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Introduction

1.1 Typographic conventions

Different fonts and styles are used throughout this manual to indicate special words or text. Computer prompts, responses and filenames will be printed in constant-spaced type. When the filename is the name of a standard header file, the name will be enclosed in angle brackets, e.g. `<stdio.h>`. These header files can be found in the INCLUDE directory of your distribution.

Samples of code, C keywords or types, assembler instructions and labels will be printed in a bold constant-space type. When used at the beginning of a sentence, a capital letter will be used. C code is case sensitive so this capitalization may need to be removed if entering the code into a program.

Particularly useful points and new terms will be emphasised using italicised type. When part of a term requires substitution, that part should be printed in the appropriate font, but in italics. For example: `#include <filename.h>`.

1.2 Using This Manual

This manual is a comprehensive guide and reference to using PICC-18. The chapters included are as follows:

▲ Tutorials to aid in the understanding and usage of HI-TECH’s C cross compilers
▲ How to use the PICC-18 command-line interface
▲ In-depth description of the C compiler
▲ How to use the assembler
▲ How to use the linker and other utilities
▲ Error messages and their meaning
▲ Description of provided library functions

For information on installing PICC-18, using the on-line manual and getting started, see the Quick Start Guide.
Tutorials

The following are tutorials to aid in the understanding and usage of HI-TECH’s C cross compilers. These tutorials should be read in conjunction with the appropriate sections in the manual as they are aimed at giving a general overview of certain aspects of the compiler. Some of the tutorials here are generic to all HI-TECH C compilers and may include information not specific for the compiler you are using.

2.1 Overview of the compilation process

This tutorial gives an overview of the compilation process that takes place with HI-TECH C compilers in terms of how the input source files are processed. The origin of files that are produced by the compiler is discussed as well as their content and function.

2.1.1 Compilation

When a program is compiled, it is done so by many separate applications whose operations are controlled by either the command-line driver (CLD) or HPD driver (HPD). In either case, HPD or the CLD take the options specified by the programmer (menu options in the case of HPD, or command-line arguments for the CLD) to determine which of the internal applications need to be executed and what options should be sent to each. When the term compiler is used, this is intended to denote the entire collection of applications and driver that are involved in the process. In the same way, compilation refers to the complete transformation from input to output by the compiler. Each application and its function is discussed further on in this document.

The compiler drivers use several files to store options and information used in the compilation process and these file types are shown in Table 2 - 1 on page 18. The HPD driver stores the compiler options into a project file which has a .prj extension. HPD itself stores its own configurational settings in an INI file, e.g. HPD51.ini in the BIN directory of your distribution. This file stores information such as colour values and mouse settings. Users who wish to use the CLD can store the command line arguments in a DOS batch file.

Some compilers come with chip info files which describe the memory arrangements of different chip types. If necessary this file can be edited to create new chip types which can then be selected with the appropriate command-line option of from the select processor... menu. This file will also have a .ini extension and is usually in the LIB directory.

1. The command line driver and HPD driver have processor-specific names, such as PICC, C51, or HPDXA, HPDPIC etc.
The compilation process is discussed in the following sections both in terms of what takes place at each stage and the files that are involved. Reference should be made to Figure 2 - 1 on page 19 which shows the block diagram of the internal stages of the HI-TECH compiler, and the tables of file types throughout this tutorial which list the filename extension\(^2\) used by different file formats and the information which the file contains. Note that some older HI-TECH compilers do not include all the applications discussed below.

<table>
<thead>
<tr>
<th>extension</th>
<th>name</th>
<th>contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>.prj</td>
<td>project file</td>
<td>compiler options stored by HPD driver</td>
</tr>
<tr>
<td>.ini</td>
<td>HPD initialisation file</td>
<td>HPD environment settings</td>
</tr>
<tr>
<td>.bat</td>
<td>batch file</td>
<td>command line driver options stored as DOS batch file</td>
</tr>
<tr>
<td>.ini</td>
<td>chip info file</td>
<td>information regarding chip families</td>
</tr>
</tbody>
</table>

The internal applications generate output files and pass these to the next application as indicated in the figure. The arrows from one application (drawn as ellipses) to another is done via temporary files that have non-descriptive names such as \$003361.001. These files are temporarily stored in a directory pointed to by the DOS environment variable TEMP. Such a variable is created by a set DOS command. These files are automatically deleted by the driver after compilation has been completed.

### 2.1.2 The compiler input

The user supplies several things to the compiler to make a program: the input files and the compiler options, whether using the CLD or HPD. The compiler accepts many different input file types. These are discussed below.

It is possible, and indeed in a large number of projects, that the only files supplied by the user are C source files and possibly accompanying header files. It is assumed that anyone using our compiler is familiar with the syntax of the C language. If not, there is a seemingly endless selection of texts which cover this topic. C source files used by the HI-TECH compiler must use the extension .c as this extension is used by the driver to determine the file's type. C source files can be listed in any order on the command line if using the CLD, or entered into the source file list... dialogue box if using HPD.

A header file is usually a file which contains information related to the program, but which will not directly produce executable code when compiled. Typically they include declarations (as opposed to definitions) for functions and data types. These files are included into C source code by a preprocessor directive and are often called include files. Since header files are referenced by a command that includes

---

\(2\). The extensions listed in these tables are in lower case. DOS compilers do not distinguish between upper- and lower-case file names and extensions, but in the interest of writing portable programs you should use lower-case extensions in file names and in references to these files in your code as UNIX compilers do handle case correctly.
the file's name and extension (and possibly a path), there are no restrictions as to what this name can be although convention dictates a `.h` extension.

Although executable C code may be included into a source file, a file using the extension `.h` is assumed to have non-executable content. Any C source files that are to be included into other source files should still retain a `.c` extension. In any case, the practise of including one source file into another is best avoided as it makes structuring the code difficult, and it defeats many of the advantages of having a
compiler capable of handling multiple-source files in the first place. Header files can also be included into assembler files. Again, it is recommended that the files should only contain assembler declarations.

Table 2 - 2 Input file types

<table>
<thead>
<tr>
<th>extension</th>
<th>name</th>
<th>content</th>
</tr>
</thead>
<tbody>
<tr>
<td>.c</td>
<td>C source file</td>
<td>C source conforming to the ANSI standard possibly with extensions allowed by HI-TECH C</td>
</tr>
<tr>
<td>.h</td>
<td>header file</td>
<td>C/assembler declarations</td>
</tr>
<tr>
<td>.as</td>
<td>assembler file</td>
<td>assembler source conforming to the HI-TECH assembler format</td>
</tr>
<tr>
<td>.obj</td>
<td>(relocatable)</td>
<td>pre-compiled C or assembler source as HI-TECH relocatable object file</td>
</tr>
<tr>
<td>.lib</td>
<td>library file</td>
<td>pre-compiled C or assembler source in HI-TECH library format</td>
</tr>
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</table>

HI-TECH compilers comes with many header files which are stored in a separate directory of the distribution. Typically user-written header files are placed in the directory that contains the sources for the program. Alternatively they can be placed into a directory which can be searched by using a \(-I\) (CPP include paths...) option.

An assembler file contains assembler mnemonics which are specific to the processor for which the program is being compiled. Assembler files may be derived from C source files that have been previously compiled to assembler, or may be hand-written and highly-prized works of art that the programmer has developed. In either case, these files must conform to the format expected of the HI-TECH assembler that is part of the compiler. This processor-dependence makes assembly files quite un-portable and they should be avoided if C source can be made to perform the task at hand. Assembler files must have a .as extension as this is used by the compiler driver to determine the file’s type. Assembler files can be listed in any order on the command line if using the CLD, or entered into the source file list... dialogue box if using HPD, along with the C source files.

The compiler drivers can also be passed pre-compiled HI-TECH object files as input. These files are discussed below in Section 2.1.2.1 on page 21. These files must have a .obj extension. Object files can be listed in any order on the command line if using the CLD, or entered into the object file list... dialogue box if using HPD. You should not enter the names of object files here that have been compiled from source files already in the project, only include object files that have been pre-compiled and have no corresponding source in the project, such as the run-time file. For example, if you have included init.c into the project, you should not include init.obj into the object file list.

Commonly used program routines can be compiled into a file called a library file. These files are more convenient to handle and can be accessed quickly by the compiler. The compiler can accept library files directly like other source files. A .lib extension indicates the type of the file and so library files must
be named in this way. Library files can be listed in any order on the command line if using the CLD, or entered into the library file list... dialogue box if using HPD.

The HI-TECH library functions come pre-compiled in a library format and are stored in the LIB directory in your distribution.

2.1.2.1 Steps before linking

Of all the different types of files that can be accepted by the compiler, it is the C source files that require the most processing. The steps involved in compiling the C source files are examined first.

For each C source file, a C listing file is produced by an application called CLIST. The listing files contain the C source lines proceeded by a line number before any processing has occurred. The C listing for a small test program called main.c is shown in Table 2 - 3 on page 21.

<table>
<thead>
<tr>
<th>C source</th>
<th>C listing</th>
</tr>
</thead>
<tbody>
<tr>
<td>#define VAL 2</td>
<td>1: #define VAL 2</td>
</tr>
<tr>
<td>int a, b = 1;</td>
<td>2:</td>
</tr>
<tr>
<td>void</td>
<td>3: int a, b = 1;</td>
</tr>
<tr>
<td>main(void)</td>
<td>4:</td>
</tr>
<tr>
<td>{</td>
<td>5: void</td>
</tr>
<tr>
<td>/* set starting value */</td>
<td>6: main(void)</td>
</tr>
<tr>
<td>a = b + VAL;</td>
<td>7: {</td>
</tr>
<tr>
<td>}</td>
<td>8: /* set starting value */</td>
</tr>
<tr>
<td></td>
<td>9: a = b + 2;</td>
</tr>
<tr>
<td></td>
<td>10: }</td>
</tr>
</tbody>
</table>

The input C source files are also passed to the preprocessor, CPP. This application has the job of preparing the C source for subsequent interpretation. The tasks performed by CPP include removing comments and multiple spaces (such as tabs used in indentation) from the source, and executing any preprocessor directives in the source. Directives may, for example, replace macros with their replacement text (e.g. #define directives) or conditionally include source code subject to certain conditions (e.g. #if, #ifdef etc. directives). The preprocessor also inserts header files, whether user- or compiler-supplied, into the source. Table 2 - 4 on page 22 shows preprocessor output for the test program.

The output of the preprocessor is C source, but it may contain code which has been included by the preprocessor from header files and conditional code may have been omitted. Thus the preprocessor output usually contains similar, but different code to the original source file. The preprocessor output is often referred to as a module or translational unit. The term "module" is sometimes used to describe the actual source file from which the "true" module is created. This is not strictly correct, but the meaning is clear enough.
The code generation that follows operates on the **CPP** output module, not the C source and so special steps must be taken to be able to reconcile errors and their position in the original C source files. The # 1 main.c line in the preprocessor output for the test program is included by the preprocessor to indicate the filename and line number in the C source file that corresponds to this position. Notice in this example that the comment and macro definition have been removed, but blank lines take their place so that line numbering information is kept intact.

Like all compiler applications, the preprocessor is controlled by the compiler driver (either the CLD or HPD). The type of information that the driver supplies the preprocessor includes directories to search for header files that are included into the source file, and the size of basic C objects (such as int, double, char *, etc.) using the -S, -SP options so that the preprocessor can evaluate preprocessor directives which contain a `sizeof(type)` expression. The output of the preprocessor is not normally seen unless the user uses the `-PRE` option in which case the compiler output can then be re-directed to file.

The output of **CPP** is passed to **P1**, the *parser*. The parser starts the first of the hard work involved with turning the description of a program written in the C language into the actual executable itself consisting of assembler instructions. The parser scans the C source code to ensure that it is valid and then replaces C expressions with a modified form of these. (The description of code generation that follows need not be followed to understand how to use the HI-TECH compiler, but has been included for curious readers.)

For example the C expression `a = b + 2` is re-arranged to a *prefix notation* like `= a + b 2`. This notation can easily be interpreted as a tree with `=` at the apex, `a` and `+` being branches below this, and `b` and `2` being sub-branches of the addition. The output of the parser is shown in Table 2 - 6 on page 23 for our small C program. The assignment statement in the C source has been highlighted as well as the output the parser generates for this statement. Notice that already the global symbols in the parser output

### Table 2 - 4 preprocessor output

<table>
<thead>
<tr>
<th>C source</th>
<th>Pre-processed output</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>#define VAL 2</code></td>
<td># 1 &quot;main.c&quot;</td>
</tr>
<tr>
<td><code>int a, b = 1;</code></td>
<td>int a, b = 1;</td>
</tr>
<tr>
<td><code>void main(void)</code></td>
<td>void main(void)</td>
</tr>
<tr>
<td>`{</td>
<td>`</td>
</tr>
<tr>
<td>/* set starting value */</td>
<td>`</td>
</tr>
<tr>
<td><code>a = b + VAL;</code></td>
<td><code>a = b + 2;</code></td>
</tr>
</tbody>
</table>
| }                                                   | }
Overview of the compilation process

Table 2 - 5 Intermediate and Support files

<table>
<thead>
<tr>
<th>extension</th>
<th>name</th>
<th>contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>.pre</td>
<td>pre-processed file</td>
<td>C source or assembler after the pre-processing stage</td>
</tr>
<tr>
<td>.lst</td>
<td>C listing file</td>
<td>C source with line numbers</td>
</tr>
<tr>
<td>.lst</td>
<td>assembler listing</td>
<td>C source with corresponding assembler instructions</td>
</tr>
<tr>
<td>.map</td>
<td>map file</td>
<td>symbol and psect relocation information generated by the linker</td>
</tr>
<tr>
<td>.err</td>
<td>error file</td>
<td>compiler warnings and errors resulting from compilation</td>
</tr>
<tr>
<td>.rlf</td>
<td>relocation listing file</td>
<td>information necessary to update list file with absolute addresses</td>
</tr>
<tr>
<td>.sdb</td>
<td>symbolic debug file</td>
<td>object names and types for module</td>
</tr>
<tr>
<td>.sym</td>
<td>symbol file</td>
<td>absolute address of program symbols</td>
</tr>
</tbody>
</table>

have had an underscore character pre-pended to their name. From now on, reference will be made to them using these symbols. The other symbols in this highlighted line relate to the constant. The ANSI standard states that the constant 2 in the source should be interpreted as a `signed int`. The parser ensures this is the case by casting the constant value. The `->` symbol represents the cast and the `'i` represents the type. Line numbering, variable declarations and the start and end of a function definition can be seen in this output.

Table 2 - 6 Parser output

<table>
<thead>
<tr>
<th>C source</th>
<th>Parsed output</th>
</tr>
</thead>
<tbody>
<tr>
<td>#define VAL 2</td>
<td></td>
</tr>
<tr>
<td>int a, b = 1;</td>
<td></td>
</tr>
<tr>
<td>void main(void)</td>
<td></td>
</tr>
<tr>
<td>{</td>
<td></td>
</tr>
<tr>
<td>/* set starting value */</td>
<td></td>
</tr>
<tr>
<td>a = b + VAL;</td>
<td></td>
</tr>
<tr>
<td>}</td>
<td></td>
</tr>
<tr>
<td>Version 3.2 HI-TECH Softwa...</td>
<td></td>
</tr>
<tr>
<td>&quot;3 main.c</td>
<td></td>
</tr>
<tr>
<td>[v_a `i l e ]</td>
<td></td>
</tr>
<tr>
<td>[v_b `i l e ]</td>
<td></td>
</tr>
<tr>
<td>[i_b</td>
<td></td>
</tr>
<tr>
<td>-&gt; l `i</td>
<td></td>
</tr>
<tr>
<td>]</td>
<td></td>
</tr>
<tr>
<td>&quot;7</td>
<td></td>
</tr>
<tr>
<td>[v_main `(v l e ]</td>
<td></td>
</tr>
<tr>
<td>{</td>
<td></td>
</tr>
<tr>
<td>[e :U _main ]</td>
<td></td>
</tr>
<tr>
<td>[f ]</td>
<td></td>
</tr>
<tr>
<td>&quot;9</td>
<td></td>
</tr>
<tr>
<td>[; ;main.c: 9: b = a + 2;</td>
<td></td>
</tr>
<tr>
<td>[e = _a + _b -&gt; 2 `i ]</td>
<td></td>
</tr>
<tr>
<td>&quot;10</td>
<td></td>
</tr>
<tr>
<td>[; ;main.c: 10: }</td>
<td></td>
</tr>
<tr>
<td>[e :UE 1 ]</td>
<td></td>
</tr>
</tbody>
</table>
Tutorials

It is the parser that is responsible for finding a large portion of the errors in the source code. These errors will relate to the syntax of the source code. The parser also reports warnings if the code is unusual.

The parser passes its output directly to the next stage in the compilation process. There are no driver options to force the parser to generate parsed-source output files as these files contain no useful information for the programmer.

Now the tricky part of the compilation: code generation. The code generator converts the parser output into assembler mnemonics. This is the first step of the compilation process which is processor-specific. Whereas all HI-TECH preprocessors and parsers have the same name and are in fact the same application, the code generators will have a specific, processor-based name, for example CGPIC, or CG51.

The code generator uses a set of rules, or productions, to produce the assembler output. To understand how a production works, consider the following analogy of a production used to generate the code for the addition expression in our test program. "If you can get one operand into a register" and "one operand is a int constant" then here is the code that will perform a 2-byte addition of them. Here, each quoted string would represent a sub-production which would have to be matched. The first string would try to get the contents of _a into a register by matching further sub-productions. If it cannot, this production cannot be used and another will be tried. If all the sub-productions can be met, then the code that they produce can be put together in the order specified by the production tree. Not all productions actually produce code, but are necessary for the matching process.

If no matching production/subproductions can be found, the code generator will produce a Can’t generate code for this expression error. This means that the original C source code was legal and that the code generator did try to produce assembler code for it, but that in this context, there are no productions which can match the expression.
Typically there may be around 800 productions to implement a full code generator. There were about a dozen matching productions used to generate code for the statement highlighted in Table 2 - 7 on page 24 using the XA code generator. It checked about 70 productions which were possible matches before finding a solution. The exact code generation process is too complex to describe in this document and is not required to be able to use the compiler efficiently.

The user can stop the compilation process after code generation by issuing a `-s` (*compile to .as*) option to the driver. In this case, the code generator will leave behind assembler files with a `.as` extension. Table 2 - 7 on page 24 shows output generated by the XA code generator. Only the assembler code for the opening brace of `_main` and the highlighted source line is shown. This output will be different for other compilers and compiler options.

The code generator may also produce debugging information in the form of an `.sdb` file. This operation is enabled by using the `-g` (*source level debug info*) option. One debug file is produced for each module that is being compiled. These ASCII files contain information regarding the symbols defined in each module and can be used by debugging programs. Table 2 - 5 on page 23 shows the debug files that can be produced by the compiler at different stages of the compilation. Several of the output formats also contain debugging information in addition to the code and data.

The code generator optionally performs one other task: optimization. HI-TECH compilers come with several different optimizer stages. The code generator is responsible for *global optimization* which can be enabled using a `-Zg` (*global optimization*) option. This optimization is performed on the parsed source. Amongst other things, this optimization stage allocates variables to registers whenever possible and looks for constants that are used consecutively in source code to avoid reloading these values unnecessarily.

Assembly files are the first files in the compilation process that make reference to *psects*, or program sections. The code generator will generate the psect directives in which code and data will be positioned. The output of the code generator is then passed to the *assembler* which converts the ASCII representation of the processor instructions - the ASCII mnemonics - to binary *machine code*. The assembler is specific for each compiler and has a processor-dependent name such as ASPIC or ASXA. Assembler code also contains *assembler directives* which will be executed by the assembler. Some of these directives are to define ROM-based constants, others define psects and others declare global symbols.

The assembler is optionally preceded by an optimization of the generated assembler. This is the *peephole optimization*. With some HI-TECH compilers the peephole optimizer is contained in the assembler itself, e.g. the PIC assembler, however others have a separate optimization application which is run before the assembler is executed, e.g. OPT51. Peephole optimization is carried out separately over the assembler code derived from each single function.
In addition to the peephole optimizer, the assembler itself may include a separate assembler optimizer step which attempts to replace long branches with short branches where possible. The -O option enables both assembler optimizers, even if they are performed by separate applications, however HPD includes menu items for both optimizer stages (Peephole optimization and Assembler optimization). If the peephole optimizer is part of the assembler, the assembler optimization item in HPD has no effect.

The output of the assembler is an object file. An object file is a formatted binary file which contains machine code, data and other information relating to the module from which it has been generated. Object files come in two basic types: relocatable and absolute object files. Although both contain machine code in binary form, relocatable object files have not had their addresses resolved to be absolute values. The binary machine code is stored as a block for each psect. Any addresses in this area are temporarily stored as 00h. Separate relocation information in the object file indicates where these unresolved addresses lie in the psect and what they represent. Object files also contain information regarding any psects that are defined within so that the linker may position these correctly.

**Table 2 - 8 Assembler output**

<table>
<thead>
<tr>
<th>C source</th>
<th>Relocatable object file</th>
</tr>
</thead>
<tbody>
<tr>
<td>#define VAL 2</td>
<td>11 TEXT 22</td>
</tr>
<tr>
<td>int a, b;</td>
<td>text 0 13</td>
</tr>
<tr>
<td>void main(void)</td>
<td>99 08 00 00 88 10 A9 12 8E 00 00 D6 80</td>
</tr>
<tr>
<td>{</td>
<td>12 RELOC 63</td>
</tr>
<tr>
<td>/* set start...</td>
<td>2 RPSECT data 2</td>
</tr>
<tr>
<td>a = b + VAL;</td>
<td>9 COMPLEX 0</td>
</tr>
<tr>
<td>}</td>
<td>Key: direct</td>
</tr>
<tr>
<td></td>
<td>0x7&gt;==(high bss)</td>
</tr>
<tr>
<td></td>
<td>9 COMPLEX 1</td>
</tr>
<tr>
<td></td>
<td>((high bss)&amp;0x7)+0x8</td>
</tr>
<tr>
<td></td>
<td>10 COMPLEX 1</td>
</tr>
<tr>
<td></td>
<td>low bss</td>
</tr>
</tbody>
</table>

Object files produced by the assembler follow a format which is standard for all HI-TECH compilers, but obviously their contents are machine specific. Table 2 - 8 on page 26 shows several sections of the HI-TECH format relocatable object file that has been converted to ASCII for presentation using the DUMP executable which comes with the compiler. The highlighted source line is represented by the highlighted machine code in the object file. This code is positioned in a psect called text. The underlined bytes in the object file are addresses that as yet are unknown and have been replaced with zeros. The lines after the text psect in the object file show the information used to resolve the addresses needed by the linker. The two bytes starting at offset 2 and the two single bytes at offset 9 and 10 are represented here and as can be seen, their address will be contained at an address derived from the position of the data and bss psects, respectively.
If a -ASMLIST (Generate assemble listing) option was specified, the assembler will generate an assembler listing file which contains both the original C source lines and the assembler code that was generated for each line. The assembler listing output is shown in Table 2 - 9 on page 27. Unresolved addresses are listed as being zero with unresolved-address markers "/'" and "*'" used to indicate that the values are not absolute. Note that code is placed starting from address zero in the new text psect. The entire psect will be relocated by the linker..

Table 2 - 9 Assembler listing

<table>
<thead>
<tr>
<th>C source</th>
<th>Assembler listing</th>
</tr>
</thead>
<tbody>
<tr>
<td>#define VAL 2</td>
<td>10 0000' psect text</td>
</tr>
<tr>
<td></td>
<td>11 0000'</td>
</tr>
<tr>
<td></td>
<td>_main:</td>
</tr>
<tr>
<td>int a, b;</td>
<td>12</td>
</tr>
<tr>
<td></td>
<td>;main.c: 9: a = b + 2;</td>
</tr>
<tr>
<td>void</td>
<td>13 0000' 99 08 0000' mov.w r0,#_b</td>
</tr>
<tr>
<td>main(void)</td>
<td>14 0004' 88 10 movc.w rl,[r0+]</td>
</tr>
<tr>
<td></td>
<td>15 0006' A9 12 adds.w rl,#2</td>
</tr>
<tr>
<td></td>
<td>16 0008' 8E 00* 00* mov.w _a,rl</td>
</tr>
</tbody>
</table>
|           | 17 /* set start...
|           | 18 000B' D6 80 ret |

Some HI-TECH assemblers also generate a relocatable listing file (extension: .rlf). This contains address information which can be read by the linker and used to update the assembler listing file, if such a file was created. After linking, the assembler listing file will have unresolved addresses and address markers removed and replaced with their final absolute addresses.

The above series of steps: pre-processing, parsing, code generation and assembly, are carried out for each C source file passed to the driver in turn. Errors in the code are reported as they are detected. If a file cannot be compiled due to an error, the driver halts compilation of that module after the application that generated the error completes and continues with the next file which was passed to it, starting again with the CLIST application.

