Broadcast Channels with Cooperation: Capacity and Duality for the Semi-Deterministic Case

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Outline

- Channel-source duality for BCs
- Semi-deterministic BC with decoder cooperation
- Source coding dual
- Capacity results
- Summary

"There is a curious and provocative duality between the properties of a source with a distortion measure and those of a channel..."
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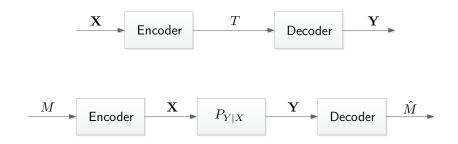
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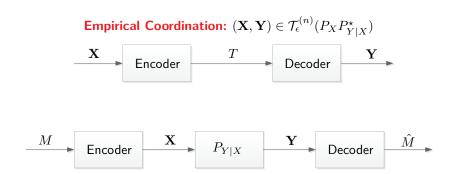
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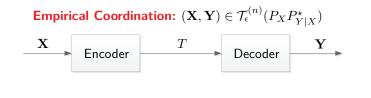
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- A formal proof of duality is still absent.
- ullet Solving one problem \Longrightarrow Valuable insight into solving dual.



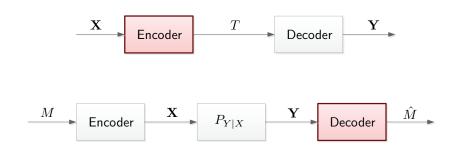


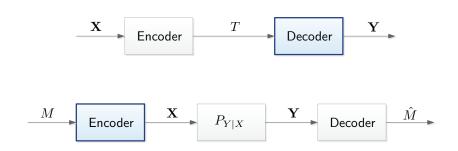
Point-to-Point Case:

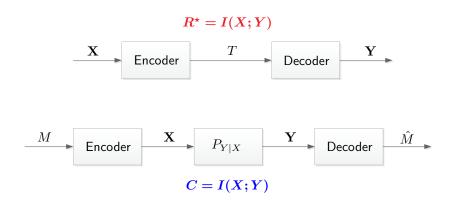


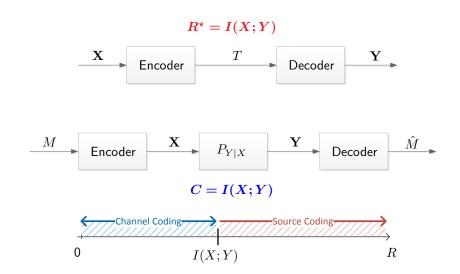


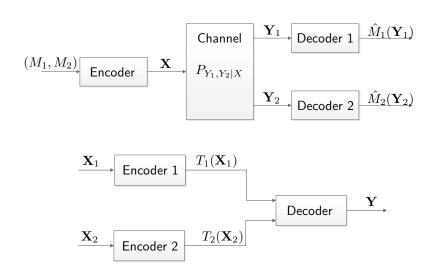
Fixed-Type Code: $(\mathbf{X},\mathbf{Y})\in\mathcal{T}^{(n)}_{\epsilon}(P_X^{\star}P_{Y|X})$

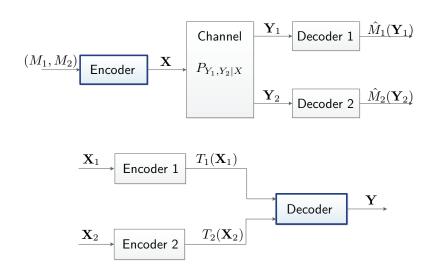


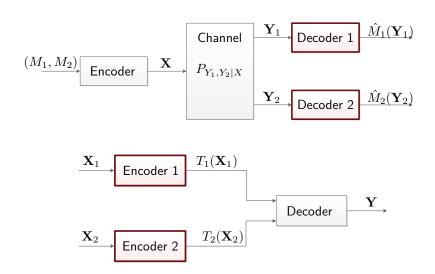












Probabilistic relations are preserved:

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Broadcast Channel

 $(\mathbf{X}, \mathbf{Y}_1, \mathbf{Y}_2) \in \mathcal{T}_{\epsilon}^{(n)} \left(P_X^{\star} P_{Y_1, Y_2 \mid X} \right)$

Dual Source Coding Setting

$$(\mathbf{X}_1, \mathbf{X}_2, \mathbf{Y}) \in \mathcal{T}_{\epsilon}^{(n)} \left(P_{X_1, X_2} P_{Y|X_1, X_2}^{\star} \right)$$

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e.g., Markov relations, deterministic functions, etc.

