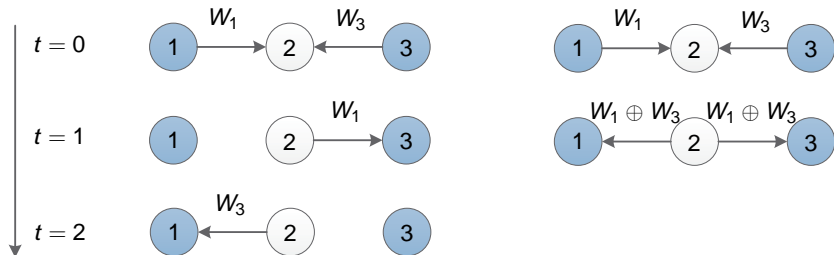
How to minimize the number of transmissions?





# Motivation

How to minimize the number of transmissions?



- Decoding:  $(W_1 \oplus W_3) \oplus W_3 = W_1$

# Motivation - Random Delay

Challenge: Random delay



- At time  $t$ : Node 2 has  $W_1(t - d_1^t)$  and  $W_3(t - d_3^t)$ , where  $d_1^t, d_3^t \sim \text{unif}(1, 2, \dots, D)$

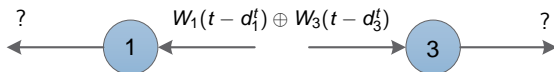
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Problem: How to decode?



- With which message to XOR with?

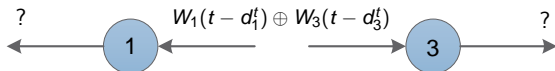
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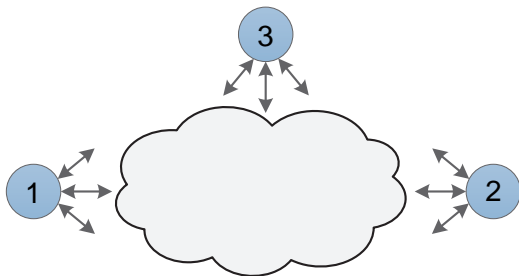
Problem: How to decode?



- With which message to XOR with?
- Which message is decoded?

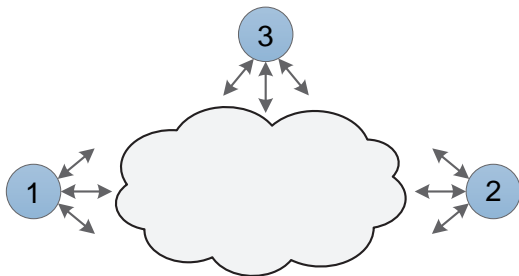
# Goal

Find a coding scheme for a bidirectional graph with random delay



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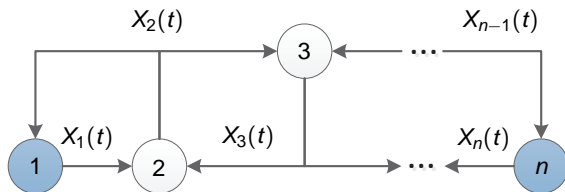


- Approach: Use building blocks

# Building Block 1 - Line Topology

One generalization of the first example is a line topology

(Y.Wu *et al.* 2004)



## Coding scheme

- Assuming  $d_1^t, d_n^t = 1, \forall t$



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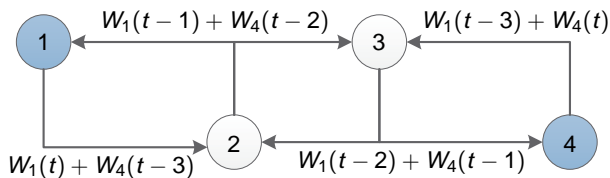
$$X_1(t) = W_1(t) + W_n(t - (n - 1))$$

- Generated by

$$X_r(t) = X_{r+1}(t - 1) + X_{r-1}(t - 1) + X_r(t - 2)$$

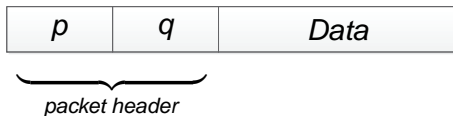
## Example

- A line topology with four nodes



A generalization to random delay

$$X_r(t) = W_1(p) + W_n(q)$$



## Summary

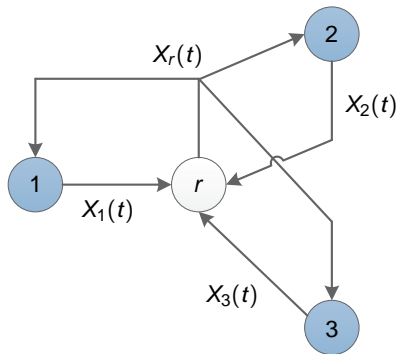
- A coding scheme for a line topology with rate of 1