For any assembler files passed to the driver, these do not require as much processing as C source files, but they must be assembled. The compiler driver will pass any .as files straight to the assembler. If the user specifies the -P (Pre-process assembler files) the assembler files are first run through the C preprocessor allowing the using of all preprocessor directives within assembly code. The output of the preprocessor is then passed to the assembler.

Object and library files passed to the compiler are already compiled and are not processed at all by the first stages of the compiler. They are not used until the link stage which is explained below.

3. The generation of this file is not shown in Figure 2 - 1 on page 19 in the interests of clarity.
If you are using HPD, *dependency information* can be saved regarding each source and header file by clicking the **save dependency information** switch. When enabled, the HPD driver determines only which files in the project need be re-compiled from the modification dates of the input source files. If the source file has not been changed, the existing object file is used.

### 2.1.2.2 The link stage

The format of relocatable object files are again processor-independent so the linker and other applications discussed below are common across the whole range of HI-TECH compilers. The linker's name is **HLINK**.4

The tasks of the linker are many. The linker is responsible for combining all the object and library files into a single file. The files operated on by the linker include all the object files compiled from the input C source files and assembler files, plus any object files or library files passed to the compiler driver, plus any run-time object files and library files that the driver supplies. The linker also performs **grouping** and **relocation** of the psects contained in all of the files passed to it, using a relatively complex set of linker options. The linker also resolves symbol names to be absolute addresses after relocation has made it possible to determine where objects are to be stored in ROM or RAM. The linker then adjusts references to these symbols - a process known as address **fixup**. If the symbol address turns out to be too large to fit into the space in the instruction generated by the code generator, a **fixup overflow** error occurs. For example, if the address of the symbol \_b in our running example was determined to be 20000h, the linker would not be able to fit this address into the first underlined two byte "hole" in the object file shown dumped in the Table Assembler output on page 26 since 20000h is larger than two bytes long.

The linker can also generate a map file which has detailed information regarding the position of the psects and the addresses assigned to symbols. The linker may also produce a symbol file. These files have a .sym extension and are generated when the -G (**Source level debug info**) option is used. This symbol file is ASCII-based and contains information for the entire program. Addresses are absolute as this file is generated after the link stage.

Although the object file produced by HLINK contains all the information necessary to run the program, the program has to be somehow transferred from the host computer to the embedded hardware. There are a number of standard formats that have been created for such a task. Emulators and chip programmers often can accept a number of these formats. The *Motorola HEX* (S record) or *Intel HEX* formats are common formats. These are ASCII formats allowing easy viewing by any text editor. They include **checksum** information which can be used by the program which downloads the file to ensure that it was transmitted without error. These formats include address information which allows those areas which do not contain data to be omitted from the file. This can make these files significantly smaller than, for example, a binary file.

---

4. Early HI-TECH linkers were called **link**.
The `OBJTOHEX` application is responsible for producing the output file requested by the user. It takes the absolute object file produced by the linker and produces an output under the direction of the compiler driver. The `OBJTOHEX` application can produce a variety of different formats to satisfy most development systems. The output types available with most HI-TECH compilers are shown in Table 2-10.

In some circumstances, more than one output file is required. In this case an application called `CROMWELL`, the reformatter, is executed to produce further output files. For example it is commonly used with the PIC compiler to read in the HEX file and the SYM file and produce a COD file.

2.2 **Psects and the linker**

This tutorial explains how the compiler breaks up the code and data objects in a C program into different parts and then how the linker is instructed to position these into the ROM and RAM on the target.

2.2.1 **Psects**

As the code generator progresses it generates an assembler file for each C source file that is compiled. The contents of these assembly files include different sections: some containing assembler instructions...
that represent the C source; others contain assembler directives that reserve space for variables in RAM; others containing ROM-based constants that have been defined in the C source; and others which hold data for special objects such as variables to be placed in non-volatile areas, interrupt vectors and configuration words used by the processor. Since there can be more than one input source file there will be similar sections of assembler spread over multiple assembler files which need to be grouped together after all the code generation is complete.

These different sections of assembler need to be grouped in special ways: It makes sense to have all the initialised data values together in contiguous blocks so they can be copied to RAM in one block move rather than having them scattered in-between sections of code; the same applies to uninitialised global objects which have to be allocated a space which is then cleared before the program starts; some code or objects have to be positioned in certain areas of memory to conform to requirements in the processor’s addressing capability; and at times the user needs to be able to position code or data at specific absolute addresses to meet special software requirements. The code generator must therefore include information which indicates how the different assembler sections should be handled and positioned by the linker later in the compilation process.

The method used by the HI-TECH compiler to group and position different parts of a program is to place all assembler instructions and directives into individual, relocatable sections. These sections of a program are known as psects - short for program sections. The linker is then passed a series of options which indicate the memory that is available on the target system and how all the psects in the program should be positioned in this memory space.

2.2.1.1 The psect directive

The PSECT assembler directives (generated by the code generator or manually included in other assembly files) define a new psect. The general form of this directive is shown below.

\[ \text{PSECT name, option, option...} \]

It consists of the token PSECT followed by the name by which this psect shall be referred. The name can be any valid assembler identifier and does not have to be unique. That is, you may have several psects with the same name, even in the same file. As will be discussed presently, psects with the same name are usually grouped together by the linker.

The directive options are described in the assembler section of the manual, but several of these will be discussed in this tutorial. The options are instructions to the linker which describe how the psect should be grouped and relocated in the final absolute object file.

Psects which all have the same name imply that their content is similar and that they should be grouped and linked together in the same way. This allows you to place objects together in memory even if they are defined in different files.
After a psect has been defined, the options may be omitted in subsequent psect directives in the same module that use the same name. The following example shows two psects being defined and filled with code and data.

```
PSECT text,global
begin:
    mov   R0,#10
    mov   R2,r4
    add   R2,#8
PSECT data
input:
    DS  8
PSECT text
next:
    mov  r4,r2
    rrc  r4
```

In this example, the psect text is defined including an option to say that this is a global psect. Three assembler instructions are placed into this psect. Another psect is created: data. This psect reserves 8 bytes of storage space for data in RAM. The last psect directive will continue adding to the first psect. The options were omitted from the second PSECT directive in this example as there has already been a psect directive in this file that defines the options for a psect of this name. The above example will generate two psects. Other assembler files in the program may also create psects which have the same name as those here. These will be grouped with the above by the linker in accordance with the PSECT directive flags.

### 2.2.1.2 Psect types

Psects can be linked in three different ways: those that will reside permanently in ROM\(^5\); those that will be allocated space in RAM after the program starts; and those that will reside in ROM, but which will be copied into another reserved space in RAM after the program starts. A combination of code - known as the run-time startup code - and psect and linker options allow all this to happen.

Typically, psects placed into ROM contain instructions and constant data that cannot be modified. Those psects allocated space in RAM only are for global data objects that do not have to assume any non-zero value when the program starts, i.e. they are uninitialised. Those psects that have both a ROM image and space reserved in RAM are for modifiable, global data objects which are initialised, that is they contain some specific value when the program begins, but that value can be changed by the program during its execution.

---

\(^5\) The term "ROM" will be used to refer to any non-volatile memory.
Objects that are initialised are usually placed into psects with the name "data" or a name based on "data". Variables that are qualified near typically use the psect rdata. The PIC data psects use names like rdata_0 to indicates that they are "near" (there is no near qualifier on the PIC - essentially all PIC objects are near by default) and the digit indicates a bank number.

Uninitialised objects are placed in psects whose name is "bss" or a name based on "bss". Again, rbss would indicate uninitialised objects that are near. The PIC compiler uses names like rbss_0, where the digit is a bank number. The abbreviation "bss" stands for block started by symbol and was a assembler pseudo-op used in IBM systems back in the days when computers were coal-fired. The continued usage of this term is still appropriate as there are some similarities in the way these schemes worked.

The following C source shows two objects being defined. The object input will be placed into a data psect; the value 22 will reside in ROM and be copied to the RAM space allocated for input by the run-time code. The object output will not contribute directly to the ROM image. A area of memory will be reserved for it in RAM and this area will be cleared by the run-time code (output will be assigned the value 0).

```c
int input = 22;  // an initialised object
int output;      // an uninitialised object
```

Snippets from the assembler listing file show how the XA compiler handles these two objects. Other compilers may produce differently structured code. The PSECT directive flags are discussed presently, but note that for the initialised object, input, the code generator used a DW (define word) directive which placed the two bytes of the int value (16 and 00) into the output which is destined for the ROM. Two bytes of storage were reserved using the DS assembler directive for the uninitialised object, output, and no values appear in the output.

```
1  0000'          PSECT data, class=CODE, space=0, align=0
2                                    GLOBAL _input
3                                    ALIGN.W
4  0000'          _input:
5  0000' 16 00    DW 22

13  0000'          PSECT bss, class=DATA, space=1, align=0
14                                    GLOBAL _output
15                                    ALIGN.W
16  0000'          _output:
17  0000'          DS 2
```
Auto variables and function parameters are local to the function in which they are defined and are handled differently by the compilers. They may be allocated space dynamically (for example on the stack) in which case they are not stored in psects by the compiler. Compilers or memory models which do not use a hardware stack, use a compiled stack which is an area of memory set aside for the auto and parameter objects for each function. These objects will be positioned in a psect. The psect in which they are allocated is defined by a `FNCONF` directive which is placed in the run-time startup code.

Two addresses are used to refer to the location of a psect: the link address and the load address. The link address is the address at which the psect (and any objects or labels within the psect) can be accessed whilst the program is executing. The load address is the address at which the psect will reside in the output file that creates the ROM image, or, alternatively, the address of where the psect can be accessed in ROM.

For the psect types that reside in ROM their link and load address will be the same, as they are never copied to a new location. Psects that are allocated space in RAM only will have a link address, but a load address is not applicable. They are assigned a load address, and you will see it listed in the map file, but it is not used. The compiler usually makes the load address of these psects the same as the link address. Since no ROM image of these psects is formed, the load address is meaningless and can be ignored. Any access to objects defined in these psects is performed using the link address. The psects that reside in ROM, but are copied to RAM have link and load addresses that are usually different. Any references to symbols or labels in these psects are always made using the link addresses.

### 2.3 Linking the psects

After the code generator and assembler\(^6\) have finished their jobs, the object files passed to the linker can be considered to be a mixture of psects that have to be grouped and positioned in the available ROM and RAM. The linker options indicate the memory that is available and the flags associated with a `PSECT` directive indicate how the psects are to be handled.

#### 2.3.1 Grouping psects

There are two `PSECT` flags that affect the grouping, or merging, of the psects. These are the `local` and `global` flags. `Global` is the default and tells the linker that the psects should be grouped together with other psects of the same name to form a single psect. `Local` psects are not grouped in this way unless they are contained in the same module. Two `local` psects which have the same name, but which are defined in different modules are treated and positioned as separate psects.

\(^6\) The assembler does not modify `PSECT` directives in any way other than encoding the details of each into the object file.
2.3.2 Positioning psects

Several PSECT flags affect how the psects are positioned in memory. Psects which have the same name can be positioned in one of two ways: they can be overlaid one another, or they can be placed so that each takes up a separate area of memory.

Psects which are to be overlaid will use the ovlrd flag. At first it may seem unusual to have overlaid psects as they might destroy other psects’ contents as they are positioned, however there are instances where this is desirable.

One case where overlaid psect are used is when the compiler has to use temporary variables. When the compiler has to pass several data objects to, say, a floating-point routine, the floats may need to be stored in temporary variables which are stored in RAM. It is undesirable to have the space reserved if it is not going to be used, so the routines that use the temporary objects are also responsible for defining the area and reserving the space in which these will reside. However several routines may be called and hence several temporary areas created. To get around this problem, the psects which contain the directives to reserve space for the objects are defined as being overlaid so that if more than one is defined, they simply overlap each other.

Another situation where overlaid psects are used is when defining the interrupt vectors. The run-time code usually defines the reset vector, but other vectors are left up to the programmer to initialize. Interrupt vectors are placed into a separate psect (often called vectors). Each vector is placed at an offset from the beginning of the vectors area appropriate for the target processor. The offset is achieved via an ORG assembler directive which moves the location counter relative to the beginning of the current psect. The macros contained in the header file <intrpt.h>, which allow the programmer to define additional interrupt vectors, also place the vectors they define into a psect with this same name, but with different offsets, depended on the interrupt vector being defined. All these psects are grouped and overlaid such that the vectors are correctly positioned from the same address - the start of the vectors psect. This merged psect is then positioned by the linker so that it begins at the start of the vectors area.

Most other compiler-generated psects are not overlaid and so they will each occupy their own unique address space. Typically these psects are placed one after the other in memory, however there are several PSECT flags that can alter the positioning of the psects. Some of these PSECT flags are discussed below.

The reloc flag is used when psects must be aligned on a boundary in memory. This boundary is a multiple of the value specified with the flag. The abs flag specifies that the psect is absolute and that it should start at address 0h. Remember, however, that if there are several psects which use this flag, then after grouping only one can actually start at address 0h unless the psects are also defined to be overlaid. Thus abs itself means that one of the psects with this name will be located at address 0h, the others following in memory subject to any other psect flags used.
2.3.3 Linker options to position p sects

The linker is told of the memory setup for a target program by the linker options that are generated by the compiler driver. The user informs the compiler driver about memory using either the -A option\(^7\) with the command line driver (CLD), or via the ROM & RAM addresses dialogue box under HPD. Additional linker options indicate how the psects are to be positioned into the available memory.

The linker options are a little confusing at first, but the following example shows how the options could be built up as a program develops, and then discusses some of the specific schemes used by HI-TECH compilers. When compiling using either the CLD or HPD, a full set of default linker options are used, based on either the -A option values, or the ROM & RAM addresses dialogue values. In most cases the linker options do not need to be modified.

2.3.3.1 Placing psects at an address

Let us assume that the processor in a target system can address 64 kB of memory and that ROM, RAM and peripherals all share this same block of memory. The ROM is placed in the top 16 kB of memory (C000h - FFFFh); RAM is placed at addresses from 0h to FFFh.

Let us also assume that three object files passed to the linker: one a run-time object file; the others compiled from the programmer's C source code. Each object file contains a compiler-generated text psect (i.e. a psect called text): the psect in one file is 100h bytes long; that from other file is 200h bytes long; that from the run-time object file is 50h bytes long. These psects are to be placed in ROM and all have the simple definition generated by the code generator:

\[
\text{PSECT text, class=CODE}
\]

The class flag is typically used with these types of psects and is considered later in this tutorial. By default, these psects are global, hence after scanning all the object files passed to it, the linker will group all the text psects together so that they are contiguous\(^8\) and form one larger text psect. The following \(-p\) linker option could be used to position the text psect at the bottom of ROM.

\[-p\text{text}=0C000h\]

There is only one address specified with this linker option since the psects containing code are not copied from ROM to RAM at any stage and the link and load addresses are the same.

---

7. \(\) The \(-A\) option on the PIC compiler serves a different purpose. Most PIC devices only have internal memory and so a memory option is not required by the compiler. High-end PICs may have external memory, this is indicated to the compiler by using a \(-\text{ROM}\) option to the CLD or by the ROM & RAM addresses... dialogue box under HPDPIC. The \(-A\) option is used to shift the entire ROM image, when using high-end devices.

8. Some processors may require word alignment gaps between code or data. These gaps can be handled by the compiler, but are not considered here.
The linker will relocate the grouped `text` psect so that it starts at address C000h. The linker will then define two global symbols with names: `__Ltext` and `__Htext` and equate these with the values: C000h and C350h which are the start and end addresses, respectively, of the `text` psect group.

Now let us assume that the run-time file and one of the programmer's files contains interrupt vectors. These vectors must be positioned at the correct location for this processor. Our fictitious processor expects its vectors to be present between locations FFC0h and FFFFh. The reset vector takes up two bytes at address FFFEh an FFFFh, and the remaining locations are for peripheral interrupt vectors. The run-time code usually defines the reset vector using code like the following.

```assembly
GLOBAL start
PSECT vectors, ovlrd
ORG 3Eh
DW start
```

This assembler code creates a new psect which is called `vectors`. This psect uses the overlaid flag (`ovlrd`) which tells the linker that any other psects with this name should be overlaid with this one, not concatenated with it. Since the psect defaults to being `global`, even `vectors` psects in other files will be grouped with this one. The `ORG` directive tells the assembler to advance 3Eh locations into this psect. It does not tell the assembler to place the object at address 3Eh. The final destination of the vector is determined by the relocation of the psect performed by the linker later in the compilation process. The assembler directive `DW` is used to actually place a word at this location. The word is the address of the (global) symbol `start`. (`start` or `powerup` are the labels commonly associated with the beginning of the run-time code.)

In one of the user's source files, the macro `ROM_VECTOR()` has been used to supply one of the peripheral interrupts at offset 10h into the vector area. The macro expands to the following in-line assembler code.

```assembly
GLOBAL_timer_isr
PSECT vectors, ovlrd
ORG 10h
DW _timer_isr
```

After the first stages of the compilation have been completed, the linker will group together all the `vectors` psects it finds in all the object files, but they will all start from the same address, i.e. they are all placed one over the other. The final `vectors` psect group will contain a word at offset 10h and another at offset 3Eh. The space from 0h to offset 0Fh and in-between the two vectors is left untouched by the linker. The linker options required to position this psect would be:

```
-pvectors=0FFC0h
```

The address given with this option is the base address of the vectors area. The `ORG` directives used to move within the `vectors` psects hence were with respect to this base address.
Both the user's files contain constants that are to be positioned into ROM. The code generator generates the following **PSECT** directive which defines the psect in which it store the values.

```
PSECT const
```

The linker will group all these **const** psects together and they can be simply placed like the **text** psects. The only problem is: where? At the moment the **text** psects end at address C34Fh so we could position the **const** psects immediately after this at address C350h, but if we modify the program, we will have to continually adjust the linker options. Fortunately we can issue a linker option like the following.

```
-ptext=0C000h,const
```

We have not specified an address for the psect **const**, so it defaults to being the address immediately after the end of the preceding psect listed in the option, i.e. the address immediately after the end of the **text** psect. Again, the **const** psect resides in ROM only, so this one address specifies both the link and load addresses.

Now the RAM psects. The user's object files contain uninitialised data objects. The code generator generates **bss** psects in which are used to hold the values stored by the uninitialised C objects. The area of memory assigned to the **bss** psect will have to be cleared before **main()** is executed.

At link time, all **bss** psects are grouped and concatenated. The psect group is to be positioned at the beginning of RAM. This is easily done via the following option.

```
-pbss=0h
```

The address 0h is the psect's link address. The load address is meaningless, but will default to the link address. The run-time code will clear the area of memory taken up by the **bss** psect. This code will use the symbols **__Lbss** and **__Hbss** to determine the starting address and the length of the area that has to be cleared.

Both the user's source files contain initialised objects like the following.

```
int init = 34;
```

The value 34 has to be loaded into the object **init** before the **main()** starts execution. Another of the tasks of the run-time code is to initialise these sorts of objects. This implies that the initial values must be stored in ROM for use by the run-time code. But the object is a variable that can be written to, so it must be present in RAM once the program is running. The run-time code must then copy the initialised values from ROM into RAM just before **main()** begins. The linker will place all the initial values into ROM in exactly the same order as they will appear in RAM so that the run-time code can simply copy the values from ROM to RAM as a single block. The linker has to be told where in ROM these values should reside as it generates the ROM image, but is must also know where in RAM these objects will be copied to so that it can leave enough room for them and resolve the run-time addresses for symbols in this area.
The complete linker options for our program, including the positioning of the data psects, might look like:

- \texttt{-ptext=0C000h,const}
- \texttt{-pvectors=0FFC0h}
- \texttt{-pbss=0h,data/const}

That is, the data psect should be positioned after the end of the bss psect in RAM. The address after the slash indicates that this psect will be copied from ROM and that its position in ROM should be immediately after the end of the const psect. As with other psects, the linker will define symbols \texttt{__Ldata} and \texttt{__Hdata} for this psect, which are the start and end link addresses, respectively, that will be used by the run-time code to copy the data psect group. However with any psects that have different link and load addresses, another symbol is also defined, in this case: \texttt{__Bdata}. This is the load address in ROM of the data psect.

2.3.3.2 Exceptional cases

The PIC compiler handles the data psects in a slightly different manner. It actually defines two separate psects: one for the ROM image of the data psects; the other for the copy in RAM. This is because the length of the ROM image is different to the length of the psect in RAM. (The ROM is wider than the RAM and values stored in ROM may be encoded as \texttt{retlw} instructions.) Other compilers may also operate this way if ROM and RAM are in different memory spaces. The linker options in this case will contain two separate entries for both psects instead of the one psect with different link and load addresses specified. The names of the data psects in RAM for the PIC compiler will be similar to \texttt{rdata_0}; those in ROM are like \texttt{idata_0}. The digit refers to a PIC RAM bank number.

The link and load addresses reported for psects that contain objects of type \texttt{bit} have slightly different meaning to ordinary link and load addresses. In the map file, the link address listed is the link address of the psect specified as a bit address. The load address is the link address specified as a byte address. Bit objects cannot be initialised, so separate link and load addresses are not required. The linker knows to handle these psects differently because of the \texttt{bit} psect flag. Bit psects will be reported in the map file as having a \texttt{scale} value of 8. This relates to the number of objects that can be positioned in an addressable unit.

2.3.3.3 Psect classes

Now let us assume that more ROM is added to our system since the programmers have been busy and filled the 16 kB currently available. Several peripheral devices were placed in the area from B000h to BFFFFh so the additional ROM is added below this from 7000h to AFFFh. Now there are two separate areas of ROM and we can no longer give a single address for the text psects.

What we can now do to take advantage of this extra memory is define several ranges of addresses that can be used by ROM-based psects. This can be done by creating a psect \textit{class}. There are several ways that psects can be linked when using classes. Classes are commonly used by HI-TECH C compilers to position the code or text psects. Different strategies are employed by different compilers to better suit
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the processor architecture for which the compilation is taking place. Some of these schemes are discussed below. If you intend to modify the default linker options or generate your own psects, check the linker options and PSECT directives generated by the code generator for the specific compiler you are using.

A class can be defined using another linker option. For example to use the additional memory added to our system we could define a class using the linker option:

\[-ACODE=7000h-AFFFh,C000h-FFFFh\]

The option is a -A immediately followed by the class name and then a comma-separated list of address ranges. Remember this is an option to the linker, not the CLD. The above example defines two address ranges for a class called CODE.

Here is how drivers for the 8051, 8051XA and Z80 compilers define the options passed to the linker to handle the class CODE psects. In large model the 8051 psect definitions for psects that contain code are as follows.

\[\text{PSECT text, class=CODE}\]

The class psect flag specifies that the psect text is a member of the class called CODE.

If a single ROM space has been specified by either not using the -ROM option with the CLD or by selecting single ROM in the ROM & RAM addresses dialogue box under HPD, no class is defined and the psects are linked using a -p option as we have been doing above. Having the psects within a class, but not having that class defined is acceptable, provided that there is a -p option to explicitly position the psects after they have been grouped. If there is no class defined and no -p option a default memory address is used which will almost certainly be inappropriate.

If multiple ROM spaces have been specified by using either the -ROMranges option with the CLD, or specifying the address ranges in the ROM & RAM addresses (after selecting the multiple ROMs button) dialogue box under HPD, a class is defined by the driver using the -A linker option similar to that shown above and the -p option is omitted from the options passed to the linker.

The PIC compiler does things a little differently as it has to contend with multiple ROM pages that are quite small. The PIC code generator defines the psects which hold code like the following.

\[\text{PSECT text0,local,class=CODE, delta=2}\]

The delta value relates to the ROM width and need not be considered here. The psects are placed in the CODE class, but note that the they are made local using the local psect flag. The psects that are generated from C functions each have unique names which proceed: text0, text1, text2 etc. Local psects are not grouped across modules, i.e. if there are two modules, each containing a local psect of the same name, they are treated are separate psects. Local psects cannot be positioned using a -p linker option as there can be more than one psect with that name. Local psects must be made members of a
class, and the class defined using a -A linker option. The PIC works in this way to assist with the placement of the code in its ROM pages. This is discussed further in Section 2.3.4 on page 42.

A few general rules apply when using classes: If, for example, you wanted to place a psect that is not already in a class into the memory that a class occupies, you can replace an address or psect name in a linker -p option with a class name. For instance, in the generic example discussed above, the const psect was placed after the text psect in memory. If you would now like this psect to be positioned in the memory assigned to the CODE class the following linker options could be used.

```plaintext
-pconst=CODE
-pvectors=0FFC0h
-pbss=0h,data/CODE
-ACODE=7000h-FFFh,C000h-FFFFh
```

Note also that the data psect’s load location has been swapped from after the end of the const psect to within the memory assigned to the CODE class to illustrate that the load address can be specified using the class name.