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Additional Principles:

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Causal/non-causal <u>encoder</u> CSI ←→ Causal/non-causal <u>decoder</u> SI

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- <u>Decoder</u> cooperation ←→ <u>Encoder</u> cooperation

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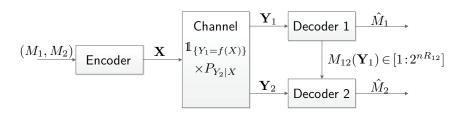
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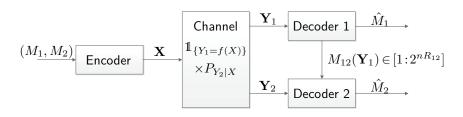
Additional Principles:

- Causal/non-causal <u>encoder</u> CSI ←→ Causal/non-causal <u>decoder</u> SI
- <u>Decoder</u> cooperation ←→ <u>Encoder</u> cooperation
- ★ Result Duality: Information measures at the corner points coincide! ★

Without cooperation: [Gelfand vs. Pinsker, 1980] and [Wyner, 1975]&[Ahlswede-Körner, 1975]



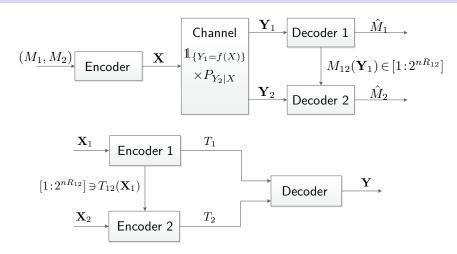
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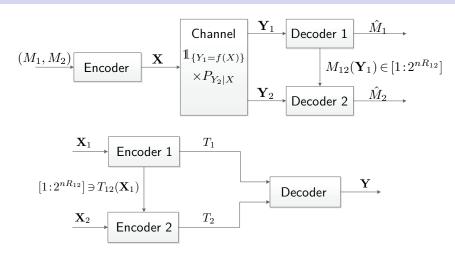
BCs with Cooperation:

- Physicaly degraded (PD) [Dabora and Servetto, 2006].
- Relay-BC [Liang and Kramer, 2007].
- State-dependent PD [Dikstein, Permuter and Steinberg, 2014].
- Degraded message sets / PD with parallel conf. [Steinberg, 2015].

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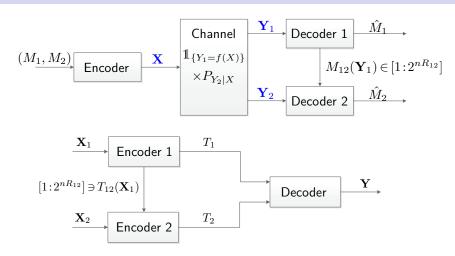


Semi-Deterministic BC

WAK Problem

$$(\mathbf{X},\mathbf{Y}_{1},\mathbf{Y}_{2})\in\mathcal{T}_{\epsilon}^{(n)}(P_{X}^{\star}\mathbb{1}_{\{Y_{1}=f(X)\}}P_{Y_{2}|X}) \quad \Longleftrightarrow \quad (\mathbf{Y},\mathbf{X}_{1},\mathbf{X}_{2})\in\mathcal{T}_{\epsilon}^{(n)}(P_{Y}\mathbb{1}_{\{X_{1}=f(Y)\}}P_{X_{2}|Y}^{\star})$$

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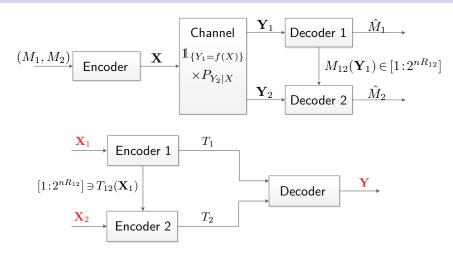


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Theorem (Coordination-Capacity Region)

For a desired coordination PMF $P_{X_2}P_{Y|X_2}\mathbb{1}_{\{X_1=f(Y)\}}$:

$$\mathcal{C}_{\textit{WAK}} = \bigcup \left\{ \begin{array}{c} R_{12} \geq I(V; X_1) - I(V; X_2) \\ R_1 \geq H(X_1 | V, U) \\ R_2 \geq I(U; X_2 | V) - I(U; X_1 | V) \\ R_1 + R_2 \geq H(X_1 | V, U) + I(V, U; X_1, X_2) \end{array} \right\}$$

where the union is over all $P_{X_1,X_2}P_{V|X_1}P_{U|X_2,V}P_{Y|X_1,U,V}$ with $P_{X_2}P_{Y|X_2}\mathbb{1}_{\{X_1=f(Y)\}}$ as marginal.