## Summary

- A coding scheme for a line topology with rate of 1
- A first building block for a general topology

## Building Block 2 - Star Topology

A star topology with three source nodes





## Coding Scheme

- Node  $r$  transmits:

$$X_r(t) = a_1 W_1(t - 1) + a_2 W_2(t - 1) + a_3 W_3(t - 1)$$

## Coding Scheme

- Node  $r$  transmits:

$$X_r(t) = a_1 W_1(t-1) + a_2 W_2(t-1) + a_3 W_3(t-1)$$
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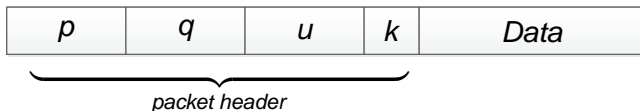
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$$X_r(t+2) = X_1(t) + X_2(t) + X_3(t)$$

A generalization to random delay

$$X_r(t) = k_1 W_1(p) + k_2 W_2(q) + k_3 W_3(u).$$

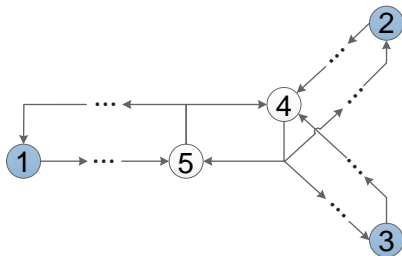


## Summary

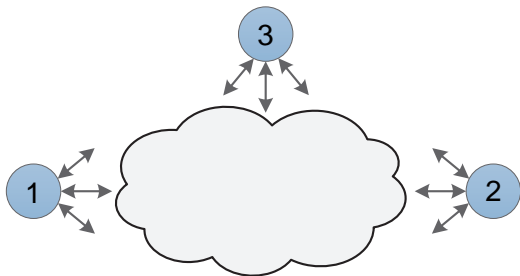
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## Summary

- A coding scheme for a star topology with rate of 0.5
- A combination is called a line-star topology

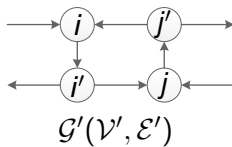
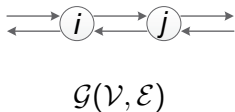


- We define a bidirectional graph  $\mathcal{G}(\mathcal{V}, \mathcal{E})$



## Wired Model

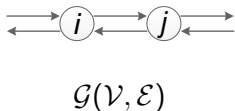
- We construct an equivalent graph  $\mathcal{G}'(\mathcal{V}', \mathcal{E}')$



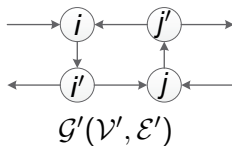


## Wired Model

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minimum transmissions



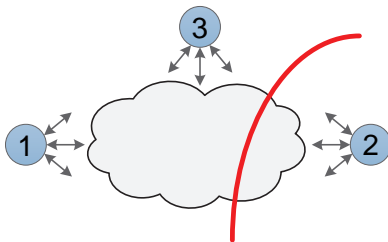
maximum rate

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- $C_{i,j}$  is the value of the cut-set bound between  $i$  and  $j$

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- $h = \min_{i \in S} C_{i, S \setminus \{i\}}$

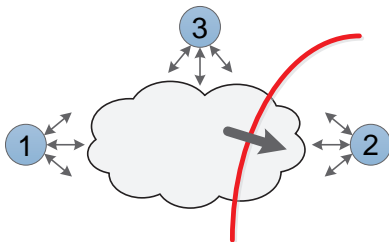


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# General Topology

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- Assumption: all rates are equal ( $R_i = R, \forall i \in \mathcal{S}$ )
- Upper bound:  $R \leq \frac{h}{2}$

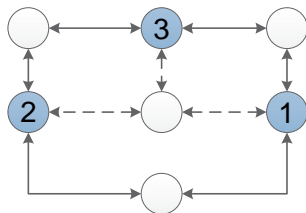


## Coding Scheme

- Partition each network  $\mathcal{G}'$  into line and star networks

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- $\mathcal{R}$  represents the line topologies
- $\mathcal{Q}$  represents the star topologies





## Lemma

There exist  $\mathcal{R}$  and  $\mathcal{Q}$  such that  $|\mathcal{R}| + \frac{|\mathcal{Q}|}{2} \geq \frac{h}{2}$

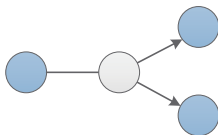
(N.Voskoboynik *et al.* 2014)

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- Element in  $\mathcal{Q}$ :