Another class definition that is sometimes seen in PIC linker options specifies three addresses for each memory range. Such an option might look something like:

```plaintext
-ENTRY=0h-FFh-1FFh
```

The first range specifies the address range in which the psect must start. The psects are allowed to continue past the second address as long as they do not extend past the last address. For the example above, all psects that are in the ENTRY class must start at addresses between 0 and FFh. The psects must end before address 1FFh. No psect may be positioned so that its starting address lies between 100h and 1FFh. The linker may, for example, position two psects in this range: the first spanning addresses 0 to 4Fh and the second starting at 50h and finishing at 138h. Such linker options are useful on some PIC processors (typically baseline PICs) for code psects that have to be accessible to instructions that modify the program counter. These instructions can only access the first half of each ROM page.

### 2.3.3.4 User-defined psects

Let us assume now that the programmer wants to include a special initialised C object that has to be placed at a specific address in memory, i.e. it cannot just be placed into, and linked with, the data psect. In a separate source file the programmer places the following code.

```plaintext
#pragma psect data=lut
int lookup_table[] = {0, 2, 4, 7, 10, 13, 17, 21, 25};
```

The pragma basically says, from here onwards in this module, anything that would normally go into the data psect should be positioned into a new psect called lut. Since the array is initialised, it would normally be placed into data and so it will be re-directed to the new psect. The psect lut will inherit any psect options (defined by the PSECT directive flags) which applied to data.
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The array is to be positioned in RAM at address 500h. The -p option above could be modified to include this psect as well.

\[-pbss=0h, data/const, lut=500h/\]

(The load address of the data psect has been returned to its previous setting.) The addresses for this psect are given as 500h/. The address 500h specifies the psect’s link address. The load address can be anywhere, but it is desirable to concatenate it to existing psects in ROM. If the link address is not followed by a load address at all, then the link and load addresses would be set to be the same, in this case 500h. The "/", which is not followed by an address, tells the linker that the load address should be immediately after the end of the previous psect’s load address in the linker options. In this case that is the data psect’s load address, which in turn was placed after the const psect. So, in ROM will be placed the const, data and lut psects, in that order.

Since this is an initialised data psect, it is positioned in ROM and must be copied to the memory reserved for it in RAM. Although different link and load addresses have been specified with the linker option, the programmer will have to supply the code that actually performs the copy from ROM to RAM. (The data psects normally created by the code generator have code already supplied in the run-time file to copy the psects.) The following is C code which could perform the copy.

```c
extern unsigned char *__Llut, *__Hlut, *__Blut;
unsigned char *i;

void copy_my_psect(void)
{
    for(i=__Llut; i<_Hlut; i++, _Blut++)
        *i = *__Blut;
}
```

Note that to access the symbols __Llut etc. from within C code, the first underscore character is dropped. These symbols hold the addresses of psects, so they are declared (not defined) as pointer objects in the C code using the `extern` qualifier. Remember that the object lookuputable will not be initialised until this C function has been called and executed. Reading from the array before it is initialized will return incorrect values.

If you wish to have initialised objects copied to RAM before main() is executed, you can write assembler code, or copy and modify the appropriate routine in the run-time code that is supplied with the compiler. You can create your own run-time object file by pre-compiling the modified run-time file and using this object file instead of the standard file that is automatically linked with user’s programs. From assembler, both the underscore characters are required when accessing the psect address symbols.
If you define your own psect based on a bss psect, then, in the same way, you will have to supply code to clear this area of memory if you are to assume that the objects defined within the psect will be cleared when they are first used.

2.3.4 Issues when linking

The linker uses a relatively complicated algorithm to relocate the psects contained in the object and library files passed to it, but the linker needs more information than that discussed above to know exactly how to relocate each psect. This information is contained in other linker options passed to the linker by the driver and in the psect flags which are used with each PSECT directive. The following explain some of the issues the linker must take into account.

2.3.4.1 Paged memory

Let’s assume that a processor has two ROM areas in which to place code and constant data. The linker will never split a psect over any memory boundary. A memory boundary is assumed to exist wherever there is a discontinuity in the address passed to the linker in the linker options. For example, if a class is specified using the addresses as follows:

    -ADATA=0h-FFh,100h-1FFh

It is assumed that some boundary exists between address FFh and 100h, even though these addresses are contiguous. This is why you will see contiguous address ranges specified like this, instead of having one range covering the entire memory space. To make it easy to specify similar contiguous address ranges, a repeat count can be used, like:

    -ADATA=0h-FFhx2

can be used; in this example, two ranges are specified: 0 to FFh and then 100h to 1FFh. Some processors have memory pages or banks. Again, a psect will not straddle a bank or page boundary.

Given that psects cannot be split over boundaries, having large psects can be a problem to relocate. If there are two, 1 kB areas of memory and the linker has to position a single 1.8 kB psect in this space, it will not be able to perform this relocation, even though the size of the psect is smaller than the total amount of memory available. The larger the psects, the more difficult it is for the linker to position them. If the above psect was split into three 0.6 kB psects, the linker could position two of them - one in each memory area - but the third would still not fit in the remaining space in either area. When writing code for processors like the PIC, which place the code generated from each C function into a separate, local psect, functions should not become too long.

If the linker cannot position a psect, it generates a Can’t find space for psect xxxx error, where xxxx is the name of the psect. Remember that the linker relocates psects so it will not report memory errors with specific C functions or data objects. Search the assembler listing file to identify which C function is associated with the psect that is reported in the error message if local psects are generated by the code generator.
Global psects that are not overlaid are concatenated to form a single psect by the linker before relocation takes place. There are instances where this grouped psect appears to be split again to place it in memory. Such instances occur when the psect class within which it is a member covers several address ranges and the grouped psect is too large to fit any of the ranges. The linker may use intermediate groupings of the psects, called *clutches* to facilitate relocation within class address ranges. Clutches are in no way controllable by the programmer and a complete understanding of their nature is not required to able to understand or use the linker options. It is suffice to say that global psects can still use the address ranges within a class. Note that although a grouped psect can be comprised of several clutches, an individual psect defined in a module can never be split under any circumstances.

### 2.3.4.2 Separate memory areas

Another issue faced by the linker is this: On some processors, there are distinct memory areas for program and data, i.e. Harvard architecture chips like the XA. For example, ROM may extend from 0h - FFFFh and separate RAM may extend from 0h - 7FFh. If the linker is asked to position a psect at address 100h via a `-p` option, how does the linker know whether this is an address in program memory or in the data space? The linker makes use of the *space* psect flag to determine this. Different areas are assigned a different *space* value. For example ROM may be assigned a *space* value of 0 and RAM a *space* flag of 1. The *space* flags for each psect are shown in the map file.

The *space* flag is not used when the linker can distinguish the destination area of an object from its address. Some processors use memory banks which, from the processor’s point of view, cover the same range of addresses, but which are within the same distinct memory area. In these cases, the compiler will assign unique addresses to objects in banked areas. For example, some PIC processors can access four banks of RAM, each bank covering addresses 0 to 7Fh. The compiler will assign objects in the first bank (bank 0) addresses 0 to 7Fh; objects in the second bank: 80h to FFh; objects in the third bank: 100h to 17Fh etc. This extra bank information is removed from the address before it is used in an assembler instruction. All PIC RAM banks use a *space* flag of 1, but the ROM area on the PIC is entirely separate and uses a different *space* flag (0). The *space* flag is not relevant to psects which reside in both memory areas, such as the data psects which are copied from ROM to RAM.

After relocation is complete, the linker will group psects together to form a *segment*. Segments, along with clutches, are rarely mentioned with the HI-TECH compiler simply because they are an abstract object used only by the linker during its operation. Segment details will appear in the map file. A segment is a collection of psects that are contiguous and which are destined for a specific area in memory. The name of a segment is derived from the name of the first psect that appears in the segment and should not be confused with the psect which has that name.

### 2.3.4.3 Objects at absolute addresses

After the psects have been relocated, the addresses of data objects can be resolved and inserted into the assembler instructions which make reference to an object’s address. There is one situation where the linker does not determine and resolve the address of a C object. This is when the object has been defined
Tutorials

as absolute in the C code. The following example shows the object DDRA being positioned at address 200h.

```
unsigned char DDRA @ 0x200;
```

When the code generator makes reference to the object DDRA, instead of using a symbol in the generated assembler code which will later be replaced with the object’s address after psect relocation, it will immediately use the value 200h. The important thing to realise is that the instructions in the assembler that access this object will not have any symbols that need to be resolved, and so the linker will simply skip over them as they are already complete. If the linker has also been told, via its linker options, that there is memory available at address 200h for RAM objects, it may very well position a psect such that an object that resides in this psect also uses address 200h. As there is no symbol associated with the absolute object, the linker will not see that two objects are sharing the same memory. If objects are overlapping, the program will most likely fail unpredictably.

When positioning objects at absolute address, it vital to ensure that the linker will not position objects over those defined as absolute. Absolute objects are intended for C objects that are mapped over the top of hardware registers to allow the registers to be easily access from the C source code. The programmer must ensure that the linker options do not specify that there is any general-purpose RAM in the memory space taken up by the hardware. Ordinary variables to be positioned at absolute addresses should be done so using a separate psect (by simply defining your own using a PSECT directive in assembler code, or by using the #pragma psect directive in C code) and linker option to position the objects. If you must use an absolute address for an object in general-purpose RAM, make sure that the linker options are modified so that the linker will not position other psects in this area.

### 2.3.5 Modifying the linker options

In most applications, the default linker options do not need to be modified. It is recommended that if you think the options should be modified, but you do not understand how the linker options work, that you seek technical assistance in regard to the problem at hand.

If you do need to modify the linker options, there are several ways to do this. If you are simply adding another option to those present by default, the option can be specified to the CLD using a `-L` option. To position the lut psect that was used in the earlier example, the following option could be used.

```
-L-plut=500/const
```

The `-L` simply passes whatever follows to the linker. If you want to add another option to the default linker options and you are using HPD and a project, then it is a simple case of opening the linker options... dialogue box and adding the option to the end of those already there. The options should be entered exactly as they should be presented to the linker, i.e. you do not need the `-L` at the front.

If you want to modify existing linker options, other than simply changing the memory address that are specified with the `-A` CLD option, then you cannot use the CLD to do this directly. What you will need
to do is to perform the compilation and link separately. Let’s say that the file main.c and extra.c are to be compiled for the 8051 with modified linker options. First we can compile up to, but not include, the link stage by using a command line something like this.

\[ \text{c51 -O -Zg -ASMLIST -C main.c extra.c} \]

The \(-C\) options stops the compilation before the link stage. Include any other options which are normally required. This will create two object files: main.obj and extra.obj, which then have to be linked together.

Run the CLD again in verbose mode by giving a \(-V\) option on the command line, passing it the names of the object files created above, and redirect the output to a file. For example:

\[ \text{c51 -V -A8000,0,100,0,0 main.obj extra.obj > main.lnk} \]

Note that if you do not give the \(-A\) CLD option, the compiler will prompt you for the memory addresses, but with the output redirected, you will not see the prompts.

The file generated (main.lnk) will contain the command line that CLD generated to run the linker with the memory values specified using the \(-A\) option. Edit this file and remove any messages printed by the compiler. Remove the command line for any applications run after the link stage, for example OBJTOHEX or CROMWELL, although you should take note of what these command lines are as you will need to run these applications manually after the link stage. The linker command line is typically very long and if a DOS batch file is used to perform the link stage, it is limited to lines 128 characters long. Instead the linker can be passed a command file which contains the linker options only. Break up the linker command line in the file you have created by inserting \textit{backslash} characters "\" followed by a \textit{return}. Also remove the name and path of the linker executable from the beginning of the command line so that only the options remain. The above command line generated a main.lnk file that was then edited as suggested above to give the following.

\[ \text{-z -pvectors=08000h,text,code,data,const,strings \}
-\text{prbit=0/20h,rbss,rdata/strings,irdata,idata/rbss \}
-\text{pbss=0100h/idata -pnvram=bss,heap -ol.obj \}
-\text{m/tmp/06206eaa /usr/hitech/lib/rt51--ns.obj main.obj \}
\text{extra.obj /usr/hitech/lib/51--nsc.lib} \]

Now, with care, modify the linker options in this file as required by your application.

Now perform the link stage by running the linker directly and redirecting its arguments from the command file you have created.

\[ \text{hlink < main.lnk} \]

This will create an output file called l.obj. If other applications were run after the link stage, you will need to run them to generate the correct output file format, for example a HEX file.
Modifying the options to HPD is much simpler. Again, simply open the linker options... dialogue box and make the required changes, using the buttons at the bottom of the box to help with the editing. Save and re-make your project.

The map file will contain the command line actually passed to the linker and this can be checked to confirm that the linker ran with the new options.

2.4
PICC18 Command Line Driver

PICC18 is the driver invoked from the DOS command line to compile and/or link C programs. PICC18 has the following basic command format:

```
PICC18 [options] files [libraries]
```

It is conventional to supply the options (identified by a leading `dash` “-”) before the filenames, but in fact this is not essential.

The options are discussed below. The files may be a mixture of source files (C or assembler) and object files. The order of the files is not important, except that it will affect the order in which code or data appears in memory. Libraries are a list of library names, or `-L` options (see page 57). Source files, object files and library files are distinguished by PICC18 solely by the file type or extension. Recognized file types are listed in Table 3-1. This means, for example, that an assembler file must always have a `.as` extension (alphabetic case is not important).

### Table 3-1 PICC-18 File Types

<table>
<thead>
<tr>
<th>File Type</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>.c</td>
<td>C source file</td>
</tr>
<tr>
<td>.as</td>
<td>Assembler source file</td>
</tr>
<tr>
<td>.obj</td>
<td>Relocatable object code file</td>
</tr>
<tr>
<td>.lib</td>
<td>Relocatable object library file</td>
</tr>
</tbody>
</table>

PICC18 will check each file argument and perform appropriate actions. C files will be compiled; assembler files will be assembled. At the end, unless suppressed by one of the options discussed later, all object files resulting from compilation or assembly, or those listed explicitly on the command line, will be linked together with the standard run-time code and libraries and any user-specified libraries. Functions in libraries will be linked into the resulting output file only if referenced in the source code.

Invoking PICC18 with only object files specified as the file arguments (i.e. no source files) will mean only the link stage is performed. It is typical in Makefiles to use PICC18 with a `-C` option to compile several source files to object files, then to create the final program by invoking PICC18 again with only the generated object files and appropriate libraries (and appropriate options).

### 3.1 Long Command Lines

The PICC18 driver is a 32-bit Windows application, thus it is able to process command lines exceeding 128 characters in length. The driver may be called from within a DOS batch file or passed options via a
command file. When using batch files, the command line to PICC18 must be contained on one line. Driver options may be spread over multiple lines in a command file by using a space character followed by a backslash “\” followed by a return to split the lines. For example a command file may contain:

```
-V -O -18C452 -UBROF -D32 \nfile1.obj file2.obj mylib.lib
```

If this was in the file xyz.cmd then PICC18 would be invoked as:

```
PICC18 < xyz.cmd
```

Since no command line arguments were supplied, PICC18 will read xyz.cmd for its command line.

The command file may also be read by using the @ symbol. For example:

```
PICC18 @xyz.cmd
```

### 3.2 Default Libraries

PICC18 will search the appropriate standard C library by default for symbol definitions. This will always be done last, after any user-specified libraries. The particular library used will be dependent on the processor selected.

The standard library contains a version of `printf()` that supports only integer length values. If you want to print long values with `printf()`, or `sprintf()` or related functions, you must specify a `-Ll` option. This will search the library containing the long version of `printf()`. For floating-point and long `printf()` support, use the `-Lf` option which will search the library containing the floating-point version of `printf()`. You do not need the `-Ll` option if you have specified the `-Lf` option.

### 3.3 Standard Run-Time Code

PICC18 will also automatically link in the standard run-time module appropriate for the processor selected, unless you have specified the option to disable this, `-NORT`. If you require any special powerup initialization, rather than replace or modify the standard run-time module, you should use the `powerup` routine feature (see page 78).

### 3.4 PICC18 Compiler Options

The compiler is configured primarily for generation of ROM code. PICC18 recognizes the compiler options listed in Table 3 - 2 on page 49. The case of the options is not important, however UNIX shells are case sensitive when it comes to names of files. The ordering or configuration of some options are important.
## PICC18 Compiler Options

**Table 3 - 2 PICC18 Options**

<table>
<thead>
<tr>
<th>Option</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-processor</td>
<td>Define the processor</td>
</tr>
<tr>
<td>-AAHEX</td>
<td>Generate an American Automation symbolic HEX file</td>
</tr>
<tr>
<td>-ASMLIST</td>
<td>Generate assembler .LST file for each compilation</td>
</tr>
<tr>
<td>-Aaddress</td>
<td>Specify offset for ROM code</td>
</tr>
<tr>
<td>-A-option</td>
<td>Specify <code>-option</code> to be passed directly to the assembler</td>
</tr>
<tr>
<td>-BIN</td>
<td>Generate a Binary output file</td>
</tr>
<tr>
<td>-BL</td>
<td>Selects large memory model</td>
</tr>
<tr>
<td>-BS</td>
<td>Selects small memory model</td>
</tr>
<tr>
<td>-C</td>
<td>Compile to object files only</td>
</tr>
<tr>
<td>-CKfile</td>
<td>Make OBJTOHEX use a checksum file</td>
</tr>
<tr>
<td>-CP16</td>
<td>Use 16-bit wide pointers to program space</td>
</tr>
<tr>
<td>-CP24</td>
<td>Use 24-bit wide pointers to program space</td>
</tr>
<tr>
<td>-CRfile</td>
<td>Generate cross-reference listing</td>
</tr>
<tr>
<td>-D24</td>
<td>Use truncated 24-bit floating point format for doubles</td>
</tr>
<tr>
<td>-D32</td>
<td>Use IEEE754 32-bit floating point format for doubles</td>
</tr>
<tr>
<td>-Dmacro</td>
<td>Define preprocessor macro</td>
</tr>
<tr>
<td>-E</td>
<td>Define format for compiler errors</td>
</tr>
<tr>
<td>-Efile</td>
<td>Redirect compiler errors to a file</td>
</tr>
<tr>
<td>-E+file</td>
<td>Append errors to a file</td>
</tr>
<tr>
<td>-FAKELOCAL</td>
<td>Produce MPLAB-specific debug information</td>
</tr>
<tr>
<td>-FDDOUBLE</td>
<td>Enables the use of faster 32-bit floating point math routines</td>
</tr>
<tr>
<td>-Gfile</td>
<td>Generate enhanced source-level symbol table</td>
</tr>
<tr>
<td>-HELP</td>
<td>Print summary of options</td>
</tr>
<tr>
<td>-ICD</td>
<td>Compile code for MPLAB-ICD</td>
</tr>
<tr>
<td>-Ipath</td>
<td>Specify a directory pathname for include files</td>
</tr>
<tr>
<td>-INTEL</td>
<td>Generate an Intel HEX format output file (default)</td>
</tr>
<tr>
<td>-Llibrary</td>
<td>Specify a library to be scanned by the linker</td>
</tr>
<tr>
<td>-L-option</td>
<td>Specify <code>-option</code> to be passed directly to the linker</td>
</tr>
<tr>
<td>-Mfile</td>
<td>Request generation of a MAP file</td>
</tr>
<tr>
<td>-MOT</td>
<td>Generate a Motorola S1/S9 HEX format output file</td>
</tr>
<tr>
<td>-MPLAB</td>
<td>Specify compilation and debugging under MPLAB IDE</td>
</tr>
<tr>
<td>-Nsize</td>
<td>Specify identifier length</td>
</tr>
<tr>
<td>-NODEL</td>
<td>Do not remove temporary/intermediate files</td>
</tr>
<tr>
<td>-NOERRATA</td>
<td>Disable errata-fix modifications of the output code</td>
</tr>
<tr>
<td>-NORT</td>
<td>Do not link standard runtime module</td>
</tr>
<tr>
<td>-O</td>
<td>Enable post-pass optimization</td>
</tr>
<tr>
<td>-Ofile</td>
<td>Specify output filename</td>
</tr>
</tbody>
</table>
### 3.4.1 -processor: Define processor

This option defines the processor which is being used. For example to compile for the PIC18C452 processor, use the command line option `-18c452`. You can also add your own processors to the compiler. For more information about this, see the section Processor Support on page 67.

### 3.4.2 -Aaddress: Specify offset for ROM

The `-A` option is used to specify a base address for the ROM image. This option may be required with debugging tools, such as bootloaders, that expect the ROM image to begin at an address other than zero. This option effects all ROM-based psects including reset and interrupt vectors, and affects the linker classes which place all other code and `const` qualified data.

If the base address is positioned in external memory, the `-ROM` option must be used to specify the range of external memory addresses available.

### 3.4.3 -A-option: Specify Extra Assembler Option

The `-A` option can also be used to specify an extra “-” option which will be passed directly to the assembler by PICC18. If `-A` is followed immediately by any text starting with a “-” character, the text

<table>
<thead>
<tr>
<th>Option</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-P</td>
<td>Preprocess assembler files</td>
</tr>
<tr>
<td>-PRE</td>
<td>Produce preprocessed source files</td>
</tr>
<tr>
<td>-PROTO</td>
<td>Generate function prototype information</td>
</tr>
<tr>
<td>-PSECTMAP</td>
<td>Display complete memory segment usage after linking</td>
</tr>
<tr>
<td>-q</td>
<td>Specify quiet mode</td>
</tr>
<tr>
<td>-RESRAMranges</td>
<td>Reserve the specified RAM address ranges.</td>
</tr>
<tr>
<td>-RESROMranges</td>
<td>Reserve the specified ROM address ranges.</td>
</tr>
<tr>
<td>-ROMranges</td>
<td>Specify program space memory for ROM-less devices</td>
</tr>
<tr>
<td>-S</td>
<td>Compile to assembler source files only</td>
</tr>
<tr>
<td>-SIGNED_CHAR</td>
<td>Make the default char signed.</td>
</tr>
<tr>
<td>-STRICT</td>
<td>Enable strict ANSI keyword conformance</td>
</tr>
<tr>
<td>-TEK</td>
<td>Generate a Tektronix HEX format output file</td>
</tr>
<tr>
<td>-Usymbol</td>
<td>Undefine a predefined preprocessor symbol</td>
</tr>
<tr>
<td>-UBROF</td>
<td>Generate an UBROF format output file</td>
</tr>
<tr>
<td>-V</td>
<td>Verbose: display compiler pass command lines</td>
</tr>
<tr>
<td>-Wlevel</td>
<td>Set compiler warning level</td>
</tr>
<tr>
<td>-Wlevel!</td>
<td>Set compiler warning level and stop on warnings</td>
</tr>
<tr>
<td>-X</td>
<td>Eliminate local symbols from symbol table</td>
</tr>
<tr>
<td>-Zg[level]</td>
<td>Enable global optimization in the code generator</td>
</tr>
</tbody>
</table>

Table 3 - 2 PICC18 Options
will be passed directly to the assembler without being interpreted by PICC18. For example, if the option -A-H is specified, the -H option will be passed on to the assembler when it is invoked which will display constant values as hexadecimal values in the assembler output.

3.4.4 -AAHEX: Generate American Automation Symbolic Hex

The -AAHEX option directs PICC18 to generate an American Automation symbolic format HEX file, producing a file with the .hex extension. This option has no effect if used with an option which specifies a .bin file output. The American Automation hex format is an enhanced Motorola S-Record format which includes symbol records at the start of the file. This option should be used if producing code which is to be debugged with an American Automation in-circuit emulator.

3.4.5 -ASMLIST: Generate Assembler .LST Files

The -ASMLIST option tells PICC18 to generate an assembler listing file for each module being compiled. The list file shows both the original C code, and the generated assembler code and the corresponding binary op-codes. The listing file will have the same name as the source file, and a file type (extension) of .lst. Provided the link stage has successfully concluded, the listing file will be updated by the linker so that it contains absolute addresses and symbol values. Thus you may use the assembler listing file to determine the position of, and exact op codes corresponding to, instructions.

3.4.6 -BIN: Generate Binary Output File

The -BIN option tells PICC18 to generate a binary image output file. The output file will be given an extension of .bin. Binary output may also be selected by specifying an output file of type .bin using the -O file option.

3.4.7 -Bl: Select Large Memory Model

The -Bl option tells PICC18 to compile using the large memory model. See section 4.6.3 on page 97 for more details about this memory model.

3.4.8 -Bs: Select Small Memory Model

The -Bs option tells PICC18 to compile using the small memory model. See section 4.6.3 on page 97 for more details about this memory model.