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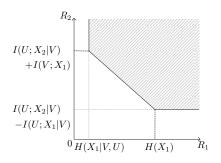
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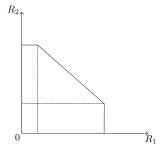
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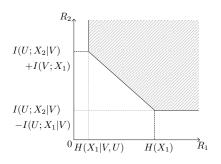
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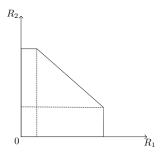
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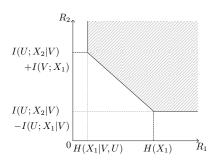


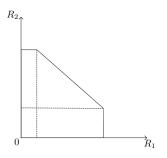




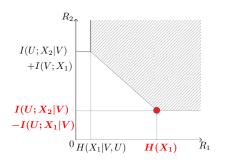


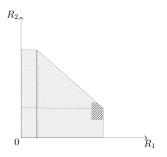
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$R_{12} = I(V; X_1) - I(V; X_2)$	
(R_1,R_2) at Lower Corner Point:	(R_1,R_2) at Lower Corner Point:
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(R_1,R_2) at Upper Corner Point:	(R_1,R_2) at Upper Corner Point:
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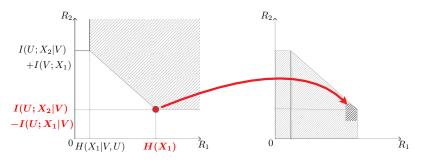


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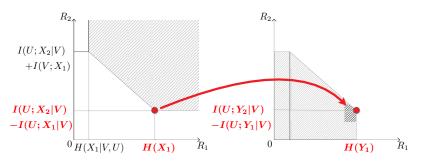




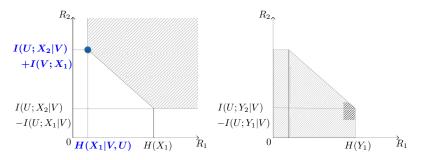
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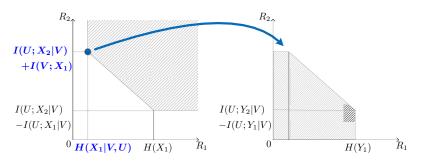
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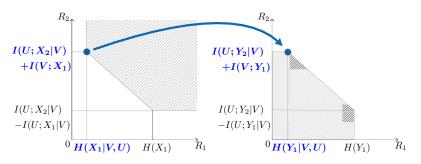
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Semi-Deterministic BC with Cooperation - Solution

Theorem (Capacity Region)

The capacity region is:

$$\mathcal{C}_{BC} = \bigcup \left\{ \begin{array}{c} R_{12} \geq I(V;Y_1) - I(V;Y_2) \\ R_1 \leq H(Y_1) \\ R_2 \leq I(V,U;Y_2) + R_{12} \\ R_1 + R_2 \leq H(Y_1|V,U) + I(U;Y_2|V) + I(V;Y_1) \end{array} \right\}$$

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• Later: Achievability and converse proofs for an alternative region.

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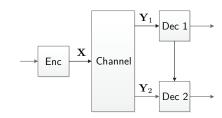
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where the union is over all $P_{V,U,Y_1,X}P_{Y_2|X}\mathbb{1}_{\{Y_1=f(X)\}}$.

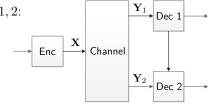
- Later: Achievability and converse proofs for an alternative region.
- \bullet $\mathcal{C}_{\mathsf{BC}}$ emphasizes duality.