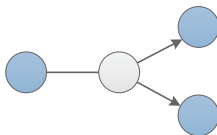


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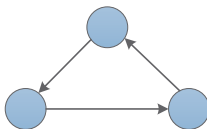
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- Element in  $\mathcal{Q}$ :



- Element in  $\mathcal{R}$ :



Construct the coding scheme

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Construct the coding scheme

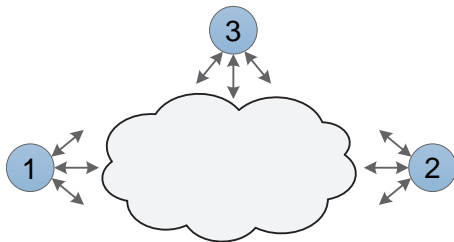
- Find  $\mathcal{R}$
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## Construct the coding scheme

- Find  $\mathcal{R}$
- At each line use the coding scheme from building block 1
- Find  $\mathcal{Q}$
- At each star use the coding scheme from building block 2

## Summary

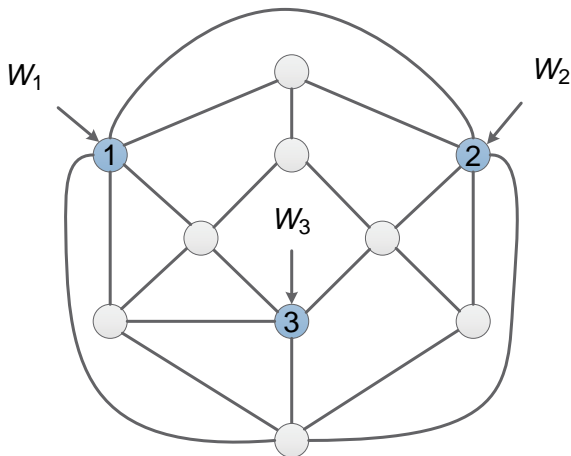
- Using line and star topologies, we construct a coding scheme for a general topology





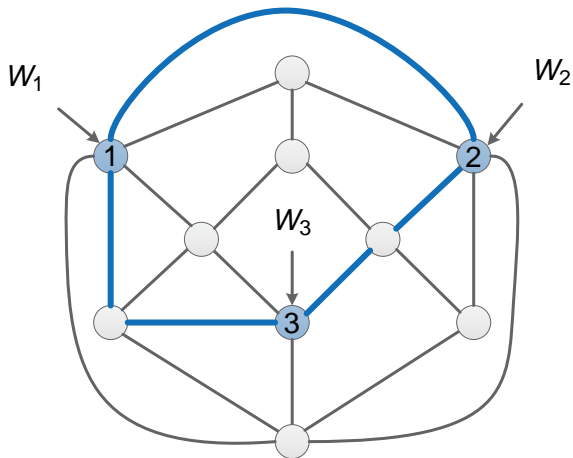
# Example 1

$h = 4$  and therefore, the upper bound is 2



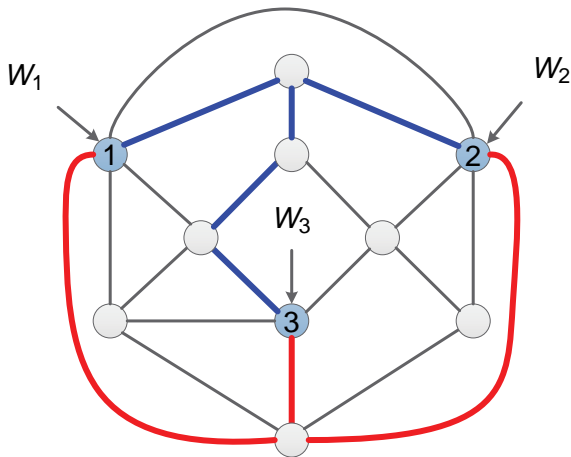
# Example 1

The set  $\mathcal{R}$  includes one ring, i.e.,  $|\mathcal{R}| = 1$



# Example 1

The set  $\mathcal{Q}$  includes two star networks, i.e.,  $|\mathcal{Q}| = 2$



## Example 2

Maximum values of  $|\mathcal{R}|$  and  $|\mathcal{Q}|$  in a graph  $\mathcal{G}(\mathcal{V}, \mathcal{E})$

Maximum value	General graph	Complete graph
$ \mathcal{R} $	$\lfloor \frac{ \mathcal{V} }{3} \rfloor$	1
$ \mathcal{Q} $	$ \mathcal{V}  - 2$	$ \mathcal{V}  - 3$

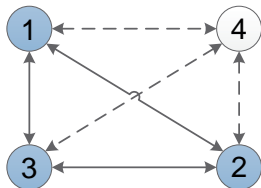


Figure: Complete graph

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Thank You!