3.4.9 -C: Compile to Object File

The -C option is used to halt compilation after generating a relocatable object file. This option is frequently used when compiling multiple source files using a “make” utility. If multiple source files are specified to the compiler each will be compiled to a separate .obj file. The object files will be placed in the directory in which PICC18 was invoked, to handle situations where source files are located in read-only directories. To compile three source files main.c, module1.c and asmcode.as to object files you could use a command similar to:
PICC18 Command Line Driver

```
PICC18 -18C452 -O -Zg -C main.c module1.c asmcode.as
```

The compiler will produce three object files `main.obj`, `module1.obj` and `asmcode.obj` which could then be linked to produce an Intel HEX file using the command:

```
PICC18 -18C452 main.obj module1.obj asmcode.obj
```

The compiler will accept any combination of `.c`, `.as` and `.obj` files on the command line. Assembler source files will be passed directly to the assembler and object files will not be used until the linker is invoked. Unless the `-Ofile` option is used to specify an output filename and type, the final output will be an Intel HEX file with the same base name as the first source or object file specified on the command line and the extension `.hex`. The example above would produce a file called `main.hex`.

### 3.4.10 -CKfile: Generate Checksum

This option causes OBJTOHEX to use `file` for checksum specifications. Table 7 - 4 on page 226 for further details.

### 3.4.11 -CP16: Use 16-bit Wide Program Space Pointers

The `-CP16` option is the default if no code space pointer size is specified. It will make all pointer objects that can point to program, or code, space 16 bits wide. The pointers affected are function pointers, pointers to `const` and pointers to `far`. This option should be used when compiling for chips that have 64k bytes of program space or less, or if only the first 64k bytes of this space is dereferenced using a pointer. See Pointer Types on page 89 for more information on pointer types.

### 3.4.12 -CP24: Use 24-bit Wide Program Space Pointers

The `-CP24` option will make all pointer objects that can point to program, or code, space 24 bits wide. The pointers affected are function pointers, pointers to `const` and pointers to `far`. This option should be used when there is more than 64k bytes of program space on the target device and pointer dereferences are made to addresses above the 64k byte limit. See Pointer Types on page 89 for more information on pointer types.

### 3.4.13 -CRfile: Generate Cross Reference Listing

The `-CR` option will produce a cross reference listing. If the `file` argument is omitted, the “raw” cross reference information will be left in a temporary file, leaving the user to run the `CREF` utility. If a filename is supplied, for example `-CRtest.crf`, PICC18 will invoke `CREF` to process the cross reference information into the listing file, in this case `test.crf`. If multiple source files are to be included in the cross reference listing, all must be compiled and linked with the one PICC18 command. For example, to generate a cross reference listing which includes the source modules `main.c`, `module1.c` and `nvram.c`, compile and link using the command:

```
PICC18 -18C452 -CRmain.crf main.c module1.c nvram.c
```
3.4.14 -D24: Use 24-bit Doubles

This option is the default if no double size option is specified on the command line. It specifies the use of truncated 24-bit floating point format for objects of type double. It has no affect for object defined as float. See Floating Point Types and Variables on page 84 for more details.

3.4.15 -D32: Use 32-bit Doubles

This option tells the compiler to use the IEEE754 32-bit floating point format for double variables. See Floating Point Types and Variables on page 84 for more details. This option is automatically enabled if the -FDOUBLE option is selected.

3.4.16 -Dmacro: Define Macro

The -D option is used to define a preprocessor macro on the command line, exactly as if it had been defined using a #define directive in the source code. This option may take one of two forms, -Dmacro which is equivalent to:

```
#define macro 1
```

placed at the top of each module compiled using this option, or -Dmacro=text which is equivalent to:

```
#define macro text
```

where text is the textual substitution required. Thus, the command:

```
PICC18 -18C252 -Ddebug -Dbuffers=10 test.c
```

will compile test.c with macros defined exactly as if the C source code had included the directives:

```
#define debug 1
#define buffers 10
```

3.4.17 -E: Define Format for Compiler Errors

If the -E option is not used, the default behaviour of the compiler is to display any errors in a “human readable” format line with a caret “^” and error message pointing out the offending characters in the source line, for example:

```
x.c: main()
  4: PORT_A = xFF;
      ^ undefined identifier: xFF
```

This standard format is perfectly acceptable to a person reading the error output, but is not usable with environments which support compiler error handling. The following sections indicate how this option may be used in such situations.
3.4.17.1 Using the -E Option

Using the -E option instructs the compiler to generate error messages in a format which is acceptable to some text editors and development environments.

If the same source code as used in the example above were compiled using the -E option, the error output would be:

    x.c 4 9: undefined identifier: xFF

indicating that the error occurred in file x.c at line 4, offset 9 characters into the statement. The second numeric value - the column number - is relative to the left-most non-space character on the source line. If an extra space or tab character were inserted at the start of the source line, the compiler would still report an error at line 4, column 9.

3.4.17.2 Modifying the Standard -E Format

If the -E option does not meet your editor’s requirement, you can redefine its format by setting two environment variables: HTC_ERR_FORMAT and HTC_WARN_FORMAT. These environment variables are in the form of a printf-style string in which you can use the specifiers shown in Table 3 - 3.

<table>
<thead>
<tr>
<th>Specifier</th>
<th>Expands To</th>
</tr>
</thead>
<tbody>
<tr>
<td>%f</td>
<td>Filename</td>
</tr>
<tr>
<td>%l</td>
<td>Line number</td>
</tr>
<tr>
<td>%c</td>
<td>Column number</td>
</tr>
<tr>
<td>%s</td>
<td>Error string</td>
</tr>
</tbody>
</table>

The column number is relative to the left-most non-space character on the source line. Here is an example of setting the environment variables from within DOS:

    set HTC_WARN_FORMAT=WARNING: file %f; line %l; column %c; %s
    set HTC_ERR_FORMAT=ERROR: file %f; line %l; column %c; %s

Using the previous source code, the output from the compiler when using the above environment variables would be:

    ERROR: file x.c; line 4; column 9; undefined identifier: xFF

Remember that if these environment variables are set in a batch file, you must prepend the specifiers with an additional percent character to stop the specifiers being interpreted immediately by DOS, e.g. the filename specifier would become %%f.
The Microchip MPLAB IDE requires the following definitions for HTC_WARN_FORMAT and HTC_ERR_FORMAT, respectively, to enable it to display the corresponding line with the error or warning in its editor.

Warning[ ] file %f %l : %s
Error[ ] file %f %l : %s

3.4.17.3 Redirecting Errors to a File

Error output, either in standard or -E format, can be redirected into files using UNIX- or DOS-style standard output redirection. The error from the example above could have been redirected into a file called errlist using the command:

PICC18 -18C242 -E x.c > errlist

Compiler errors can also be appended onto existing files using the redirect and append syntax. If the error file specified does not exist it will be created. To append compiler errors to a file, use a command like:

PICC18 -18C242 -E x.c >> errlist

3.4.18 -Efile: Redirect Compiler Errors to a File

Some editors do not allow the standard command line redirection facilities to be used when invoking the compiler. To work with these editors, PICC18 allows the error listing filename to be specified as part of the -E option. Error files generated using this option will always be in -E format. For example, to compile x.c and redirect all errors to x.err, use the command:

PICC18 -18C242 -Ex.err x.c

The -E option also allows errors to be appended to an existing file by specifying an addition character “+” at the start of the error filename, for example:

PICC18 -18C242 -E+x.err y.c

If you wish to compile several files and combine all of the errors generated into a single text file, use the -E option to create the file then use -E+ when compiling all the other source files. For example, to compile a number of files with all errors combined into a file called project.err, you could use the -E option as follows:

PICC18 -18C242 -Eproject.err -O -Zg -C main.c
PICC18 -18C242 -E+project.err -O -Zg -C part1.c
PICC18 -18C242 -E+project.err -C asmodecode.as

The file project.err will contain any errors from main.c, followed by the errors from part1.c and then asmodecode.as, for example:
3.4.19 -FDOUBLE

This option is used to enable the use of faster 32-bit floating point routines. These alternative routines make floating point multiplication and division many times faster than the default routines, but do require more ROM and RAM space. These routines are only available for 32-bit double types. The -D32 option is automatically selected when using the -FDOUBLE option.

3.4.20 -FAKELOCAL

This option should be used in conjunction with the -G option to produce debug information that is specific to Microchip’s MPLAB environment. It will allow the user to debug using variables local to a function, e.g. in watch windows. The debug information associated with source-level single stepping is also modified so that a better correlation between the source and instructions in the program memory window is obtained. This option also adjusts the format for compiler errors so that they can be more readily interpreted by the MPLAB IDE. See also Section 4.15 on page 117.

3.4.21 -G: Generate source-level Symbol File

The -G option generates a source-level symbol file (i.e. a file which allows tools to determine which line of source code is associated with machine code instructions, and determine which source-level variable names correspond with areas of memory, etc.) for use with HI-TECH Software debuggers and simulators such as Lucifer. If no filename is given, the symbol file will have the same base name as the first source or object file specified on the command line, and an extension of .sym. For example the option -GTEST.SYM generates a symbol file called test.sym. Symbol files generated using the -G option include source-level information for use with source-level debuggers.

Note that all source files for which source-level debugging is required should be compiled with the -G option. The option is also required at the link stage, if this is performed separately. For example:

    PICC18 -18C252 -G -C test.c
    PICC18 -18C252 -C module1.c
    PICC18 -18C252 -Gtest.sym test.obj module1.obj

will include source-level debugging information for test.c only because module1.c was not compiled with the -G option.
PICC18 Compiler Options

3.4.22 -HELP: Display Help

When used with no other options present on the command line, the -HELP option displays information on the PICC18 options.

3.4.23 -ICD

This option can be used to indicate that the output code is to be downloaded to the MPLAB In-Circuit Debugger. It will make appropriate adjustments to the linker options required by the ICD. When used, this option defines a macro called MPLAB_ICD.

3.4.24 -I path: Include Search Path

Use -I to specify an additional directory to use when searching for header files which have been included using the #include directive. The -I option can be used more than once if multiple directories are to be searched. The default include directory containing all standard header files are always searched even if no -I option is present. The default search path is searched after any user-specified directories have been searched. For example:

```
PICC18 -18C252 -C -Ic:\include -Id:\myapp\include test.c
```

will search the directories c:\include and d:\myapp\include for any header files included into the source code, then search the default include directory which is typically c:\ht-pic18\include.

3.4.25 -INTEL: Generate INTEL Hex File

The -INTEL option directs PICC18 to generate an Intel HEX file producing a file with .hex extension. This option has no effect if used with an option which specifies a .bin file output.

3.4.26 -L library: Scan Library

The -L option is used to specify additional libraries which are to be scanned by the linker. Libraries specified using the -L option are scanned before the standard C library, allowing additional versions of standard library functions to be accessed. For example when compiling for the PIC18C452 processor, the floating-point version of printf() can be linked by searching the library pic800-f.lib using the option -Lf.

The argument to -L is a library keyword to which the prefix PIC; digits representing the processor range, the number of ROM banks and the number of RAM banks; and the suffix .lib are added. Thus the option -Ll will, for example, scan the library pic800-l.lib and the option -Lxx will scan a library called pic800-xx.lib. All libraries must be located in the LIB subdirectory of the compiler installation directory. As indicated, the argument to the -L option is not a complete library filename.

If you wish the linker to scan libraries whose names do not follow the above naming convention or whose locations are not in the LIB subdirectory, simply include the libraries’ names on the command
line along with your source files. Alternatively, the linker may be invoked directly allowing the user to manually specify all the libraries to be scanned.

**3.4.26.1 Printf with Additional Support for Longs and Floats**

By default, `printf()` and related functions contained in the standard libraries support the printing of integers only. If additional support for `longs` and `floats` is required different libraries must be used during the link stage. For complete information on `printf()`, see page 270.

To use the version of `printf()` which has additional support for `longs`, you must include a supplementary library by using the following option when linking (or when you are compiling and linking in one step):

```
-Ll
```

To use the version of `printf()` which has additional support for `longs` and `floats`, use the option:

```
-Lf
```

In the above options, `l` and `f` are merely specifying the library type as described in the section Standard Libraries on page 76. In addition to these options, no modification of your source code is required to print `longs` or `floats`. If these options are not specified, `printf()` will not know about the `long` and `float` placeholders and any attempt to use them will print the placeholder character(s) as text with no substitution.

An alternative version of `printf()` with additional flags and functionality can be selected by using the command line option:

```
-Lw
```

This version of `printf()` can print `long` and `float` variables by default, and is larger than the standard version of `printf()`.

**3.4.27 -L-option: Specify Extra Linker Option**

The `-L` option can also be used to specify an extra “-” option which will be passed directly to the linker by PICC18. If `-L` is followed immediately by any text starting with a dash character “-”, the text will be passed directly to the linker without being interpreted by PICC18. For example, if the option `-L-FOO` is specified, the `-FOO` option will be passed on to the linker when it is invoked.

The `-L` option is especially useful when linking code which contains extra program sections (or `psects`), as may be the case if the program contains C code which makes use of the `#pragma psect` directive or assembler code which contains user-defined psects. See Section The `#pragma psect` Directive for more information. If the `-L` option did not exist, it would be necessary to invoke the linker manually to link code which uses the extra psects.
One commonly used linker option is \(-N\), which sorts the symbol table in the map file by address, rather than by name. This would be passed to PICC18 as the option \(-L-N\).

The \(-L\) option can also be used to replace default linker options. For example, \(-L-ARAM=0-35Fh\) will inform the linker that the available memory address for objects in the RAM psect class range from 0 to 35Fh. By default, this range is 0 to 5FFh for the PIC18C452 processor. The default option that you are replacing must contain an equal character.

3.4.28 -Mfile: Generate Map File

The \(-M\) option is used to request the generation of a map file. If no filename is specified, the map information is displayed on the screen, otherwise the filename specified with \(-M\) will be used.

3.4.29 -MPLAB: Compile and Debug using MPLAB IDE

The \(-MPLAB\) option informs the HI-TECH C that both compilation and subsequent debugging will be performed from within the Microchip MPLAB IDE. This option turns on source level debugging (\(-G\)), turns on the \(-FAKELOCAL\) option to allow enhanced source and variable tracking, and adjusts the compiler's error message format (\(-E\)) to be that required by the MPLAB IDE.

If compilation is performed under a separate make facility, but debugging is performed under the MPLAB IDE, then the \(-G\), \(-E\) and \(-FAKELOCAL\) options can be used separately.

3.4.30 -MOT: Generate Motorola S-Record HEX File

The \(-MOT\) option directs PICC18 to generate a Motorola S-Record HEX file if producing a file with .hex extension. This option has no effect if used with an option which specifies a .bin file output.

3.4.31 -Nsize: Identifier Length

This option allows the C identifier length to be increased from the default value of 31. Valid sizes for this option are from 32 to 255. The option has no effect for all other values.

3.4.32 -NODEL Do not remove temporary/intermediate files

When this option is enabled, the intermediate files created during the build process will not be deleted. Having access to these files allows for greater customization over the project build process.

3.4.33 -NOERRATA Disable errata-fix modifications

By default there may be modifications made to the output code of the compiler to work around errata documented by Microchip. Some chips may not have any errata and not all errata is corrected by the compiler. The use of this options disables any errata fixes made by the compiler. There may be situations when the errata changes can be safely disabled. See the release notes for more information.
3.4.34 -NORT: Do Not Link Standard Runtime Module

Using this option will not link in the standard runtime startup module. The user should then supply their own version of the runtime startup module in the list of input files on the command line. Even if the required startup module does not contain executable code, it will almost certainly require symbol and psect definitions for successful compilation, so this module cannot simply be omitted completely. The source for the standard runtime module is supplied in the SOURCES directory of your distribution and this should be used as the basis for your own runtime module.

See Section 4.3.5 on page 76 for information on customizing the runtime code.

3.4.35 -O: Invoke Optimizer

The -O option invokes the post-pass optimizer after the code generation pass. This is an optimization of the generated assembler code.

3.4.36 -Ofile: Specify Output File

This option allows the name and type of the output file to be specified to the compiler. If no -O option is specified, the output file’s name will be derived from the first source or object file specified on the command line. You can use the -O option to specify an output file of type HEX, BIN or UBR, containing Intel or Motorola HEX records, binary code or UBROF respectively. For example:

```
PICC18 -18C452 -Otest.bin prog1.c part2.c
```

will produce a binary file named test.bin.

3.4.37 -P: Preprocess Assembly Files

The -P option causes the assembler files to be preprocessed before they are assembled thus allowing the use of preprocessor directives, such as `#include`, with assembler code. By default, assembler files are not preprocessed.

3.4.38 -PRE: Produce Preprocessed Source Code

The -PRE option is used to generate preprocessed C source files with an extension .pre. It may be useful to ensure that preprocessor macros have expanded to what you think they should. Use of this option can also create C source files which do not require any separate header files. This is useful when sending files for technical support.

3.4.39 -PROTO: Generate Prototypes

The -PROTO option is used to generate .pro files containing both ANSI and K&R style function declarations for all functions within the specified source files. Each .pro file produced will have the
same base name as the corresponding source file. Prototype files contain both ANSI C-style prototypes and old-style C function declarations within conditional compilation blocks.

The extern declarations from each .pro file should be edited into a global header file which is included in all the source files comprising a project. The .pro files may also contain static declarations for functions which are local to a source file. These static declarations should be edited into the start of the source file. To demonstrate the operation of the -PROTO option, enter the following source code as file test.c:

```c
#include <stdio.h>
add(arg1, arg2)
int * arg1;
int * arg2;
{
    return *arg1 + *arg2;
}

void printlist(int * list, int count)
{
    while (count--)
        printf("%d ", *list++);
    putchar(\n');
}
```

If compiled with the command:

```
PICC18 -18C252 -PROTO test.c
```

PICC18 will produce test.pro containing the following declarations which may then be edited as necessary:

```c
/* Prototypes from test.c */
/* extern functions - include these in a header file */
#if PROTOTYPES
extern int add(int *, int *);
extern void printlist(int *, int);
#else /* PROTOTYPES */
extern int add();
extern void printlist();
#endif /* PROTOTYPES */
```
## 3.4.40 -PSECTMAP: Display Complete Memory Usage

The -PSECTMAP option is used to display a complete memory and psect (program section) dump after linking the user code. The information provided by this option is more detailed than the standard memory usage map which is normally printed after linking. The -PSECTMAP option causes the compiler to print a listing of every compiler and user generated psect, followed by the standard memory usage map. For example:

Psect Usage Map:

<table>
<thead>
<tr>
<th>Psect</th>
<th>Contents</th>
<th>Memory Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>powerup</td>
<td>Power on reset code</td>
<td>$000000 - $000003</td>
</tr>
<tr>
<td>text</td>
<td>Program and library code</td>
<td>$000008 - $000017</td>
</tr>
<tr>
<td>init</td>
<td>Initialization code</td>
<td>$000018 - $00003B</td>
</tr>
<tr>
<td>end_init</td>
<td>Initialization code</td>
<td>$00003C - $00003F</td>
</tr>
<tr>
<td>text</td>
<td>Program and library code</td>
<td>$000040 - $000123</td>
</tr>
<tr>
<td>text</td>
<td>Program and library code</td>
<td>$000124 - $000137</td>
</tr>
<tr>
<td>clrtext</td>
<td>Memory clearing code</td>
<td>$000138 - $00014B</td>
</tr>
<tr>
<td>bss</td>
<td>RAM variables</td>
<td>$0000F6 - $0000F7</td>
</tr>
<tr>
<td>param</td>
<td>Parameter auto blocks</td>
<td>$0000F8 - $0000FB</td>
</tr>
<tr>
<td>data</td>
<td>RAM initialised data</td>
<td>$0000FC - $0000FF</td>
</tr>
<tr>
<td>temp</td>
<td>Temporary RAM data</td>
<td>$000000 - $000001</td>
</tr>
<tr>
<td>idata</td>
<td>ROM image of initialised data</td>
<td>$000004 - $000007</td>
</tr>
</tbody>
</table>

Memory Usage Map:

Program ROM | $000000 - $000003 | $000004 (4) bytes
Program ROM | $000008 - $00014B | $000144 (324) bytes
| $000148 (328) bytes total Program ROM

RAM data    | $0000F6 - $0000FF | $00000A (10) bytes total RAM data
Near RAM    | $000000 - $000001 | $000002 (2) bytes total Near RAM
ROM data    | $000004 - $000007 | $000004 (4) bytes total ROM data

Program statistics:

Total ROM used 332 bytes (1.0%)
Total RAM used 12 bytes (0.8%) Near RAM used 2 bytes (1.6%)
3.4.41 -q: Quiet Mode

If used, this option must be the first option specified on the command line. It places the compiler in a quiet mode which suppresses the HI-TECH Software copyright notice from being displayed.

3.4.42 -RESRAM\texttt{ranges},\texttt{ranges}/

The -RESRAM option is used to reserve a particular section of RAM space. The address ranges must be specified in HEX. The syntax for this option is a comma-separated list of address ranges. For example:

-RESRAM\texttt{20-40}

This will reserve the RAM address range from 0x20 to 0x40. Reserved memory will not be considered by the linker for the placement of objects, thus they can be used by the programmer for any purpose.

3.4.43 -RESROM\texttt{ranges},\texttt{ranges}/

The -RESROM option is used to reserve a particular section of program space. The address ranges must be specified in HEX. The syntax for this option is a comma-separated list of address ranges. For example:

-RESROM\texttt{1000-10FF,2000-20FF}

This will reserve the ROM address ranges 0x1000 to 0x10FF and 0x2000 to 0x20FF. Reserved memory will not be considered by the linker for the placement of objects, thus they can be used by the programmer for any purpose.

3.4.44 -ROM\texttt{ranges}: Specify External Memory

Some PIC18 devices allow external memory to be placed in the chip’s program space. To indicate the additional program space memory present in such systems, the -ROM option can be used. The syntax for this option is a comma-separated list of address ranges, e.g.

-ROM\texttt{0-2FFF,6000-7FFF}

This option specifies the \texttt{CODE} psect class definition passed to the linker, in which is placed all code and ROM data, such as objects qualified as \texttt{const}. Memory specified by this option is appended to any memory already defined for the target device in the \texttt{pic18.ini} file.

If no -ROM option is passed to the PICC18 driver when compiling for a ROM-less device, an error is generated. All ROM-less devices specify a ROM size of zero in the \texttt{pic18.ini} file.

This option can be used for any PIC18 device, even those which do not have external memory interfaces. It is up to the programmer to ensure that the memory selected by this option matches their hardware setup.
3.4.45 -S: Compile to Assembler Code

The -S option stops compilation after generating an assembler source file. An assembler file will be generated for each C source file passed on the command line. The command:

```
PICC18 -18C252 -O -Zg -S test.c
```

will produce an assembler file called `test.as` which contains the code generated from `test.c`. The optimization options -O and -Zg can be used with -S, making it possible to examine the compiler output for any given set of options. This option is particularly useful for checking function calling conventions and signature values when attempting to write external assembly language routines. The file produced by this option differs to that produced by the -ASMLIST option in that it does not contain op-codes or addresses and it may be used as a source file and subsequently passed to the assembler to be assembled.

3.4.46 -SIGNED_CHAR: Make Char Type Signed

Unless this option is used, the default behaviour of the compiler is to make all character values and variables of type unsigned char unless explicitly declared or cast to signed char. This option will make the default char type signed char. When using this option, any unsigned character object will have to be explicitly declared unsigned char.

The range of a signed character type is -128 to +127 and the range of similar unsigned objects is 0 to 255

3.4.47 -STRICT: Strict ANSI Conformance

The -STRICT option is used to enable strict ANSI conformance of all special keywords. HI-TECH C supports various special keywords (for example the persistent type qualifier). If the -STRICT option is used, these keywords are changed to include two underscore characters at the beginning of the keyword (e.g. __persistent) so as to strictly conform to the ANSI standard. Be warned that use of this option may cause problems with some standard header files (e.g. <intrpt.h>).

3.4.48 -TEK: Generate Tektronix HEX File

The -TEK option tells the compiler to generate a Tektronix format HEX file if producing a file with .hex extension. This option has no effect if used with an option which specifies a .bin file output.

3.4.49 -Umacro: Undefine a Macro

The -U option, the inverse of the -D option, is used to undefine predefined macros. This option takes the form -Umacro. The option, -Udraft, for example, is equivalent to:

```
#undef draft
```

placed at the top of each module compiled using this option.
3.4.50 -UBROF: Generate UBROF Format Output File

The -UBROF option tells the compiler to generate a UBROF format output file suitable for use with certain in-circuit emulators. The output file will be given an extension .ubr. UBROF output may also be selected by specifying an output file of type .ubr using the -O option. This option has no effect if used with an option which specifies a .bin file output.

3.4.51 -V: Verbose Compile

The -V is the verbose option. The compiler will display the full command lines used to invoke each of the compiler applications or compiler passes. This option may be useful for determining the exact linker options if you need to directly invoke the HLINK command.

3.4.52 -W level[!] : Set Warning Level

The -W option is used to set the compiler warning level. Allowable warning levels range from -9 to 9. The warning level determines how “picky” the compiler is about dubious type conversions and constructs. The default warning level -W0 will allow all normal warning messages. Warning level -W1 will suppress the message Func() declared implicit int. -W3 is recommended for compiling code originally written with other, less strict, compilers. -W9 will suppress all warning messages. Negative warning levels -W-1, -W-2 and -W-3 enable special warning messages including compile-time checking of arguments to printf() against the format string specified.