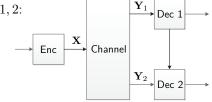
• Rate Splitting: $M_j=(M_{j0},M_{jj}),\ j=1,2$:

Enc X Channel Y₁ Dec 1

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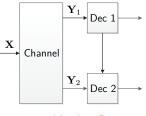
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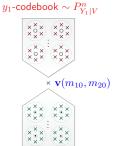
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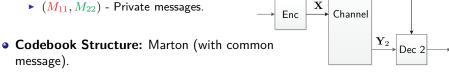




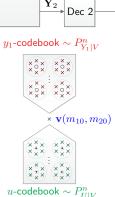
u-codebook $\sim P_{U|V}^n$

Enc

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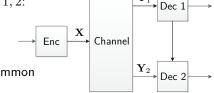


Cooperation:

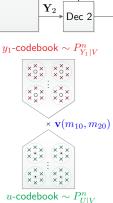


Dec 1

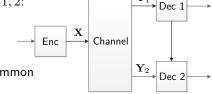
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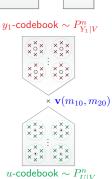
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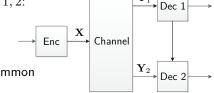
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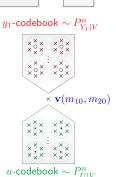


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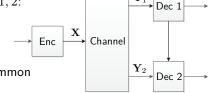


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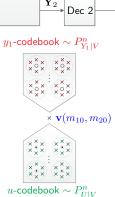


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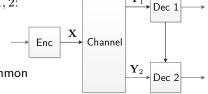


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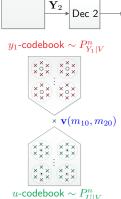
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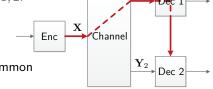
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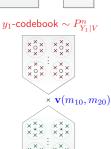


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u-codebook $\sim P_{U|V}^n$

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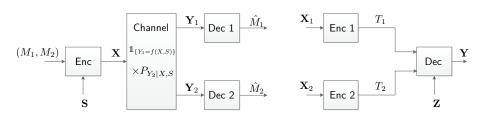
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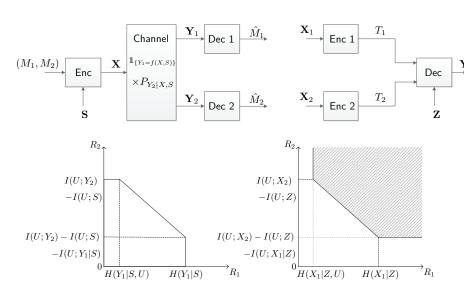
Thank you!

State-Dependant Semi-Deterministic BC vs. Dual:

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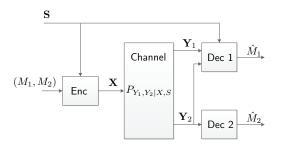


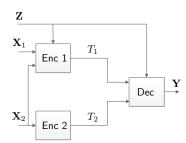
State-Dependant Semi-Deterministic BC vs. Dual:



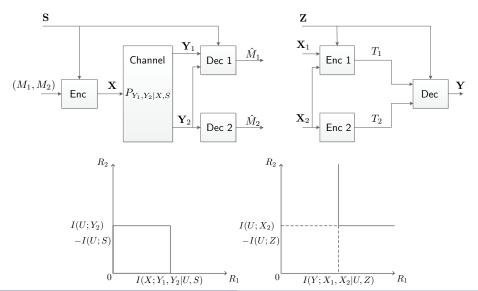
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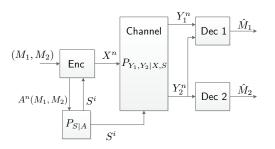


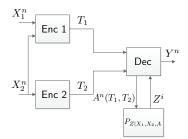
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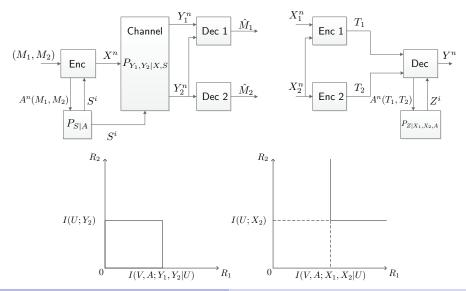
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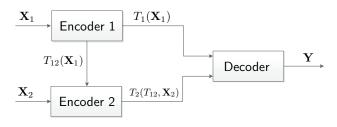
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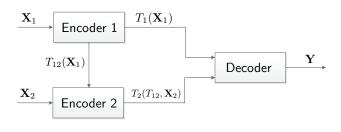