Use this option with care as some warning messages indicate code that is likely to fail during execution.

If it is desired to halt execution on any warning messages, concluding the -Wlevel option with an exclamation mark (!) will cause the compiler to stop if any warnings are produced.

3.4.53 -X: Strip Local Symbols

The option -X strips local symbols from any files compiled, assembled or linked. Only global symbols will remain in any object files or symbol files produced.

3.4.54 -Zg[level]: Global Optimization

The -Zg option invokes global optimization during the code generation pass. This can result in significant reductions to code size and internal RAM usage. This optimizer is less critical than the post-pass optimizer, but can still significantly reduce the code size.

Global optimization attempts to optimize register usage on a function-by-function basis. It also takes advantage of constant propagation in code to avoid un-necessary accesses to memory.

The default level for this option is 1 (the least optimization). The level can be set anywhere from 1 to 9 (the most optimization). The number indicates how hard the optimizer tries to reduce code size. For PICC18, there is usually little advantage in using levels above 3.
Features and Runtime Environment

PICC-18 supports a number of special features and extensions to the C language which are designed to ease the task of producing ROM-based applications. This chapter documents the compiler options and special language features which are specific to the Microchip PIC 18 family of processors.

4.1 ANSI Standard Issues

4.1.1 Divergence from the ANSI C Standard

PICC-18 diverges from the ANSI C standard in one area: function recursion. Due to the PIC18’s hardware limitations of no easily-usable stack and limited memory, function recursion is unsupported.

4.1.2 Implementation-defined behaviour

Certain sections of the ANSI standard have implementation-defined behaviour. This means that the exact behaviour of some C code can vary from compiler to compiler. Throughout this manual are sections describing how the PICC-18 compiler behaves in such situations.

4.2 Processor-related Features

PICC-18 has many features which relate directly to the PIC18 family of processors. These are detailed in the following sections.

4.2.1 Processor Support

PICC-18 supports the full range of Microchip PIC 18 processors. Additional code-compatible processors may be added by editing the pic18.ini file in the LIB directory. User-defined processors should be placed at the end of the file. The header of the file explains how to specify a processor. Newly added processors will be available the next time you compile by selecting the name of the new processor on the command line in the usual way.

4.2.2 Configuration Fuses

The PIC18 processor’s have several locations which contain the configuration bits or fuses. These bits may be set using the configuration macro. The macro has the form:

```
__CONFIG(n, x)
```
(there are two leading underscore characters) where \( n \) is the configuration register number and \( x \) is the value that is to be the configuration word. The macro is defined in `<pic18.h>` so be sure to include that into each module that uses this macro.

The configuration macro programs the upper and lower half of each register, i.e. it programs 16 bits with each call. Special named quantities are defined in the header file appropriate for the processor you are using to help you enable the required features. Table 4 - 1 on page 69 for the available configuration bit settings associated with the 18Cxxx chip types, and Table 4 - 2 on page 70 for the settings for flash devices.

For example, to set a PIC18Cxx1 chip to have an RC type oscillator, an 8-bit bus width, the powerup timer disabled, the watchdog timer enabled with a post scale factor of 1:1, and the stack full/underflow reset disabled, the following could be used.

```c
#include <pic18.h>
__CONFIG(1, RC);
__CONFIG(2, BW8 & PWRTDIS & WDTPS1 & WDTEN);
__CONFIG(4, STVRDIS);
```

Note that the individual selections are ANDed together. Any bits which are not selected in these macros will remain unprogrammed. You should ensure that you have specified all bits correctly to ensure proper operation of the part when programmed. Consult your PIC datasheet for more details.

The `__CONFIG` macro does not produce executable code and should be placed outside function definitions.

### 4.2.3 ID Locations

The PIC18 devices have locations outside the addressable memory area that can be used for storing program information, such as an ID number. The `__IDLOC` macro may be used to place data into these locations. The macro is used in a manner similar to:

```c
#include <pic18.h>
__IDLOC(x);
```

where \( x \) is a list of nibbles which are to be positioned in to the ID locations. The upper nibble of each location is programmed as 0FH so that the whole byte is treated as a `nop` instruction if executed. The following:

```c
__IDLOC(15F01);
```

will attempt to fill five ID locations with the hexadecimal values: F1H, F5H, FFH, F0H and F1H. The base address of the ID locations is specified by the `idloc` psect which will be automatically assigned an address dependent on the type of processor selected.
Table 4 - 1 Configuration bit settings for 18Cxxx parts

<table>
<thead>
<tr>
<th>Meaning</th>
<th>Config register</th>
<th>18Cxx1</th>
<th>18Cxx2</th>
<th>18Cxx8</th>
</tr>
</thead>
<tbody>
<tr>
<td>Code protection</td>
<td>1 low byte</td>
<td>n/a</td>
<td>PROTECT, UNPROTECT</td>
<td>PROTECT, UNPROTECT</td>
</tr>
<tr>
<td>Oscillator system clock switch</td>
<td>1 high byte</td>
<td>n/a</td>
<td>OSCSEN, OSCSDIS</td>
<td>OSCSEN, OSCSDIS</td>
</tr>
<tr>
<td>Oscillator types</td>
<td>1 high byte</td>
<td>RC, HS, EC, LP</td>
<td>RCIO, HSPLL, ECIO, EC, RC, HS, XT, LP</td>
<td>RCIO, HSPLL, ECIO, EC, RC, HS, XT, LP</td>
</tr>
<tr>
<td>Bus width</td>
<td>2 low byte</td>
<td>BW16, BW8</td>
<td>n/a</td>
<td>n/a</td>
</tr>
<tr>
<td>Powerup timer enable</td>
<td>2 low byte</td>
<td>PWRTEN, PWRTDIS</td>
<td>PWRTEN, PWRTDIS</td>
<td>PWRTEN, PWRTDIS</td>
</tr>
<tr>
<td>Brown out reset voltage</td>
<td>2 low byte</td>
<td>n/a</td>
<td>BORV25, BORV27, BORV42, BORV45</td>
<td>BORV25, BORV27, BORV42, BORV45</td>
</tr>
<tr>
<td>Brown out reset enable</td>
<td>2 low byte</td>
<td>n/a</td>
<td>BOREN, BORDIS</td>
<td>BOREN, BORDIS</td>
</tr>
<tr>
<td>Watchdog timer post scale select</td>
<td>2 high byte</td>
<td>WDTPS1 - WDTPS128</td>
<td>WDTPS1 - WDTPS128</td>
<td>WDTPS1 - WDTPS128</td>
</tr>
<tr>
<td>Watchdog timer enable</td>
<td>2 high byte</td>
<td>WDTEN, WDTDIS</td>
<td>WDTEN, WDTDIS</td>
<td>WDTEN, WDTDIS</td>
</tr>
<tr>
<td>CCP2 Mux bit</td>
<td>3 high byte</td>
<td>n/a</td>
<td>CCP2RC1, CCP2RB3</td>
<td>n/a</td>
</tr>
<tr>
<td>Stack full/underflow reset enable</td>
<td>4 low byte</td>
<td>STVREN, STVRDIS</td>
<td>STVREN, STVRDIS</td>
<td>STVREN, STVRDIS</td>
</tr>
</tbody>
</table>

### 4.2.4 EEPROM Data

For those PIC 18 devices that support external programming of their EEPROM data area, the `__EEPROM_DATA()` macro can be used to place the initial EEPROM data values into the HEX file ready for programming. The macro is used as follows.
### Table 4 - 2 Configuration bit settings for 18Fxxx parts

<table>
<thead>
<tr>
<th>Meaning</th>
<th>Config register</th>
<th>18Fxx2</th>
<th>18Fxx8</th>
<th>18Fxx20</th>
</tr>
</thead>
<tbody>
<tr>
<td>Oscillator system clock switch</td>
<td>1 high byte</td>
<td>OSCSEN, OSCSDIS</td>
<td>OSCSEN, OSCSDIS</td>
<td>OSCSEN, OSCSDIS</td>
</tr>
<tr>
<td>Oscillator types</td>
<td>1 high byte</td>
<td>RCRA6, HSPLL, ECRA6, ECDB4, RC, HS, XT, LP</td>
<td>RCRA6, HSPLL, ECRA6, ECDB4, RC, HS, XT, LP</td>
<td>RCRA6, HSPLL, ECRA6, ECDB4, RC, HS, XT, LP</td>
</tr>
<tr>
<td>Powerup timer enable</td>
<td>2 low byte</td>
<td>PWRTEN, PWRTDIS</td>
<td>PWRTEN, PWRTDIS</td>
<td>PWRTEN, PWRTDIS</td>
</tr>
<tr>
<td>Brown out reset voltage</td>
<td>2 low byte</td>
<td>BORV20, BORV27, BORV42, BORV45</td>
<td>BORV20, BORV27, BORV42, BORV45</td>
<td>BORV25, BORV27, BORV42, BORV45</td>
</tr>
<tr>
<td>Brown out reset enable</td>
<td>2 low byte</td>
<td>BOREN, BORDIS</td>
<td>BOREN, BORDIS</td>
<td>BORDIS</td>
</tr>
<tr>
<td>Watchdog timer post scale select</td>
<td>2 high byte</td>
<td>WDTPS1 - WDTPS128</td>
<td>WDTPS1 - WDTPS128</td>
<td>WDTPS1 - WDTPS128</td>
</tr>
<tr>
<td>Watchdog timer enable</td>
<td>2 high byte</td>
<td>WDTEN, WDTDIS</td>
<td>WDTEN, WDTDIS</td>
<td>WDTEN, WDTDIS</td>
</tr>
<tr>
<td>External bus data wait enable</td>
<td>3 low byte</td>
<td>n/a</td>
<td>n/a</td>
<td>WAITEN*, WAITDIS*</td>
</tr>
<tr>
<td>*18F8x20 only</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Microcontroller/Microprocessor mode</td>
<td>3 high byte</td>
<td>n/a</td>
<td>n/a</td>
<td>MCU*, MPU*, MPUBB*, XMCU*</td>
</tr>
<tr>
<td>*18F8x20 only</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CCP2 Mux bit</td>
<td>3 high byte</td>
<td>CCP2RC1, CCP2RB3</td>
<td>n/a</td>
<td>CCP2RC1, CCP2RE7, CCP2RB3 *</td>
</tr>
</tbody>
</table>
Table 4 - 2 Configuration bit settings for 18Fxxx parts

<table>
<thead>
<tr>
<th>Meaning</th>
<th>Config register</th>
<th>18Fxx2</th>
<th>18Fxx8</th>
<th>18Fxx20</th>
</tr>
</thead>
<tbody>
<tr>
<td>Stack full/underflow reset enable</td>
<td>4 low byte</td>
<td>STVREN, STVRDIS</td>
<td>STVREN, STVRDIS</td>
<td>STVREN, STVRDIS</td>
</tr>
<tr>
<td>Debug enable</td>
<td>4 low byte</td>
<td>DEBUGEN, DEBUGDIS</td>
<td>DEBUGEN, DEBUGDIS</td>
<td>DEBUGEN, DEBUGDIS</td>
</tr>
<tr>
<td>Low voltage ICSP</td>
<td>4 low byte</td>
<td>LVPEN, LVPDIS</td>
<td>LVPEN, LVPDIS</td>
<td>LVPEN, LVPDIS</td>
</tr>
<tr>
<td>Code protection *18Fx720 only **18Fx58 only</td>
<td>5 low byte</td>
<td>CPA, CP3, CP2, CP1, CP0</td>
<td>CPA, CP3**, CP2**, CP1, CP0</td>
<td>CPA, CP7*, CP6*, CP5*, CP4*, CP3, CP2, CP1, CP0</td>
</tr>
<tr>
<td>Data protection</td>
<td>5 high byte</td>
<td>CPD</td>
<td>CPD</td>
<td>CPD</td>
</tr>
<tr>
<td>Boot code protection</td>
<td>5 high byte</td>
<td>CPB</td>
<td>CPB</td>
<td>CPB</td>
</tr>
<tr>
<td>Protect all blocks</td>
<td>5 high byte</td>
<td>CPALL</td>
<td>CPALL</td>
<td>CPALL</td>
</tr>
<tr>
<td>Write protect/enable *18Fx720 only **18Fx58 only</td>
<td>6 low byte</td>
<td>WP3, WP2, WP1, WP0</td>
<td>WPA, WP3**, WP2**, WP1, WP0</td>
<td>WPA, WP7*, WP6*, WP5*, WP4*, WP3, WP2, WP1, WP0</td>
</tr>
<tr>
<td>Write protect/enable boot block</td>
<td>6 high byte</td>
<td>WPB, WRTEN</td>
<td>WPB, WPU</td>
<td>WPB, WPU</td>
</tr>
<tr>
<td>Write protect/enable configuration registers</td>
<td>6 high byte</td>
<td>WPC, WRTEN</td>
<td>WPC, WPU</td>
<td>WPC, WPU</td>
</tr>
<tr>
<td>Write protect/enable data block</td>
<td>6 high byte</td>
<td>WPD, WRTEN</td>
<td>WPD, WPU</td>
<td>WPD, WPU</td>
</tr>
</tbody>
</table>
### Table 4 - 2 Configuration bit settings for 18Fxxx parts

<table>
<thead>
<tr>
<th>Meaning</th>
<th>Config register</th>
<th>18Fxx2</th>
<th>18Fxx8</th>
<th>18Fxx20</th>
</tr>
</thead>
<tbody>
<tr>
<td>Write protect all blocks</td>
<td>6</td>
<td>WPALL</td>
<td>WPALL</td>
<td>WPALL</td>
</tr>
<tr>
<td>Table read protection/enable</td>
<td>7 low byte</td>
<td>TRP3, TRP2, TRP1, TRP0</td>
<td>TRPA, TRP3**, TRP2**, TRP1, TRP0</td>
<td>TRPA, TRP7*, TRP6*, TRP5*, TRP4*, TRP3, TRP2, TRP1, TRP0</td>
</tr>
<tr>
<td>* 18Fx720 only</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>** 18Fx58 only</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Table read protection boot block</td>
<td>7 high byte</td>
<td>TRPB</td>
<td>TRPB</td>
<td>TRPB</td>
</tr>
<tr>
<td>Protect all blocks from table reads</td>
<td>7</td>
<td>TRPALL</td>
<td>TRPALL</td>
<td>TRPALL</td>
</tr>
</tbody>
</table>
#include <pic18.h>
__EEPROM_DATA(0, 1, 2, 3, 4, 5, 6, 7);

The macro accepts eight parameters, being eight data values. Each value should be a byte in size. Unused values should be specified as a parameter of zero. The macro may be called multiple times to define the required amount of EEPROM data. It is recommended that the macro be placed outside any function definitions.

The macro defines, and places the data within, a psect called `eeprom_data`. This psect is positioned by a linker option in the usual way.

This macro is not used to write to EEPROM locations during run-time. The macros `EEPROM_READ()` and `EEPROM_WRITE()`, and the function versions of these macros, can be called to read from, and write to, the EEPROM during program execution.

### 4.2.5 EEPROM and Flash Runtime Access

EEPROM and flash memory macros are defined for convenience and are available for chips that have EEPROM or flash memory. The predefined EEPROM and flash memory macros can be used in the following manner.

To write a byte-size value to an address in EEPROM memory:

```c
EEPROM_WRITE(address, value);
```

To read a byte of data from an address in EEPROM memory, and store it in a variable:

```c
variable=EEPROM_READ(address);
```

For convenience, `EEPROM_SIZE` predefines the total size of data EEPROM available on chip.

To copy a block of code/data to an area in flash memory:

```c
flash_write(source_pointer, length, dest_pointer);
```

To read a byte of data from an address in flash memory, and store in a variable:

```c
variable=FLASH_READ(address);
```

### 4.2.6 Bit Instructions

Wherever possible, PICC-18 will attempt to use the PIC18 bit instructions. For example, when using a bitwise operator and a mask to alter a bit within an integral type, the compiler will check the mask value to determine if a bit instruction can achieve the same functionality.

```c
int foo;
foo |= 0x40;
```
will produce the instruction

bsf _foo, 6

To set or clear individual bits within integral types, the following macros could be defined and used.

```c
#define bitset(var, bitno) ((var) |= 1 << (bitno))
#define bitclr(var, bitno) ((var) &= ~(1 << (bitno)))
```

To perform the same operation as the above example, the bitset macro could be employed as follows.

```c
bitset(foo, 6);
```

### 4.2.7 Multi-byte SFRs

Some of the SFRs associated with the PIC18 can be grouped to form multi-byte values, e.g. the TMRxH and TMRxL register together form a 16-bit timer count value. Depending on the device and mode of operation, there may be hardware requirements to read these registers correctly, e.g. the TMRxL register often must be read before trying to read the TMRxH register to obtain a valid 16-bit result.

Although it is possible to define an absolute non-`char` C variable to map over such registers, the order in which PICC-18 reads the bytes of a multi-byte object varies depending on the context of the variable in an expression, i.e. it may read the most significant byte first, or the least. Thus, it highly recommended that the existing SFR `char` definitions in the chip header files be used. Each SFR should be accessed directly and in the required order by the programmer’s code. This will ensure a much higher degree of portability.

The following code copies the two byte registers into C `unsigned` variable `i` for subsequent use.

```c
i = TMR0L;
i += TMR0H << 8;
```

### 4.3 Files

#### 4.3.1 Source Files

The extension used with source files is important as it is used by the compiler drivers to determine their content. Source files containing C code should have the extension `.c`, assembler files should have extensions of `.as`, relocatable object files require the `.obj` extension, and library files should be named with a `.lib` extension. See the tutorial Section 2.1.2 on page 18 for more information on how these input files are processed by the compiler.

#### 4.3.2 Output File Formats

The compiler is able to directly produce a number of the output file formats which are used by common PROM programmers and in-circuit emulators.
The default behaviour of the **PICC18** command is to produce *Bytecraft COD* and *Intel HEX* output. If no output filename or type is specified, **PICC18** will produce a *Bytecraft COD* and *Intel HEX* file with the same base name as the first source or object file specified on the command line. Table 4 - 3 on page 75 shows the output format options available with **PICC18**. The *File Type* column lists the filename extension which will be used for the output file.

### Table 4 - 3 Output File Formats

<table>
<thead>
<tr>
<th>Format Name</th>
<th>Description</th>
<th>PICC18 Option</th>
<th>File Type</th>
</tr>
</thead>
<tbody>
<tr>
<td><em>Motorola HEX</em></td>
<td>S1/S9 type hex file</td>
<td>-MOT</td>
<td>.hex</td>
</tr>
<tr>
<td><em>Intel HEX</em></td>
<td><em>Intel</em> style hex records (default)</td>
<td>-INTEL</td>
<td>.hex</td>
</tr>
<tr>
<td>Binary</td>
<td>Simple binary image</td>
<td>-BIN</td>
<td>.bin</td>
</tr>
<tr>
<td>UBRof</td>
<td>“Universal Binary Image Relocatable Format”</td>
<td>-UBROF</td>
<td>.ubr</td>
</tr>
<tr>
<td><em>Tektronix HEX</em></td>
<td><em>Tektronix</em> style hex records</td>
<td>-TEK</td>
<td>.hex</td>
</tr>
<tr>
<td><em>American Automation HEX</em></td>
<td>Hex format with symbols for <em>American Automation emulators</em></td>
<td>-AAHEX</td>
<td>.hex</td>
</tr>
<tr>
<td><em>Bytecraft .COD</em></td>
<td><em>Bytecraft</em> code format (default)</td>
<td>n/a (default)</td>
<td>.cod</td>
</tr>
<tr>
<td>Library</td>
<td>HI-TECH library file</td>
<td>n/a</td>
<td>.lib</td>
</tr>
</tbody>
</table>

In addition to the options shown, the `-O` option may be used to request generation of binary or UBRof files. If you use the `-O` option to specify an output filename with a `.bin` type, for example `-Otest.bin`, **PICC18** will produce a binary file. Likewise, if you need to produce UBRof files, you can use the `-O` option to specify an output file with type `.ubr`, for example `-Otest.ubr`.

### 4.3.3 Symbol Files

The **PICC18** `-G` option tells the compiler to produce a symbol file which can be used by debuggers and simulators to perform symbolic and source-level debugging. This option produces symbol files which contain both assembler- and C-level information. If no symbol filename is specified, by default a file called `file.sym` will be produced, where `file` is the basename of the first source file specified on the command line. For example, to produce a symbol file called `test.sym` which includes C source-level information:

```
PICC18 -18C252 -Gtest.sym test.c init.c
```

This option will also generate a different symbol file for each module compiled. These files do not contain absolute address and have the type `.sdb`. The base name will be the same as the base name of the module being compiled. Thus the above command line would also generate symbols files with the names `test.sdb` and `init.sdb`. 
4.3.4 Standard Libraries

PICC-18 includes a number of standard libraries, each with the range of functions described in the Library Functions chapter.

Figure 4 - 1 on page 76 illustrates the naming convention used for the standard libraries. The meaning of each field is described here, where:

- **Processor Type** is always `pic`.
- **Processor Range** is 8 for the PIC18 family.
- **Configuration** is a digit, bit 0 of which is either 1 for 24-bit wide program space pointers; otherwise 0. Bit 1 is 0 to disallow the use of the LFSR instruction; 1 to allow this instruction.
- **Memory Model** is either `l` for large or `s` for small model.
- **Double Type** is `.` for 24-bit doubles, and `d` for 32-bit doubles.
- **Library Type** is `c` for standard library, `l` for the library which contains only printf-related functions with additional support for longs, and `f` for the library which contains only printf-related functions with additional support for longs and floats.

![Figure 4 - 1 PICC-18 Standard Library Naming Convention](image)

### 4.3.4.1 Limitations of printf

The `printf()` function is provided but some features have been removed. For more details on this function, see the documentation on page 270.

### 4.3.5 Run-time startup Modules

A C program requires certain objects to be initialised and the processor to be in a particular state before it can begin execution. It is the job of the run-time startup code to ready the program for execution. Since this is code that executed before the C program, it is necessarily written in assembler code. The run-time startup code is executed almost immediately after reset. In fact it is called by a special `powerup` routine,
described below, that is directly located at the reset vector address. For the PIC18 processors, the principle job of the run-time startup code is to clear uninitialized variables and assign values to those variables that have been initialised.

The run-time startup code will clear, or assign the value zero, any variables which are uninitialized at their definition and which are non-auto. This amounts to those objects which have been placed in the bss, rbss or rbit psects. Since these psects are defined as a contiguous block of memory, the run-time startup code calls a routine to clear a block of memory for each psect. In the following example, all but the object loc will be cleared by the startup code since it is an auto object. The initial value of loc is unknown.

```c
int i;
near int ni;
bit b;
void main(void)
{
    static int sloc;
    int loc;

...```

The code which clears these psects is only included if it is necessary. The modules which contain the clear routines can be found in the SOURCES directory of your distribution. The file clr.as contains the code to clear a block of memory. This is called, if required, by code in the files clrbss.as clrbbss.as and clrbit.as which initiate the clearing of each psect.

Each of these runtime modules define a symbol, or label, that forces the module to be linked when required. By including code which defines these symbols into your project the standard library routines to clear or copy data on startup can be replaced or removed.

For example to replace the code that clears bit variables, include into your project a module that defines the symbol clear_bit. This label should be the entry point to your version of the code that would normally clear the bit variables. To completely remove the code that clears bit variables, include into your project a module that merely defines this same symbol. There does not need to be any code associated with the symbol in this instance.

If you require variables to be untouched by the runtime code, you would normally qualify them as persistent. This method does not require any modification of the runtime code.

The other function of the run-time startup code is to initialise those variables assigned a value at their definition. This amounts to a block copy of the initial values from ROM to the RAM areas designated for those objects. Code to perform the copy is only included if required. In this example:
The objects `i`, `ni` and `sloc` will be initialised by the run-time startup code. Note that you cannot initialise `bit` objects and that initialised automatic variables are assigned their starting value by code placed within the function in which they are defined.

Any objects defined in assembler code, i.e. they have memory reserved using any of the `DS`, `DB` or `DW` assembler directives, will also be cleared or initialised at startup providing that the directives are placed within the compiler-generated psect used for C variables, such as `bss`, `rbss`, `data` or `rbit` etc.

The run-time startup code jumps to the function `main()`, which is referred to as `_main` by the run-time startup code. Note the underscore “_” prepended to the function name. A jump rather than a call is executed to save one level of stack. The function `main()` is, by definition of the C language, the “main program”.

The run-time startup code is provided by the pre-compiled, standard module, `picrt800.obj`, found in the LIB directory.

The source code used to generate the run-time startup module is called `picrt18x.as` which is in the SOURCES directory of your distribution. This module can be excluded from a program by using the `-NORT` command line option. A user-defined version of this module can then be written and used to replace the default module.

In addition to this module will be the routines, mentioned above, to copy data or clear memory as required. These routines are not called by name, but are linked in, if required, to a position in `picrt18x.as` indicated by a comment.

**4.3.5.1 The powerup Routine**

Some hardware configurations require special initialisation, often within the first few cycles of execution after reset. Rather than having to modify the run-time startup module to achieve this there is a hook to the reset vector provided via the `powerup` routine. This is a user-supplied assembler module that will be executed immediately on reset. Often this can be embedded in a C module as embedded assembler code. A “dummy” powerup routine is included in the file `powerup.as`. The file can be copied, modified and included into your project to replace the default powerup routine that is present in the standard libraries.
The powerup routine should be written assuming that little or no RAM is working and should only use system resources after it has tested and enabled them. The following example code shows the default powerup routine which are in the standard library:

```c
#include "sfr.h"

global powerup,start
psect powerup,class=CODE,delta=1

powerup:
    goto start
end powerup
```

The assembler header file `sfr.h` has been included so reference to the PIC’s registers can be made. You will need to specify the `-p` command-line driver option.

The powerup routine is intended to be relatively small, and since it is linked to an address lower than the interrupt vectors, it may interfere with them if it grows too large in size. The lower interrupt vector address is 08h on the PIC18 devices. To avoid overwriting the interrupt vectors, the powerup routine can be made to jump to a separate function, which will be linked at a different location, and this separate function can call the jump to `start`. The following gives an example of this:

```c
global powerup,start,big_powerup
psect powerup,class=CODE,delta=1

powerup:
    goto big_powerup

psect big_powerup,class=CODE,delta=1
big_powerup:
    ; powerup code...
    goto start
```

In this example, the `big_powerup` psect will be positioned somewhere in the memory allocated to the `CODE` class unless an explicit linker option is added by the user.

### 4.4 Supported Data Types and Variables

The PICC-18 compiler supports basic data types of 1, 2 and 4 byte size. All multi-byte types follow _least significant byte first_ format, also known as _little-endian_. Word size values thus have the least significant byte at the lower address, and double word size values have the least significant byte and least significant word at the lowest address.
Table 4 - 4 shows the data types and their corresponding size and arithmetic type.

Table 4 - 4 Data Types

<table>
<thead>
<tr>
<th>Type</th>
<th>Size (in bits)</th>
<th>Arithmetic Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>bit</td>
<td>1</td>
<td>boolean</td>
</tr>
<tr>
<td>char</td>
<td>8</td>
<td>signed or unsigned integer&lt;sup&gt;a&lt;/sup&gt;</td>
</tr>
<tr>
<td>unsigned char</td>
<td>8</td>
<td>unsigned integer</td>
</tr>
<tr>
<td>short</td>
<td>16</td>
<td>signed integer</td>
</tr>
<tr>
<td>unsigned short</td>
<td>16</td>
<td>unsigned integer</td>
</tr>
<tr>
<td>int</td>
<td>16</td>
<td>signed integer</td>
</tr>
<tr>
<td>unsigned int</td>
<td>16</td>
<td>unsigned integer</td>
</tr>
<tr>
<td>long</td>
<td>32</td>
<td>signed integer</td>
</tr>
<tr>
<td>unsigned long</td>
<td>32</td>
<td>unsigned integer</td>
</tr>
<tr>
<td>float</td>
<td>24</td>
<td>real</td>
</tr>
<tr>
<td>double</td>
<td>24 or 32&lt;sup&gt;b&lt;/sup&gt;</td>
<td>real</td>
</tr>
</tbody>
</table>

<sup>a</sup>A char is unsigned by default, and signed if the PICC18 -SIGNED_CHAR option is used.

<sup>b</sup>A double defaults to 24-bit, but becomes 32-bit with the PICC18 -D32 option.

4.4.1 Radix Specifiers and Constants

The format of integral constants specifies their radix. PICC-18 supports the ANSI standard radix specifiers as well as one which enables binary constants to specified in C code. The format used to specify the radices are given in Table 4 - 5 on page 80. The letters used to specify binary or hexadecimal radices are case insensitive, as are the letters used to specify the hexadecimal digits.

Table 4 - 5 Radix Formats

<table>
<thead>
<tr>
<th>Radix</th>
<th>Format</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>binary</td>
<td>0bnumber or 0Bnumber</td>
<td>0b10011010</td>
</tr>
<tr>
<td>octal</td>
<td>0number</td>
<td>0763</td>
</tr>
<tr>
<td>decimal</td>
<td>number</td>
<td>129</td>
</tr>
<tr>
<td>hexadecimal</td>
<td>0xnumber or 0Xnumber</td>
<td>0x2F</td>
</tr>
</tbody>
</table>

Any integral constant will have a type which is the smallest type that can hold the value without overflow. The suffix 1 or L may be used with the constant to indicate that it must be assigned either a
**Supported Data Types and Variables**

**signed long** or **unsigned long** type, and the suffix *u* or *U* may be used with the constant to indicate that it must be assigned an unsigned type, and both *l* or *L* and *u* or *U* may be used to indicate **unsigned long int** type.

Floating-point constants have **double** type unless suffixed by *f* or *F*, in which case it is a **float** constant. The suffixes *l* or *L* specify a **long double** type which is considered an identical type to **double** by PICC-18.

Character constants are enclosed by single quote characters “’”, for example ‘a’. A character constant has char type. Multi-byte character constants are not supported.

String constants or string literals are enclosed by double quote characters “ “, for example “hello world”. The type of string constants is **const char** * and the strings are stored in ROM. Assigning a string constant to a non-**const** char pointer will generate a warning from the compiler. For example:

```c
char * cp = "one";  // "one" in ROM, produces warning
const char * ccp = "two";  // "two" in ROM
char ca[] = "two";  // "two" different to the above
```

A non-**const** array initialised with a string, for example the last statement in the above example, produces an array in RAM which is initialised at startup time with the string “two” (copied from ROM), whereas a constant string used in other contexts represents an unnamed **const**-qualified array, accessed directly in ROM.

PICC-18 will use the same storage location and label for strings that have identical character sequences, except where the strings are used to initialise an array residing in RAM as indicated in the last statement in the above example.

Two adjacent string constants (i.e. two strings separated only by white space) are concatenated by the compiler. Thus:

```c
const char * cp = "hello world";
```

assigned the pointer with the string “hello world”.

**4.4.2 Bit Data Types and Variables**

PICC-18 supports **bit** integral types which can hold the values 0 or 1. Single **bit** variables may be declared using the keyword **bit**. **Bit** objects declared within a function, for example:

```c
static bit init_flag;
```

will be allocated in the bit-addressable psect **rbit**, and will be visible only in that module or function. When the following declaration is used outside any function:

```c
bit init_flag;
```
**Features and Runtime Environment**

init_flag will be globally visible, but located within the same psect.

Bit variables cannot be auto or parameters to a function. A function may return a bit object by using
the bit keyword in the function’s prototype in the usual way.

Bit variables behave in most respects like normal unsigned char variables, but they may only
contain the values 0 and 1, and therefore provide a convenient and efficient method of storing boolean
flags without consuming large amounts of internal RAM. It is not possible to declared pointers to bit
variables or statically initialise bit variables. Operations on bit objects are performed using the single
bit instructions wherever possible, thus the generated code to access bit objects is very efficient.

Note that when assigning a larger integral type to a bit variable, only the least-significant bit is used.
For example, if the bit variable bitvar was assigned as in the following:

```c
int data = 0x54;
bit bitvar;

bitvar = data;
```

it will be cleared by the code since the least significant bit of data is zero.

If you want to set a bit variable to be 0 or 1 depending on whether the larger integral type is zero (false)
or non-zero (true), use the form:

```c
bitvar = data != 0;
```

The psects in which bit objects are allocated storage are declared using the bit psect directive flag.
Eight bit objects will take up one byte of storage space which is indicated by the bit psects’ scale value
of 8 in the map file. The length given in the map file for bit psects is in units of bits, not bytes. All
addresses specified for bit objects are also bit addresses.

The bit psects are cleared on startup, but are not initialised. To create a bit object which has a non-zero
initial value, explicitly initialise it at the beginning of your code.

If the PICC18 flag -STRICT is used, the bit keyword becomes unavailable.

**4.4.2.1 Using Bit-Addressable Registers**

The bit variable facility may be combined with absolute variable declarations (see page 93) to access
bits at specific addresses. Absolute bit objects are numbered from 0 (the least significant bit of the first
byte) up. Therefore, bit number 3 (the fourth bit in the byte since numbering starts with 0) in byte number
5 is actually absolute bit number 43 (that is 8bits/byte * 5 bytes + 3 bits).

For example, to access the power down detection flag bit in the RCON register, declare RCON to be a C
object at absolute address 03h, then declare a bit variable at absolute bit address 27:
static unsigned char RCON @ 0xFD0;

static near bit PD @ (unsigned)&RCON*8+2;

Note that all standard registers and bits within these registers are defined in the header files provided. The only header file you need to include to have access to the PIC registers is `<pic18.h>` - at compile time this will include the appropriate header for the selected chip.

### 4.4.3 8-Bit Integer Data Types and Variables

PICC-18 supports both `signed char` and `unsigned char` 8-bit integral types. If the `signed` or `unsigned` keyword is absent, the default type is `unsigned char` unless the `PIC18 -SIGNED_CHAR` option is used, in which case it is `signed char`. The `signed char` type is an 8-bit two’s complement signed integer type, representing integral values from -128 to +127 inclusive. The `unsigned char` is an 8-bit unsigned integer type, representing integral values from 0 to 255 inclusive. It is a common misconception that the C `char` types are intended purely for ASCII character manipulation. This is not true, indeed the C language makes no guarantee that the default character representation is even ASCII. The `char` types are simply the smallest of up to four possible integer sizes, and behave in all respects like integers.

The reason for the name “char” is historical and does not mean that `char` can only be used to represent characters. It is possible to freely mix `char` values with `short`, `int` and `long` values in C expressions. With PICC-18 the `char` types will commonly be used for a number of purposes, as 8-bit integers, as storage for ASCII characters, and for access to I/O locations. The default `unsigned char` type is the most efficient data type on the PIC and maps directly onto the 8-bit bytes which are most efficiently manipulated by PIC instructions. It is suggested that `char` types be used wherever possible so as to maximize performance and minimize code size.

Variables may be declared using the `signed char` and `unsigned char` keywords, respectively, to hold values of these types. Where only `char` is used in the declaration, the type will be `unsigned char` unless the option, mentioned above, to specify `signed char` as default is used.

### 4.4.4 16-Bit Integer Data Types

PICC-18 supports four 16-bit integer types. `int` and `short` are 16-bit two’s complement signed integer types, representing integral values from -32,768 to +32,767 inclusive. `Unsigned int` and `unsigned short` are 16-bit unsigned integer types, representing integral values from 0 to 65,535 inclusive. All 16-bit integer values are represented in little endian format with the least significant byte at the lower address. Both `int` and `short` types are 16 bits wide as this is the smallest integer size allowed by the ANSI standard for C. The sizes of the integer types were chosen so as not to violate the ANSI standard. Allowing a smaller integer size, such as 8 bits would lead to a serious incompatibility with the C standard. 8-bit integers are already fully supported by the `char` types and should be used in place of `int` types wherever possible.
Variables may be declared using the `signed int`, `unsigned int`, `signed short int` and `unsigned short int` keyword sequences, respectively, to hold values of these types. Where only `int` is used in the declaration, the type will be `signed int`. When specifying a `short int` type, the keyword `int` may be omitted. Thus a variable declared as `short` will contain a `signed short int` and a variable declared as `unsigned short` will contain an `unsigned short int`.

### 4.4.5 32-Bit Integer Data Types and Variables

PICC-18 supports two 32-bit integer types. `Long` is a 32-bit two's complement signed integer type, representing integral values from \(-2,147,483,648\) to \(+2,147,483,647\) inclusive. `Unsigned long` is a 32-bit unsigned integer type, representing integral values from 0 to 4,294,967,295 inclusive. All 32-bit integer values are represented in *little endian* format with the least significant word and least significant byte at the lowest address. `Long` and `unsigned long` occupy 32 bits as this is the smallest long integer size allowed by the ANSI standard for C.

Variables may be declared using the `signed long int` and `unsigned long int` keyword sequences, respectively, to hold values of these types. Where only `long int` is used in the declaration, the type will be `signed long`. When specifying this type, the keyword `int` may be omitted. Thus a variable declared as `long` will contain a `signed long int` and a variable declared as `unsigned long` will contain an `unsigned long int`.

### 4.4.6 Floating Point Types and Variables

Floating point is implemented using the IEEE 754 32-bit format and a modified IEEE 754 (truncated) 24-bit format.

The truncated 24-bit format is used for all `float` values. For `double` values, the truncated 24-bit format is the default, but may be explicitly invoked with the `PICC18 -D24` option. The 32-bit format is used for doubles by using the `PICC18 -D32` option. The `long double` type is identical to the `double` type.

Both of these formats are described in Table 4 - 6, where:

- `sign` is the sign bit
- `exponent` is an 8-bit exponent which is stored as *excess 127* (i.e. an exponent of 0 is stored as 127)
- `mantissa` is the mantissa, which is to the right of the radix point. There is an implied bit to the left of the radix point which is always 1 except for a zero value, where the implied bit is zero. A zero value is indicated by a zero exponent.

The value of this number is \((-1)^{\text{sign}} \times 2^{(\text{exponent}-127)} \times 1.\text{mantissa}\).

Here are some examples of the IEEE 754 32-bit and modified IEEE 754 24-bit formats:
Supported Data Types and Variables

Table 4 - 6 Floating Point Formats

<table>
<thead>
<tr>
<th>Format</th>
<th>Sign</th>
<th>biased exponent</th>
<th>mantissa</th>
</tr>
</thead>
<tbody>
<tr>
<td>IEEE 754 32-bit</td>
<td>x</td>
<td>xxxx xxxx</td>
<td>xxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx</td>
</tr>
<tr>
<td>Modified IEEE 754 24-bit</td>
<td>x</td>
<td>xxxx xxxx</td>
<td>xxx xxxx xxxx xxxx</td>
</tr>
</tbody>
</table>

Table 4 - 7 IEEE 754 32-bit and 24-bit Examples

<table>
<thead>
<tr>
<th>Format</th>
<th>Number</th>
<th>biased exponent</th>
<th>1.mantissa</th>
<th>decimal</th>
</tr>
</thead>
<tbody>
<tr>
<td>IEEE 754 32-bit</td>
<td>7DA6B69Bh</td>
<td>11111011b (251)</td>
<td>1.01001101011011010011011b (1.302447676659)</td>
<td>2.77000e+37</td>
</tr>
<tr>
<td>Modified IEEE 754 24-bit</td>
<td>42123Ah</td>
<td>10000100b (132)</td>
<td>1.001001000111010b (1.142395019531)</td>
<td>36.557</td>
</tr>
</tbody>
</table>

Note that the most significant bit of the mantissa column in Table 4 - 7 on page 85 (that is the bit to the left of the radix point) is the implied bit, which is assumed to be 1 unless the exponent is zero (in which case the float is zero).

The 32-bit example in Table 4 - 7 on page 85 can be calculated manually as follows.

The sign bit is zero; the biased exponent is 251, so the exponent is 251-127=124. Take the binary number to the right of the decimal point in the mantissa. Convert this to decimal and divide it by $2^{23}$ where 23 is the number of bits taken up by the mantissa, to give 0.302447676659. Add one to this fraction. The floating-point number is then given by:

$$(-1)^0 \times 2^{124} \times 1.302447676659 = 1 \times 2.126764793256e+37 \times 1.302447676659 \approx 2.77000e+37$$

Variables may be declared using the float and double keywords, respectively, to hold values of these types. Floating point types are always signed and the unsigned keyword is illegal when specifying a floating point type. Types declared as long double will use the same format as types declared as double.

4.4.7 Structures and Unions

PICC-18 supports struct and union types of any size from one byte upwards. Structures and unions only differ in the memory offset applied for each member. The members of structures and unions may not be objects of type bit, but bit fields are fully supported.

Structures and unions may be passed freely as function arguments and return values. Pointers to structures and unions are fully supported.
4.4.7.1 Bit Fields in Structures

PICC-18 fully supports *bit fields* in structures.

Bit fields are always allocated within 8-bit words. The first bit defined will be the least significant bit of the word in which it will be stored. When a bit field is declared, it is allocated within the current 8-bit unit if it will fit, otherwise a new 8-bit byte is allocated within the structure. Bit fields can never cross the boundary between 8-bit allocation units. For example, the declaration:

```c
struct {
    unsigned lo : 1;
    unsigned dummy : 6;
    unsigned hi : 1;
} foo;
```

will produce a structure occupying 1 byte. If `foo` was ultimately linked at address 10H, the field `lo` will be bit 0 of address 10H, `hi` will be bit 7 of address 10H. The least significant bit of `dummy` will be bit 1 of address 10H and the most significant bit of `dummy` will be bit 6 of address 10H.

Unnamed bit fields may be declared to pad out unused space between active bits in control registers. For example, if `dummy` is never used the structure above could have been declared as:

```c
struct {
    unsigned lo : 1;
    unsigned          : 6;
    unsigned hi : 1;
} foo;
```

If a bit field is declared in a structure that is assigned an absolute address, no storage will be allocated for the structure. Absolute structures would be used when mapping a structure over a register to allow a portable method of accessing individual bits within the register.

A structure with bitfields may be initialised by supplying a comma-separated list of initial values for each field. For example:

```c
struct {
    unsigned lo  : 1;
    unsigned mid : 6;
    unsigned hi  : 1;
} foo = {1, 8, 0};
```

4.4.7.2 Structure and Union Qualifiers

PICC-18 supports the use of type qualifiers on structures. When a qualifier is applied to a structure, all of its members will inherit this qualification. In the following example the structure is qualified `const`.

```c
struct {
    unsigned lo : 1;
    unsigned mid : 6;
    unsigned hi : 1;
} foo;
```
const struct {
    int number;
    int *ptr;
} record = { 0x55, &i};

In this case, the structure will be placed into ROM and each member will, obviously, be read-only. Remember that all members must be initialised if a structure is **const**.

If the members of the structure were individually qualified **const** but the structure was not, then the structure would be positioned into RAM, but each member would be read-only. Compare the following structure with the above.

```c
struct {
    const int number;
    int * const ptr;
} record = { 0x55, &i};
```

### 4.4.8 Standard Type Qualifiers

*Type qualifiers* provide information regarding how an object may be used, in addition to its type which defines it storage size and format. PICC-18 supports both ANSI qualifiers and additional special qualifiers which are useful for embedded applications and which take advantage of PIC18 architecture.

#### 4.4.8.1 Const and Volatile Type Qualifiers

PICC-18 supports the use of the ANSI type qualifiers **const** and **volatile**.

The **const** type qualifier is used to tell the compiler that an object has a constant value and will not be modified. If any attempt is made to modify an object declared **const**, the compiler will issue a warning. User defined objects declared **const** are placed in a special sections in ROM. Obviously, a **const** object must be initialised when it is declared as it cannot be assigned a value at any point in the code following. For example:

```c
const int  version = 3;
```

The **volatile** type qualifier is used to tell the compiler that an object cannot be guaranteed to retain its value between successive accesses. This prevents the optimizer from eliminating apparently redundant references to objects declared **volatile** because it may alter the behaviour of the program to do so. All Input/Output ports and any variables which may be modified by interrupt routines should be declared **volatile**, for example:

```c
volatile static near unsigned char  PORTA @ 0xF80;
```

**Volatile** objects may be accessed in a different way to non-**volatile** objects. For example, when assigning a non-**volatile** object the value 1, the object will be cleared and then incremented, but the
same operation performed on a **volatile** object will load the W register with 1 and then store this to the appropriate address.

### 4.4.9 Special Type Qualifiers

PICC-18 supports special type qualifiers, **persistent**, **near** and **far** to allow the user to control placement of **static** and **extern** class variables into particular address spaces. If the PICC18 option, `-STRICT` is used, these type qualifiers are changed to **__persistent**, **__near** and **__far**, respectively. These type qualifiers may also be applied to pointers. These type qualifiers may not be used on variables of class **auto**; if used on variables local to a function they must be combined with the **static** keyword. For example, you may not write:

```c
void test(void)
{
    persistent int intvar;  /* WRONG! */

    .. other code ..
}
```

because `intvar` is of class **auto**. To declare `intvar` as a **persistent** variable local to function `test()`, write:

```c
static persistent int intvar;
```

PICC-18 also supports the keywords **bank1**, **bank2** and **bank3**. These keywords have been included to allow code to be easily ported from PICC. These keywords are accepted by PICC-18, but have no effect in terms of the object’s storage or how they are accessed. These keywords do, however, affect the storage of objects when compiling with the PIC C compiler - see your PIC C manual for more details.

#### 4.4.9.1 Persistent Type Qualifier

By default, any C variables that are not explicitly initialised are cleared to zero on startup. This is consistent with the definition of the C language. However, there are occasions where it is desired for some data to be preserved across resets or even power cycles (on-off-on).

The **persistent** type qualifier is used to qualify variables that should not be cleared on startup. In addition, any **persistent** variables will be stored in a different area of memory to other variables. **Persistent** objects are placed within one of the non-volatile psects. If the **persistent** object is also qualified **near**, it placed in the **nvram** psect. **Persistent bit** objects are placed within the **nvbit** psect. All other **persistent** objects are placed in the **nvram** psect.

There are some library routines provided to check and initialise **persistent** data - see page 268 for more information, and for an example of using **persistent** data.
4.4.9.2 Near Type Qualifier

The near type qualifier is used to place static variables in the access bank of the PIC18. Near objects are represented by 8 bit addresses and the access bank is always accessible regardless of the currently selected RAM bank so accessing near objects may be faster than accessing other objects, and typically results in smaller code sizes.

Here is an example of an unsigned char object placed within the access bank:

```
static near unsigned char fred;
```

4.4.9.3 Far Type Qualifier

The far type qualifier is used to place static variables into external program space of the PIC18 for those devices which can support additional memory. Accesses to these variables are less efficient than accesses to internal variables and extensive accesses to far variables will result in larger code sizes.

Here is an example of an unsigned char object placed into the device’s external code space:

```
farm unsigned int farvar;
```

Note that if the external memory area uses addresses greater than FFFFh (as in most cases) the -CP24 command line option will be required in order to access these variables correctly.

4.4.10 Bdata Type qualifier

The bdata type qualifier only has significance when compiling in the small memory model. In this model all static and extern class variables are placed in the access bank, but this qualifier specifies that the object is to be placed outside the access bank in the banked data area of the device. The object then behaves like an unqualified object in large model. This qualifier is useful when the access bank has overflowed by a small amount as it allows some objects to moved back into the banked memory and prevents having to revert to the large memory model.

4.4.11 Pointer Types

There are two basic pointer types supported PICC-18: data pointers and function pointers. Data pointers hold the address of data objects which can be read and/or written by the program. Function pointers hold the address of an executable routine which can be called indirectly via the pointer.

Of the data pointers, RAM pointers are limited to accessing only the data space (RAM) of the PIC18 device, but the const and far pointer types can access the data and program space (typically ROM, although hardware using devices with an external memory interface may implement any type of memory in this space).
4.4.11.1 RAM Pointers

All RAM pointer objects on these PIC18 devices are 16 bits wide, with the exception of pointers to objects qualified as `near` which are 8 bits wide.

A pointer to RAM, for example:

```c
char * cp;
```

is 16 bits wide and can access all of the RAM available on the PIC18 devices.

A pointer to `near` is only 8 bits wide and can access the general-purpose RAM area of the access bank. In other words they can be used to dereference any variable qualified as `near`. The amount of general purpose RAM in the access bank varies from device to device. Being smaller in size, using pointers to `near` result in smaller code sizes. If a pointer only ever accesses `near`-qualified objects, then that pointer should be qualified as a pointer to `near`.

The operation of RAM pointers are unaffected by the `-CP24/-CP16` switch, nor are they affected by the choice of memory model.

4.4.11.2 Const and Far Pointers

`Const` and `far` pointers can either be 16 or 24 bits wide. Their size can be toggled with the `-CP24` or `-CP16` command line option. The code used to dereference them also changes with their size. The same pointer size must be used for all modules in a project.

A pointer to `far` is identical to a pointer to `const`, except that pointers to `far` may be used to write to the address they hold. A pointer to `const` objects cannot be used to write as the `const` qualifier imposes that the object is read-only.

`Const` and `far` pointers which are 16 bits wide can access all RAM areas and most of the program space. At runtime when dereferenced, the contents of the pointer are examined. For addresses above the upper limit of RAM the program space is accessed using table read or table write instructions. Addresses below the upper limit of RAM access the data space. Even if the address held by a pointer to `const` is in RAM, the RAM location may not be changed.

The default linker options always place `const` data at addresses above the upper limit of the data space so that the correct memory space is accessed when dereferencing with pointers.

If the target device selected has more than 64k bytes of program space memory, then only the lower 64k bytes may be accessed with 16-bit wide pointers. Provided that all program space objects that need to be dereferenced are in the lower 64k bytes, 16-bit pointers to `const` and `far` objects may still be used. The smaller pointer size results in less RAM required and less code produced and so should be used whenever possible.
**Supported Data Types and Variables**

**Const** and **far** pointers which are 24 bits wide can access all RAM areas and all of the program space. At runtime when dereferenced, the contents of the pointer are examined. If bit number 21 in the address is set, the address is assumed to be a RAM address. Bit number 21 of the address is then ignored. If Bit number 21 is clear, then the address is assumed to be of an object in the program space and the access is performed using table read or table write instructions. Again, no writes to objects are permitted using a pointer to **const**.

Note that when dereferencing a 24-bit pointer, the most significant implemented bit (bit number 21) of the TBLPTRU register may be overwritten. This bit may be used to enable access to the configuration area of the PIC18 device. If loading the table pointer registers from hand-written assembler code, make no assumptions about the state of bit number 21 prior to executing table read or write instructions.

### 4.4.11.3 Function Pointers

Function pointers can be defined to indirectly call functions or routines in the program space. The size of these pointers are 16 or 24 bits wide and is controlled by the `-CP24/-CP16` command line option. When 16-bit wide function pointers are used only routines within the lower 64k bytes can be indirectly called. The larger 24-bit function pointer allows indirect calls to be made to any routine, but at the expense of increased code size and RAM usage.

It should be stressed that the `-CP16` option affects function pointer sizes and that it does not affect code that calls functions directly, i.e. by their name rather than indirectly via a pointer. Thus you can still directly call functions residing at any location even if you are using the `-CP16` option, however you can only indirectly call functions that reside in the lower 64k byte area of the program space.

The addresses for all code labels are always shown in the map file as an untruncated byte address regardless of the options used.

### 4.4.11.4 Combining Type Qualifiers and Pointers

Pointers can be qualified like any other C object, but care must be taken when doing so as there are two quantities associated with pointers. The first is the actual *pointer* itself, which is treated like any ordinary C variable and has memory reserved for it. The second is the *object* that the pointer references, or to which the pointer points. The general form of a pointer definition looks like the following.

“object’s type & qualifiers”∗“pointer’s qualifiers”∗“pointer’s name”;

Here are three examples, highlighting the fields with spacing:

```c
near int * nip ;
int * nearinp ;
near int * nearninp;
```
The first example is a pointer called `nip`. It contains the address of an `int` object that is qualified `near`. Since a `near` object is in the access bank the pointer is only 8 bits wide as discussed above. The pointer itself (i.e. the 8-bit value the pointer holds) will reside somewhere in the main banked memory.

The second example is a pointer called `inp` which contains the address of an `int` object. Since this object is not qualified `near`, the pointer needs 16 bits to access the object’s location. The `near` keyword after the `*` indicates that the pointer itself has been qualified `near` and so the pointer (i.e. the 16-bit value the pointer holds) will reside in the access bank, but the object whose address the pointer holds is located in the main banked memory.

The last example is of a pointer called `ninp` which is itself qualified `near` and which also holds the address of an object that is also qualified `near`. In this example, both the pointer and the object that the pointer references will be located in the access bank. The pointer will be 8 bits wide.

The rule is as follows: if the modifier is to the left of the `*` in the pointer declaration, it applies to the object which the pointer addresses. If the modifier is to the right of the `*`, it applies to the pointer variable itself.

The `const`, `volatile`, `far` and `persistent` modifiers may also be applied to pointers in the above manner.

To allow portability between PICC and PICC-18 code, the use of the `bank1`, `bank2` and `bank3` keywords is allowed with PICC-18 pointer definitions in the manner described above. These keywords have no effect and do not alter the way in which indirect accesses are made. The use of these keywords with pointers defined in PICC does affect the operation of pointer dereferences - see your PICC manual for more details.

### 4.5 Storage Class and Object Placement

Objects are positioned in different memory areas dependant on their storage class and declaration. This is discussed in the following sections.

#### 4.5.1 Local Variables

A `local variable` is one which only has scope within the block in which it was defined. That is, it may only be referenced within that block. C supports two classes of local variables in functions: `auto` variables which are normally allocated in the function’s auto-variable block, and `static` variables which are always given a fixed memory location and have permanent duration.

#### 4.5.1.1 Auto Variables

`Auto` (short for `automatic`) variables are the default type of local variable. Unless explicitly declared to be `static` a local variable will be made `auto`, however the `auto` keyword may be used if desired. `Auto` variables are allocated in the `auto-variable block` and referenced by indexing off the symbol that represents that block. The variables will not necessarily be allocated in the order declared - in contrast
to parameters which are always in lexical order. Note that most type qualifiers cannot be used with auto variables, since there is no control over the storage location. The exceptions are const and volatile.

All auto variables are allocated memory within one bank of RAM. At present, all functions share the same bank of memory for auto objects. The size of a functions auto-variable block may not exceed the size of one bank, which is 100H bytes.

The auto-variable blocks for a number of functions are overlapped by the linker if those functions are never called at the same time.

Auto objects are referenced with a symbol that consists of a question mark, “?”, concatenated with a_function plus some offset, where function is the name of the function in which the object is defined. For example, if the int object test is the first object placed in main()’s auto-variable block it will be accessed using the addresses ?a_main and ?a_main+1 since an int is two bytes long.

Auto variables may be accessed using the banked instructions of the PIC18. When accessing auto objects with banked instructions, the compiler will ensure that the bank of the auto-variable block is selected using a movlb instruction, and then access the locations using the appropriate instructions. In essence this amounts to an 8 bit access within the selected bank.

4.5.1.2 Static Variables

Uninitialized static variables are allocated in one of the bss, rbss or bigbss psects. Objects qualified near appear in the rbss psect; objects larger than one bank in size or byte long objects are placed in the bigbss psect and the remainder in the bss psect. They will occupy fixed memory locations which will not be overlapped by storage for other functions. Static variables are local in scope to the function in which they are declared, but may be accessed by other functions via pointers since they have permanent duration. Static variables are guaranteed to retain their value between calls to a function, unless explicitly modified via a pointer. Static variables are not subject to any architectural limitations on the PIC18.

Static variables which are initialised are only done so once during the program’s execution. Thus, they may be preferable over initialised auto objects which are assigned a value every time the block in which the definition is placed is executed.

4.5.2 Absolute Variables

A global or static variable can be located at an absolute address by following its declaration with the construct @ address, for example:

```
volatile unsigned char Portvar @ 0x06;
```

will declare a variable called Portvar located at 06h. Note that the compiler does not reserve any storage, but merely equates the variable to that address, the compiler-generated assembler will include a line of the form:
Note also that the compiler and linker do not make any checks for overlap of absolute variables with other variables of any kind, so it is entirely the programmer’s responsibility to ensure that absolute variables are allocated only in memory not in use for other purposes.

This construct is primarily intended for equating the address of a C identifier with a microprocessor register. To place a user-defined variable at an absolute address, define it in a separate psect and instruct the linker to place this psect at the required address. See “The #pragma psect Directive” on page 112.

Absolute variables have their address supplied by the code generator, not the linker, and hence no symbols are used which require fixup by the linker. This means that the name of the object will not be present in the map file, or any symbol information produced by the linker.

4.5.3 Objects in Program Space

Const objects are usually placed in program space. On the PIC18 devices, the program space is byte-wide, the compiler stores one character per byte location and values are read using the table read instructions. All const-qualified data objects and string literals are placed in the const psect. The const psect is placed at an address above the upper limit of RAM since RAM and const pointers use this address to determine if an access to ROM or RAM is required. See section 4.4.11 on page 89.

4.6 Functions

4.6.1 Function Argument Passing

The method used to pass function arguments depends on the size of the argument or arguments.

If there is only one argument, and it is one byte in size, it is passed in the W register.

If there is only one argument, and it is greater than one byte in size, it is passed in the argument area of the called function. If there are subsequent arguments, these arguments are also passed in the argument area of the called function. The argument area is referenced by an offset from the symbol ?_function, where function is the name of the function concerned.

If there is more than one argument, and the first argument is one byte in size, it is passed in the W register, with subsequent arguments being passed in the argument area of the called function.

In the case of a variable argument list, which is defined by the ellipsis symbol . . . , the calling function builds up the variable argument list and passes a pointer to the variable part of the argument list in btemp. Btemp is the label at the start of the temp psect (the psect used for temporary data).

Take, for example, the following ANSI-style function:
void test(char a, int b)
{
}

The function test() will receive the parameter b in its function argument block and a in the W register. A call:

test( 'a', 8);

would generate code similar to:

movlw   08h
movff   wreg,?_test
movlw   0h
movff   wreg,?_test+1
movlw   061h
call    (_test)

In this example, the parameter b is held in the memory locations ?_test and ?_test+1.

If you need to determine, for assembler code for example, the exact entry or exit code within a function or the code used to call a function, it is often helpful to write a dummy C function with the same argument types as your assembler function, and compile to assembler code with the PICC -S option, allowing you to examine the assembler code.

4.6.2 Function Return Values

Function return values are passed to the calling function as follows:

4.6.2.1 8-Bit Return Values

Eight-bit values are returned from a function in the W register. For example, the function:

char return_8(void)
{
    return 0;
}

will exit with the following code:

movlw   0
return
4.6.2.2 16-Bit and 32-bit Return Values

16-bit and 32-bit values are returned in temporary memory locations, with the least significant word in the lowest memory location. For example, the function:

```c
int return_16(void)
{
    return 0x1234;
}
```

will exit with the following code:

```
    movlw   34h
    movwf   btemp
    movlw   12h
    movwf   btemp+1
    return
```

4.6.2.3 Structure Return Values

Composite return values (struct and union) of size 4 bytes or smaller are returned in memory as with 16-bit and 32-bit return values. For composite return values of greater than 4 bytes in size, the structure or union is copied into the struct psect. Data is copied using the library routine structcopy which uses FSR0 for the source address, FSR1 for the destination address and W for the structure size. For example:

```c
struct fred
{
    int ace[4];
} ;

struct fred return_struct(void)
{
    struct fred wow;

    return wow;
}
```

will exit with the following code:

```
    movlw   low(?a_func+0)
    movwf   fsr0l
    movlw   high(?a_func+0)
```
Functions

movwf    fsr0h
movlw    structret
movwf    fsrl1
clrf     fsrlh
movlw    24

4.6.3 Memory Models and Usage

The compiler makes few assumptions about memory. With the exception of variables declared using the \texttt{@address} construct, absolute addresses are not allocated until link time.

The memory used is based upon information in the chipinfo file (which defaults to \texttt{pic18.ini} in the LIB directory). The linker will automatically locate code and \texttt{const}-qualified data into all the available memory pages and ensure that psects do not straddle any memory boundary.

Temporary variables created and used by the compiler are placed in the access bank to increase efficiency.

There are two memory models available for PICC-18: small and large. The default memory model is large. The memory model is selected via the \texttt{-Bx} command line option. See section 3.4.7 on page 51.

In large memory model, all objects qualified \texttt{near} are placed in a “near” psect (e.g. \texttt{rbss}, \texttt{rdata}) which are positioned in the access bank. These objects can be efficiently accessed using less generated code than other objects. All other objects are placed in the PIC18’s banked memory space in the following manner. Each module can allocate up to one bank of initialized (\texttt{data} psect), and one bank of uninitialized (\texttt{bss} psect), global or \texttt{static} local objects. Any object that is larger than one bank in size (e.g. an array) is placed in a separate area (in one of the “big” psects) and this area can grow across bank boundaries to a size limited only by the available space on the device. Any single byte objects are also placed in one of the “big” psects (“big” refers to the size of the psect, not the size of the objects within the psect). All \texttt{auto} and parameter variables from all functions are overlapped by the linker if possible and then placed into an available RAM bank.

In small model, all objects qualified \texttt{near} are placed in a “near” psect (e.g. \texttt{rbss}, \texttt{rdata}) which are positioned in the access bank as per the large model. The global and \texttt{static} local initialized and uninitialized objects are also placed in the access bank as are all single byte objects. Objects larger than the access bank size are positioned in a separate area (in one of the “big” psects) and this area can grow across bank boundaries to a size limited only by the available space on the device. All \texttt{auto} and parameter variables are positioned as per the large model.
4.7 Register Usage

The W register is used for register-based function argument passing and for function return values. This register should be preserved by any assembly language routines which are called.

4.8 Operators

PICC-18 supports all the ANSI operators. The exact results of some of these are implementation defined. The following sections illustrate code produced by the compiler.

4.8.1 Integral Promotion

When there is more than one operand to an operator, they typically must be of exactly the same type. The compiler will automatically convert the operands, if necessary, so they have the same type. The conversion is to a “larger” type so there is no loss of information. Even if the operands have the same type, in some situations they are converted to a different type before the operation. This conversion is called integral promotion. PICC-18 performs these integral promotions where required. If you are not aware that these changes of type have taken place, the results of some expressions are not what would normally be expected.

Integral promotion is the implicit conversion of enumerated types, signed or unsigned varieties of char, short int or bitfield types to either signed int or unsigned int. If the result of the conversion can be represented by an signed int, then that is the destination type, otherwise the conversion is to unsigned int.

Consider the following example.

```c
unsigned char count, a=0, b=50;
if(a - b < 10)
    count++;
```

The unsigned char result of a - b is 206 (which is not less than 10), but both a and b are converted to signed int via integral promotion before the subtraction takes place. The result of the subtraction with these data types is -50 (which is less than 10) and hence the body of the if() statement is executed. If the result of the subtraction is to be an unsigned quantity, then apply a cast. For example:

```c
if((unsigned int)(a - b) < 10)
    count++;
```

The comparison is then done using unsigned int, in this case, and the body of the if() would not be executed.

Another problem that frequently occurs is with the bitwise compliment operator, “~”. This operator toggles each bit within a value. Consider the following code.
```c
unsigned char count, c;
c = 0x55;
if( ~c == 0xAA)
    count++;
```

If `c` contains the value 55h, it often assumed that `~c` will produce AAh, however the result is FFAAh and so the comparison in the above example would fail. The compiler may be able to issue a mismatched comparison error to this effect in some circumstances. Again, a cast could be used to change this behaviour.

The consequence of integral promotion as illustrated above is that operations are not performed with `char`-type operands, but with `int`-type operands. However there are circumstances when the result of an operation is identical regardless of whether the operands are of type `char` or `int`. In these cases, PICC-18 will not perform the integral promotion so as to increase the code efficiency. Consider the following example.

```c
unsigned char a, b, c;
a = b + c;
```

Strictly speaking, this statement requires that the values of `b` and `c` should be promoted to `unsigned int`, the addition performed, the result of the addition cast to the type of `a`, and then the assignment can take place. Even if the result of the `unsigned int` addition of the promoted values of `b` and `c` was different to the result of the `unsigned char` addition of these values without promotion, after the `unsigned int` result was converted back to `unsigned char`, the final result would be the same. An 8-bit addition is more efficient than an a 16-bit addition and so the compiler will encode the former.

If, in the above example, the type of `a` was `unsigned int`, then integral promotion would have to be performed to comply with the ANSI standard.

### 4.8.2 Shifts applied to integral types

The ANSI standard states that the result of right shifting (`>>` operator) signed integral types is implementation defined when the operand is negative. Typically, the possible actions that can be taken are that when an object is shifted right by one bit, the bit value shifted into the most significant bit of the result can either be zero, or a copy of the most significant bit before the shift took place. The latter case amounts to a sign extension of the number.

PICC-18 performs a sign extension of any `signed` integral type (for example `signed char`, `signed int` or `signed long`). Thus an object with the `signed int` value 0124h shifted right one bit will yield the value 0092h and the value 8024h shifted right one bit will yield the value C012h.

Right shifts of `unsigned` integral values always clear the most significant bit of the result.

Left shifts (`<<` operator), `signed` or `unsigned`, always clear the least significant bit of the result.
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### 4.8.3 Division and modulus with integral types

The sign of the result of division with integers when either operand is negative is implementation specific. Table 4 - 8 shows the expected sign of the result of the division of operand 1 with operand 2 when compiled with PICC-18.

**Table 4 - 8 Integral division**

<table>
<thead>
<tr>
<th>Operand 1</th>
<th>Operand 2</th>
<th>Quotient</th>
<th>Remainder</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
</tr>
<tr>
<td>-</td>
<td>+</td>
<td>-</td>
<td>-</td>
</tr>
<tr>
<td>+</td>
<td>–</td>
<td>–</td>
<td>+</td>
</tr>
<tr>
<td>–</td>
<td>–</td>
<td>+</td>
<td>–</td>
</tr>
</tbody>
</table>

In the case where the second operand is zero (division by zero), the result will always be zero.

### 4.9 Psects

The compiler splits code and data objects into a number of standard program sections referred to as *psects*. The HI-TECH assembler allows an arbitrary number of named psects to be included in assembler code. The linker will group all data for a particular psect into a single segment.

If you are using **PICC18** to invoke the linker, you don’t need to worry about the information documented here, except as background knowledge. If you want to run the linker manually (this is not recommended), or write your own assembly language subroutines, you should read this section carefully.

A psect can be created in assembler code by using the **PSECT** assembler directive, see. In C, user-defined psects can be created by using the **#pragma psect** preprocessor directive.

#### 4.9.1 Compiler-generated Psects

The code generator places code and data into psects with standard names which are subsequent positioned by the default linker options. These psects are described below.

- **powerup** Which contains executable code for the standard or user-supplied power-up routine.
- **ida** These psects contain the ROM image of any initialised variables. These psects are copied into the **data** psects at startup.
- **irda** These psects contain the ROM image of any initialised **near** variables. These psects are copied into the **rdata** psects at startup.
**Psects**

**ibigdata** These psects contain the ROM image of initialised objects which at runtime reside in the **bigdata** psect. This includes global or static local **char** objects or **char** arrays, and arrays whose size exceeds the size of a RAM bank.

**text** These psects (where \( n \) is a number) contain all executable code. Typically \( n \) is incremented for each new C function compiled.

**text** Is a global psect used for executable code for some library functions.

**const** These psects hold objects that are declared **const** and string literals which are not modifiable.

**config** Used to store the configuration words.

**idloc** Used to store the ID location words.

**eeprom_data** Use to store data to be programmed into the EEPROM data area.

**intcode** Is the psect which contains the executable code for the default or high-priority interrupt service routine. This psect is linked to interrupt vector at address 08H.

**intcodelol** the psect which contains the executable code for the low-priority interrupt service routine. This psect is linked to interrupt vector at address 018H.

**init** Used by initialisation code which, for example, clears RAM.

**end_init** Used by initialisation code which, for example, clears RAM.

**clrtext** Used by some startup routines for copying the **data** psects.

The compiler-generated psects which are placed in RAM are:

**rbss** These psects contain any uninitialized **near** variables. They reside in the access bank.

**bigbss** These psects contain any uninitialized global or **static** local **char** objects or **char** arrays, and arrays whose size exceeds the size of a RAM bank. This psect is linked into a psect class which does not have RAM bank boundaries. Accessing objects in this area may be less efficient that accessing objects in the data psect.

**farbss** This psect contains any uninitialized objects which have been declared as far to be positioned in external code space. By default this psect is linked after the top of program memory.

**bss** These psects contain any uninitialized variables not contained in the above psects.

**rdata** These psects contain any initialised **near** variables. They reside in the access bank.

**bigdata** These psects contain any initialized global or **static** local **char** objects or **char** arrays, and arrays whose size exceeds the size of a RAM bank. This psect is linked into a psect class
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which does not have RAM bank boundaries. Accessing objects in this area may be less efficient that accessing objects in the data psect.

**data** These psects contain any initialised variables not contained in the above psects. These psects will be wholly placed within a RAM bank and so can be accessed more efficiently.

**nvrram** This psect holds **near persistent** variables. It is not cleared or otherwise modified by the runtime startup code.

**nvbit** This psect hold **persistent bit** objects. It is not cleared or otherwise modified at startup.

**nvram** This psect is used to store **persistent** variables. It is not cleared or otherwise modified at startup.

**rbit** These psects are used to store all **bit** variables. All bit objects are **near** by default and are placed in the access bank.

**struct** Contains any structure larger than 4 bytes in size which is returned from a function.

**intsave_regs** Holds the registers (including temporary locations) saved by the interrupt service routine.

**temp** Is used to store scratch variables used by the compiler. These include function return values larger than a byte and values passed to and returned from library routines. This psect will be positioned in the access bank.

### 4.10 Interrupt Handling in C

The compiler incorporates features allowing the PIC18 interrupts to be handled without writing any assembler code. The PIC18 devices have two separate interrupt vectors and a priority scheme to dictate how the interrupt code is called.

#### 4.10.1 Interrupt Functions

The function qualifier **interrupt** may be applied to at most two functions to allow them to be called directly from the hardware interrupts. The compiler will process the **interrupt** function differently to any other functions, generating code to save and restore any registers used and exit using the **retfie** instruction instead of a **retlw** or **return** instructions at the end of the function.

(If the PIC18 option `-STRICT` is used, the **interrupt** keyword becomes **__interrupt**. Wherever this manual refers to the **interrupt** keyword, assume **__interrupt** if you are using `-STRICT`.)

The PIC18 devices have two interrupts, each with their own vector location. These have different priorities and are known as **low-priority** and **high-priority interrupts**. If the PIC18 is placed in compatibility mode, only one interrupt is available and this defaults to being the high-priority interrupt. An **interrupt** function must be declared as type **interrupt void** and may not have parameters. In
addition, the keyword `low_priority` may be used to indicate that the `interrupt` function is to be linked with the low-priority vector when not in compatibility mode. `Interrupt` functions may not be called directly from C code, but they may call other functions itself, subject to certain limitations. Once defined, the corresponding interrupt vector is linked to the `interrupt` function.

An example of a high-priority (default) `interrupt` function is shown here.

```c
long tick_count;

void interrupt tc_int(void)
{
    ++tick_count;
}
```

A low-priority interrupt function may be defined as in the following example.

```c
void interrupt low_priority tc_clr(void)
{
    tick_count = 0;
}
```

It is up to the user to determine and set the priority levels associated with each interrupt source on the PIC18 devices. Defining a low-priority `interrupt` function does not put the PIC into interrupt-priority mode.

Low- and high-priority `interrupt` functions have their own separate areas of memory in which to save context, thus a high-priority `interrupt` function may interrupt a low-priority `interrupt` function with no loss of data. The high-priority `interrupt` can also employ the devices’ shadow registers to enable rapid context switching during the entry and exit of the service routine.

The `interrupt_level` pragma may be used with either or both `interrupt` functions in the usual way.

### 4.10.2 Context Saving on Interrupts

The PIC18 processor only saves the program counter on its stack whenever an interrupt occurs. Other registers and objects must be saved in software. PICC-18 automatically determines which registers and objects are used by an interrupt function and saves these appropriately.

If the interrupt routine calls other functions and these functions are defined before the interrupt code in the same module, then any registers used by these functions will be saved as well. If the called functions have not been seen by the compiler, a worst case scenario is assumed and all registers and objects will be saved.
PICC-18 does not scan assembly code which is placed in-line within the interrupt function for register usage. Thus, if you include in-line assembly code into an interrupt function, you may have to add extra assembly code to save and restore any registers or locations used if they are not already saved by the interrupt entry routine.

By default, the high-priority interrupt function will utilize a fast interrupt save/restore technique where the W, STATUS and BSR registers are saved and restored via the devices’ internal shadow registers. This minimizes code size and reduces the instruction cycles to access the high-priority service routine. Note that fast interrupt save/restore is not available if compiling for the MPLAB ICD2 debugger.

The high-priority or compatibility-mode interrupt function places a small routine in a psect called intcode which is linked directly to the interrupt vector. This code saves the STATUS (if fast interrupts are not used) and PCLATH registers then jumps to code placed in a text psect. This code will save further context if it is necessary and then jump to code directly related to the interrupt function. The interrupt function code is also placed in a text psect.

All objects saved are done so to locations at an offset to a symbol called saved_regsh, except for the BSR register. If fast interrupts are not used, BSR is saved to a location symbol called saved_bshr.

The low-priority interrupt function places the code to save the STATUS and PCLATH registers in a psect called intcodelo, which is directly linked to the low-priority interrupt vector. Operation is then similar to the high-priority interrupt case, only with objects being saved offset to the symbol saved_regsl and the BSR register saved to a location symbol called saved_bsr1.

**4.10.3 Context Retrieval**

Any objects saved by the compiler are automatically restored before the interrupt function returns. The restoration code is placed into a text psect. The retfie instruction placed at the end of the interrupt code will reload the program counter and the program will return to the location at which it was when the interrupt occurred.

**4.10.4 Interrupt Levels**

Normally it is assumed by the compiler that any interrupt may occur at any time, and an error will be issued by the linker if a function appears to be called by an interrupt function and by main-line code, or another interrupt. Since it is often possible for the user to guarantee this will not happen for a specific routine, the compiler supports an interrupt level feature to suppress the errors generated.

This is achieved with the #pragma interrupt_level directive. There are two interrupt levels available, and any interrupt functions at the same level will be assumed by the compiler to be mutually exclusive. This exclusion must be guaranteed by the user, i.e. the compiler is not able to control interrupt priorities. Each interrupt function may be assigned a single level, either 0 or 1.
In addition, any non-\texttt{interrupt} functions that are called from an \texttt{interrupt} function and also from main-line code may also use the \texttt{#pragma interrupt\_level} directive to specify that they will never be called by interrupts of one or more levels. This will prevent linker from issuing an error message because the function was included in more than one call graph. Note that it is entirely up to the user to ensure that the function is \textit{not} called by both main-line and interrupt code at the same time. This will normally be ensured by disabling interrupts before calling the function. It is not sufficient to disable interrupts inside the function after it has been called.

An example of using the interrupt levels is given below. Note that the \texttt{#pragma} directive applies to only the immediately following function. Multiple \texttt{#pragma interrupt\_level} directives may precede a non-interrupt function to specify that it will be protected from multiple interrupt levels.