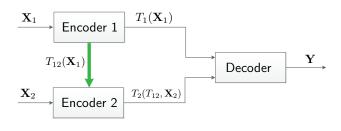
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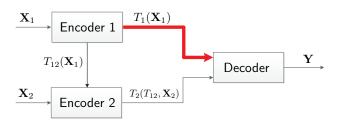


Rate	Corner Point 1	Corner Point 2
R_{12}	$I(V;X_1) - I(V;X_2)$	$I(V;X_1) - I(V;X_2)$
R_1	$H(X_1)$	$H(X_1 V,U)$
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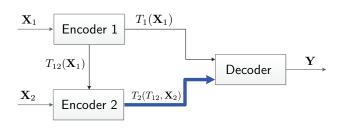
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AK Problem with Cooperation - Proof Outline

Converse:

AK Problem with Cooperation - Proof Outline

Converse:

Standard techniques while defining

$$V_i = (T_{12}, X_1^{n \setminus i}, X_{2,i+1}^n),$$

 $U_i = T_2,$

for every $1 \le i \le n$.

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Time mixing properties.

$\begin{tabular}{ll} Semi-Deterministic BC with Cooperation - Achievability \\ Outline \end{tabular}$

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• Enc X^n Channel Y_2^n Dec 2

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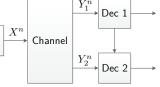
 (M_{11}, M_{22}) Private messages.

 Channel

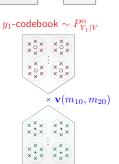
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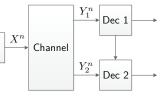
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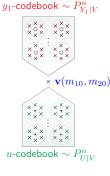
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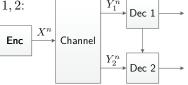


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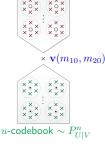


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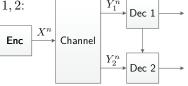


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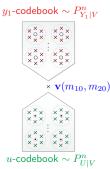


 y_1 -codebook $\sim P_{V, |V}^n$

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- **Gain:** Dec. 2 reduces search space of V by R_{12} .



Via telescoping identities:

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$$= \sum_{i=1}^{n} \left[I(M_{2}; Y_{2,i}^{n} | M_{12}, Y_{1}^{i-1}) - I(M_{2}; Y_{2,i+1}^{n} | M_{12}, Y_{1}^{i}) \right] + I(M_{2}; M_{12})$$

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$$\begin{split} &H(M_2) - n\epsilon_n \leq I(M_2; Y_2^n | M_{12}) + I(M_2; M_{12}) \\ &= \sum_{i=1}^n \Big[I(M_2; Y_{2,i}^n | M_{12}, Y_1^{i-1}) - I(M_2; Y_{2,i+1}^n | M_{12}, Y_1^i) \Big] + I(M_2; M_{12}) \\ &= \sum_{i=1}^n \Big[I(M_2; Y_{2,i} | M_{12}, Y_1^{i-1}, Y_{2,i+1}^n) - I(M_2; Y_{1,i} | M_{12}, Y_1^{i-1}, Y_{2,i+1}^n) \Big] \\ &\quad + I(M_2; M_{12}) \end{split}$$

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- 1. Auxiliaries: $V_i = (M_{12}, Y_1^{i-1}, Y_{2,i+1}^n)$ and $U_i = M_2$.
- 2. Telescoping identities [Kramer, 2011], e.g.,

$$\begin{split} &H(M_2) - n\epsilon_n \leq I(M_2; Y_2^n | M_{12}) + I(M_2; M_{12}) \\ &= \sum_{i=1}^n \Big[I(M_2; Y_{2,i}^n | M_{12}, Y_1^{i-1}) - I(M_2; Y_{2,i+1}^n | M_{12}, Y_1^i) \Big] + I(M_2; M_{12}) \\ &= \sum_{i=1}^n \Big[I(M_2; Y_{2,i} | M_{12}, Y_1^{i-1}, Y_{2,i+1}^n) - I(M_2; Y_{1,i} | M_{12}, Y_1^{i-1}, Y_{2,i+1}^n) \Big] \\ &\quad + I(M_2; M_{12}) \end{split}$$

★ Replaces 2 uses of Csiszár Sum Identity.