```c
/* non-interrupt function called by interrupt and main-line code */
#pragma interrupt\_level 1
void bill()
{
    int i;
    i = 23;
}

/* two interrupt functions calling the same non-interrupt function */
#pragma interrupt\_level 1

void interrupt fred(void)
{
    bill();
}

#pragma interrupt\_level 1
void interrupt joh()
{
    bill();
}

main()
{
    bill();
}
```

Both the low- and high-priority \texttt{interrupt} functions may use the interrupt level feature.
4.10.5 Interrupt Registers

It is up to the user how they want the interrupt source configured. All the registers and bits associated with interrupts are defined in the specific header file which can be accessed by including `<pic18.h>`. The following is an example of setting up the interrupts associated with the change-on-PORTB source. Interrupt priorities are used and the interrupt source is made a low priority. See your PIC18 datasheet for more information.

```c
void main(void)
{
    TRISB = 0x80;  // only RB7 will interrupt on change
    IPEN = 1;     // interrupt priorities enabled
    PEIE = 1;     // enable peripheral interrupts
    RBIP = 0;     // make this a low priority interrupt
    RBIE = 1;     // enable PORTB change interrupts
    RBIF = 0;     // clear any pending events
    GIEL = 1;     // enable low-priority interrupts
    while(1)
    {
    }
}

void interrupt low_priority_B_change(void)
{
    if(RBIF && RBIE) {
        PORTB;          // read PORTB to clear any mismatch
        RBIF = 0;       // clear the flag
        // process interrupt here
    }
}
```

4.11 Mixing C and Assembler Code

Assembly language code can be mixed with C code using three different techniques.

4.11.1 External Assembly Language Functions

Entire functions may be coded in assembly language as separate `.as` source files, assembled by the assembler (`ASPIC`) and combined into the binary image using the linker. This technique allows arguments and return values to be passed between C and assembler code.
To access an external function, first include an appropriate C `extern` declaration in the calling C code. For example, suppose you have an assembly language function to double a unsigned byte value:

```c
extern char    twice(char);
```

declares an external function called `twice()` which has a return value type of `char` and takes a single argument of type `char`. The actual code for `twice()` will be supplied by an external `.as` file which will be separately assembled with ASPIC. The full PIC18 assembler code for `twice()` would be something like:

```assembly
PROCESSOR    18C452

PSECT text0,class=CODE,local,delta=1
GLOBAL _twice
SIGNAT _twice,4201

_twice:
    ; parameter is passed in the W register - assign it to ?a_twice.
    movlb   ?a_twice shr (8)  ; select local bank
    movwf   ?a_twice & 0ffh
    addwf   ?a_twice & 0ffh,w
    ; The return is already in the W register as required.
    return

FNSIZE _twice,1,0
GLOBAL ?a_twice
END
```

The name of the assembly language function is the name declared in C, with an underscore prepended. The `GLOBAL` pseudo-op is the assembler equivalent to the C `extern` keyword and the `SIGNAT` pseudo-op is used to enforce link time calling convention checking. Signature checking and the `SIGNAT` pseudo-op are discussed in more detail later in this chapter.

Note that in order for assembly language functions to work properly they must look in the right place for any arguments passed and must correctly set up any return values. Local variable allocation (via the `FNSIZE` directive), argument and return value passing mechanisms are discussed in the Section Functions on page 94 and should be understood before attempting to write assembly language routines.

### 4.11.2 Accessing C objects from within assembler

Global C objects may be directly accessed from within assembly code using their name prepended with an underscore character. For example, the object `foo` defined globally in a C module:
int foo;

may be access from assembler as follows.

    GLOBAL _foo
    movwf _foo

If the assembler is contained in a different module, then the `GLOBAL` assembler directive should be used in the assembler code to make the symbol name available, as above. If the object is being accessed from in-line assembly in another module, then an `extern` declaration for the object can be made in the C code, for example:

    extern int foo;

This declaration will only take effect in the module if the object is also accessed from within C code. If this is not the case then, an in-line `GLOBAL` assembler directive should be used. Care should be taken in the object is defined in a bank other than 0. The address of a C object includes the bank information which must be stripped before the address can be used in most PIC18 instructions. The exceptions are the `movff` and `lsfr` instructions. Failure to do this may result in fixup errors issued by the linker. If in doubt as to writing assembler which access C objects, write code in C which performs a similar task to what you intend to do and study the assembler listing file produced by the compiler.

### 4.11.3 #asm, #endasm and asm()

PIC18 instructions may also be directly embedded in C code using the directives `#asm`, `#endasm` and the statement `asm()`. The `#asm` and `#endasm` directives are used to start and end a block of assembler instructions which are to be embedded inside C code. The `asm()` statement is used to embed a single assembler instruction in the code generated by the C compiler. The following example shows both methods used to rotate a byte left through carry:

```c
unsigned char var;
void main(void)
{

    var = 1;
    #asm    // like this...
        movlb (_var) >> 8
        rlcf (_var)&0ffh,f
    #endasm  // or like this
        asm("movlb (_var) >> 8");
        asm("rlcf (_var)&0ffh,f");
}
```
When using in-line assembler code, great care must be taken to avoid interacting with compiler-generated code. If in doubt, compile your program with the `PICC18 -S` option and examine the assembler code generated by the compiler.

IMPORTANT NOTE: the `#asm` and `#endasm` construct is not syntactically part of the C program, and thus it does not obey normal C flow-of-control rules. For example, you cannot use a `#asm` block with an `if` statement and expect it to work correctly. If you use in-line assembler around any C constructs such as `if`, `while`, `do` etc. they you should use only the `asm("\")` form, which is a C statement and will correctly interact with all C flow-of-control structures.

### 4.12 Preprocessing

All C source files are preprocessed before compilation. Assembler files can also be preprocessed if the `-p` command-line option is issued.

#### 4.12.1 Preprocessor Directives

PICC-18 accepts several specialised preprocessor directives in addition to the standard directives. These are listed in Table 4 - 9 on page 110.

Macro expansion using arguments can use the `#` character to convert an argument to a string, and the `##` sequence to concatenate tokens.

#### 4.12.2 Predefined Macros

The compiler drivers define certain symbols to the preprocessor (CPP), allowing conditional compilation based on chip type etc. The symbols listed in Table 4 - 10 on page 111 show the more common symbols defined by the drivers. Each symbol, if defined, is equated to 1 unless otherwise stated.

#### 4.12.3 Pragma Directives

There are certain compile-time directives that can be used to modify the behaviour of the compiler. These are implemented through the use of the ANSI standard `#pragma` facility. The format of a pragma is:

```
#pragma keyword options
```

where `keyword` is one of a set of keywords, some of which are followed by certain `options`. A list of the keywords is given in Table 4 - 11 on page 112. Those keywords not discussed elsewhere are detailed below.

#### 4.12.3.1 The `#pragma jis` and `nojis` Directives

If your code includes strings with two-byte characters in the JIS encoding for Japanese and other national characters, the `#pragma jis` directive will enable proper handling of these characters, specifically not interpreting a backslash “\” character when it appears as the second half of a two byte
### Table 4 - 9 Preprocessor directives

<table>
<thead>
<tr>
<th>Directive</th>
<th>Meaning</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>#</code></td>
<td>preprocessor null directive, do nothing</td>
<td><code>#</code></td>
</tr>
<tr>
<td><code>#assert</code></td>
<td>generate error if condition false</td>
<td><code>#assert SIZE &gt; 10</code></td>
</tr>
</tbody>
</table>
| `#asm`    | signifies the begining of in-line assembly | `#asm
movlw 10h
#endasm` |
| `#define` | define preprocessor macro | `#define SIZE 5
#define FLAG
#define add(a,b) ((a)+(b))` |
| `#elif`   | short for `#else` `#if` | see `#ifdef` |
| `#else`   | conditionally include source lines | see `#if` |
| `#endasm` | terminate in-line assembly | see `#asm` |
| `#endif` | terminate conditional source inclusion | see `#if` |
| `#error`  | generate an error message | `#error Size too big` |
| `#if`     | include source lines if constant expression true | `#if SIZE < 10
c = process(10)
#else
skip();
#endif` |
| `#ifndef` | include source lines if preprocessor symbol not defined | `#ifndef FLAG
jump();
#endif` |
| `#ifdef`   | include source lines if preprocessor symbol defined | `#ifdef FLAG
do_loop();
#else
skip_loop();
#endif` |
| `#ifndef` | include source lines if preprocessor symbol not defined | `#ifndef FLAG
jump();
#endif` |
| `#include` | include text file into source | `#include <stdio.h>
#include "project.h"` |
| `#line`   | specify line number and filename for listing | `#line 3 final` |
| `#nn`     | (where `nn` is a number) short for `#line nn` | `#20` |
| `#pragma` | compiler specific options | See section 4.12.3 on page 109 |
| `#undef`  | undefines preprocessor symbol | `#undef FLAG` |
Preprocessing

Table 4-9 Preprocessor directives

<table>
<thead>
<tr>
<th>Directive</th>
<th>Meaning</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>#warning</td>
<td>generate a warning message</td>
<td>#warning Length not set</td>
</tr>
</tbody>
</table>

character. The nojis directive disables this special handling. JIS character handling is disabled by default.

Table 4-10 Predefined CPP Symbols

<table>
<thead>
<tr>
<th>Symbol</th>
<th>When set</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>HI_TECH_C</td>
<td>Always</td>
<td>To indicate that the compiler in use is HI-TECH C.</td>
</tr>
<tr>
<td><em>HTC_VER_MAJOR</em></td>
<td>Always</td>
<td>To indicate the integer component of the compiler’s version number.</td>
</tr>
<tr>
<td><em>HTC_VER_MINOR</em></td>
<td>Always</td>
<td>To indicate the decimal component of the compiler’s version number.</td>
</tr>
<tr>
<td><em>HTC_VER_PATCH</em></td>
<td>Always</td>
<td>To indicate the patch level of the compiler’s version number.</td>
</tr>
<tr>
<td>LARGE_DATA</td>
<td>-CP24</td>
<td>To indicate that program space pointers are 24 bits in size.</td>
</tr>
<tr>
<td>SMALL_DATA</td>
<td>-CP16</td>
<td>To indicate that program space pointers are 16 bits in size.</td>
</tr>
<tr>
<td>LARGE_MODEL</td>
<td>-Bl</td>
<td>To indicate that code is compiled in large memory model.</td>
</tr>
<tr>
<td>SMALL_MODEL</td>
<td>-Bs</td>
<td>To indicate that code is compiled in small memory model.</td>
</tr>
<tr>
<td><em>MPC</em></td>
<td>Always</td>
<td>To indicate the code is compiled for the Microchip PIC family.</td>
</tr>
<tr>
<td>_PIC18</td>
<td>Always</td>
<td>To indicate that this is a PIC18 device.</td>
</tr>
<tr>
<td>_18CXXX</td>
<td>When chip selected</td>
<td>To indicate the specific chip type selected</td>
</tr>
<tr>
<td>MPLAB_ICD</td>
<td>-ICD</td>
<td>To indicate that code is being generated for the MPLAB In-Circuit Debugger</td>
</tr>
</tbody>
</table>

4.12.3.2 The #pragma printf_check Directive

Certain library functions accept a format string followed by a variable number of arguments in the manner of printf(). Although the format string is interpreted at run-time, it can be compile-time checked for consistency with the remaining arguments. This directive enables this checking for the
named function, e.g. the system header file `<stdio.h>` includes the directive `#pragma printf_check(printf) const` to enable this checking for `printf()`. You may also use this for any user-defined function that accepts printf-style format strings. The qualifier following the function name is to allow automatic conversion of pointers in variable argument lists. The above example would cast any pointers to strings in RAM to be pointers of the type `(const char *)`

Note that the warning level must be set to -1 or below for this option to have effect.

### 4.12.3.3 The `#pragma psect` Directive

Normally the object code generated by the compiler is broken into the standard psects as already documented. This is fine for most applications, but sometimes it is necessary to redirect variables or code into different psects when a special memory configuration is desired. Code and data for any of the standard C psects may be redirected using a `#pragma psect` directive. For example, if all the uninitialised global data in a particular C source file is to be placed into a psect called `otherram`, the following directive should be used:

```
#pragma psect bss=otherram
```

This directive tells the compiler that anything which would normally be placed in the `bss` psect should now be placed in the `otherram` psect.

<table>
<thead>
<tr>
<th>Directive</th>
<th>Meaning</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>interrupt_level</td>
<td>Allow interrupt function to be called from main-line code. See section 4.10.4 on page 104</td>
<td><code>#pragma interrupt_level 1</code></td>
</tr>
<tr>
<td>jis</td>
<td>Enable JIS character handling in strings</td>
<td><code>#pragma jis</code></td>
</tr>
<tr>
<td>nojis</td>
<td>Disable JIS character handling (default)</td>
<td><code>#pragma nojis</code></td>
</tr>
<tr>
<td>printf_check</td>
<td>Enable printf-style format string checking</td>
<td><code>#pragma printf_check(printf) const</code></td>
</tr>
<tr>
<td>psect</td>
<td>Rename compiler-defined psect</td>
<td><code>#pragma psect text=mytext</code></td>
</tr>
<tr>
<td>regsused</td>
<td>Specify registers which are used in an interrupt</td>
<td><code>#pragma regsused w</code></td>
</tr>
<tr>
<td>switch</td>
<td>Specify code generation for switch statements</td>
<td><code>#pragma switch direct</code></td>
</tr>
</tbody>
</table>

Placing code in a different psect is done by redirecting the text psect. To redirect code the following preprocessor directive can be used.

```
#pragma psect text=othercode
```
where `othercode` is the name of the new psect to be created and filled.

This example will define the psect `othercode0` for `function()`’s code, and `othercode1` for `another()`’s code.

Any given psect should only be redirected once in a particular source file, and all psect redirections for a particular source file should be placed at the top of the file, below any `#include` statements and above any other declarations. For example, to declare a group of uninitialized variables which are all placed in a psect called `otherram`, the following technique should be used:

```
--File OTHERRAM.C
#pragma psect bss=otherram
char buffer[5];
int var1, var2, var3;
```

Any files which need to access the variables defined in `otherram.c` should `#include` the following header file:

```
--File OTHERRAM.H
extern char buffer[5];
extern int var1, var2, var3;
```

The `#pragma psect` directive allows code and data to be split into arbitrary memory areas. Definitions of code or data for non-standard psects should be kept in separate source files as documented above. When linking code which uses non-standard psect names, you will need to use the `PICC18 -L` option to specify an extra linker option, drive the linker manually. If you want a nearly standard configuration with the addition of only an extra psect like `otherram`, you can use the `PICC18 -L` option to add an extra `-P` specification to the linker command. For example:

```
PICC18 -L-Potherram=200h -18C452 test.obj otherram.obj
```

will link `test.obj` and `otherram.obj` with a standard configuration, and the extra `otherram` psect at 200h in RAM.

### 4.12.3.4 The `#pragma regsused` Directive

PICC-18 will automatically save context when an interrupt occurs. The compiler will determine only those registers and objects which need to be saved for the particular interrupt function defined. The `#pragma regsused` directive allows the programmer to further limit the registers and objects that the compiler might save and retrieve on interrupt.

Table 4 - 12 on page 114 shows registers names that would commonly be used with this directive. The register names are not case sensitive and a warning will be produced if the register name is not recognised.
This pragma affects the first interrupt function following in the source code. Code for High-End devices which contains multiple interrupt functions should include one directive for each interrupt function.

For example, to limit the compiler to saving no registers other than the W register and FSR register for an interrupt function, use:

```
#pragma regsused w fsr
```

Even if a register, other than W or FSR, has been used and that register would normally be saved, it will not be saved if this pragma is in effect. The W and/or FSR register will only be automatically saved by the compiler if required.

### 4.12.3.5 The #pragma switch Directive

Normally the compiler decides the code generation method for switch statements which results in the smallest possible code size. Specifying the `direct` option to the `#pragma switch` directive forces the compiler to generate the table look-up style switch method. This is mostly useful where either timing or code size is an issue for switch statements (ie: state machines) and a jump table is preferred over direct comparison or vice versa. This pragma affects all code generated onwards. The `auto` option may be used to revert to the default behaviour.

### 4.13 Linking Programs

The compiler will automatically invoke the linker unless requested to stop after producing assembler code (`PICC18 -S` option) or object code (`PICC18 -C` option).

`PICC18`, by default, generates *Intel HEX* files and *Bytecraft COD*. If you use the `-BIN` option or specify an output file with a `.BIN` file-type using the `PICC18 -O` option the compiler will generate a binary image instead. After linking, the compiler will automatically generate a memory usage map which shows the
address used by, and the total sizes of, all the psects which are used by the compiled code. Note that \textbf{bit} objects are shown separately. For example:

Memory Usage Map:

<table>
<thead>
<tr>
<th>Program ROM</th>
<th>$000000 - $000003</th>
<th>$000004 (4) bytes</th>
</tr>
</thead>
<tbody>
<tr>
<td>Program ROM</td>
<td>$000018 - $00006F</td>
<td>$000058 (88) bytes</td>
</tr>
<tr>
<td>Program ROM</td>
<td>$000086 - $0001CB</td>
<td>$000146 (326) bytes</td>
</tr>
<tr>
<td></td>
<td>$0001A2</td>
<td>418 bytes total Program ROM</td>
</tr>
</tbody>
</table>

Program statistics:

- Total ROM used: 466 bytes (1.4%)
- Total RAM used: 64 bytes (4.2%) Near RAM used: 18 bytes (14.1%)

The program statistics shown after the summary provides more concise information based on each memory area of the device. This can be used as a guide to the available space left in the device. The access bank is shown in this summary which is indicated as near RAM. This is the memory area used by near-qualified objects. Such objects should be used wherever possible to reduce code size.

More detailed memory usage information, listed in ascending order of individual psects, may be obtained by using the \texttt{PICC18 -PSECTMAP} option.

\subsection*{4.13.1 Replacing Library Modules}

Although PICC-18 comes with a librarian (\texttt{LIBR}) which allows you to unpack a library files and replace modules with your own modified versions, you can easily replace a module within a library without having to do this. If you add the source file which contains the library routine you wish to replace on the command-line list of source files then the routine will replace the routine in the library file with the same name. For example, if you wished to make changes to the library function \texttt{max()} which resides in the file \texttt{max.c} in the SOURCES directory, you could make a copy of this source file, make the appropriate changes and then compile and use it as follows.

\begin{verbatim}
PICC18 -18c452 main.c init.c max.c
\end{verbatim}

The code for \texttt{max()} in \texttt{max.c} will be linked into the program rather than the \texttt{max()} function contained in the standard libraries. Note, that if you replace an assembler module, you may need the \texttt{-P} option to preprocess assembler files as the library assembler files often contain C preprocessor directives.

\subsection*{4.13.2 Signature Checking}

The compiler automatically produces signatures for all functions. A signature is a 16-bit value computed from a combination of the function’s return data type, the number of its parameters and other
information affecting the calling sequence for the function. This signature is output in the object code of any function referencing or defining the function.

At link time the linker will report any mismatch of signatures. Thus if a function is declared in one module in a different way (for example, as returning a `char` instead of `short`) then the linker will report an error.

It is sometimes necessary to write assembly language routines which are called from C using an `extern` declaration. Such assembly language functions need to include a signature which is compatible with the C prototype used to call them. The simplest method of determining the correct signature for a function is to write a dummy C function with the same prototype and compile it to assembly language using the `PICC18 -S` option. For example, suppose you have an assembly language routine called `_widget` which takes two `int` arguments and returns a `char` value. The prototype used to call this function from C would be:

```c
extern char widget(int, int);
```

Where a call to `_widget` is made in the C code, the signature for a function with two `int` arguments and a `char` return value would be generated. In order to match the correct signature the source code for `widget` needs to contain an ASPIC `SIGNAT` pseudo-op which defines the same signature value. To determine the correct value, you would write the following code:

```c
char widget(int arg1, int arg2)
{
}
```

and compile it to assembler code using

```
PICC18 -S x.c
```

The resultant assembler code includes the following line:

```c
SIGNAT _widget,8297
```

The `SIGNAT` pseudo-op tells the assembler to include a record in the `.obj` file which associates the value 8297 with symbol `_widget`. The value 8297 is the correct signature for a function with two `int` arguments and a `char` return value. If this line is copied into the `.as` file where `_widget` is defined, it will associate the correct signature with the function and the linker will be able to check for correct argument passing. For example, if another `.c` file contains the declaration:

```c
extern char widget(long);
```

then a different signature will be generated and the linker will report a signature mis-match which will alert you to the possible existence of incompatible calling conventions.
4.13.3 Linker-Defined Symbols

The link address of a psect can be obtained from the value of a global symbol with name \_Lname where name is the name of the psect. For example, \_Lbss is the low bound of the bss psect. The highest address of a psect (i.e. the link address plus the size) is symbol \_Hname. If the psect has different load and link addresses, as may be the case if the data psect is linked for RAM operation, the load address is \_Bname.

4.14 Standard I/O Functions and Serial I/O

A number of the standard I/O functions are provided in the C library with the compiler, specifically those functions intended to read and write formatted text on standard output and input. A list of the available functions is in Table 4 - 13. More details of these functions are in the Library Functions chapter.

Table 4 - 13 Supported STDIO Functions

<table>
<thead>
<tr>
<th>Function name</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>printf(const char * s, ...)</td>
<td>Formatted printing to stdout</td>
</tr>
<tr>
<td>sprintf(char * buf, const char * s, ...)</td>
<td>Writes formatted text to buf</td>
</tr>
</tbody>
</table>

Before any characters can be written or read using these functions, the putch() and getch() functions must be written. Other routines which may be required include getche() and kbhit().

You will find samples of serial code which implements the putch() and getch() functions in the file serial.c in the SAMPLES directory.

4.15 Debugging Information

4.15.1 MPLAB-specific information

Certain options and compiler features are specifically intended to help MPLAB perform symbolic debugging. The -FAKELOCAL switch performs two functions, both specific to MPLAB. Since MPLAB does not read the local symbol information produced by the compiler, this options generates additional global symbols which can be used to represent most local symbols in a program. The format for the symbols is function_name.symbol_name. Thus, if a variable called foo was defined inside the function main(), MPLAB would allow access to a global object called main.foo. This symbol format is not available in assembler code. References to this object in assembler would be via the symbol _main$foo. Although this information allows access to most local objects, if there are two or more objects with the same name in the same function, then you will not be able to examine both as they redefine the same symbol.

The -FAKELOCAL switch also alters the line numbering information produced so that MPLAB can better follow the C source when performing source-level stepping.
This option also adjusts the format for compiler errors so that they can be more readily interpreted by the MPLAB IDE. This error format can still be overridden by setting the appropriate environment variables as described in section -E: Define Format for Compiler Errors on page 53.
**PICC-18 Macro Assembler**

The HI-TECH PICC-18 Macro Assembler assembles source files for the Microchip PIC18 series of microprocessors. This chapter describes the usage of the assembler and the directives (assembler pseudo-ops and controls) accepted by the assembler.

The PICC-18 Macro Assembler package includes a linker, librarian, cross reference generator and an object code converter.

### 5.1 Assembler Usage

The assembler is called ASPIC18 and is available to run on PC and UNIX machines. Note that the assembler will not produce any messages unless there are errors or warnings - there are no “assembly completed” messages.

The usage of the assembler is similar under all of these operating systems. All command line options are recognised in either upper or lower case. The basic command format is shown:

```
ASPIC18 [ options ] files ...
```

*Files* is a space-separated list of one or more assembler source files. Where more than one source file is specified the assembler treats them as a single module, i.e. a single assembly will be performed on the concatenation of all the source files specified. The files must be specified in full, no default extensions or suffixes are assumed.

*Options* is an optional space-separated list of assembler options, each with a minus sign “-” as the first character. A full list of possible options is given in Table 5 - 1 on page 120, and a full description of each option follows.

### 5.2 Assembler Options

The command line options recognised by ASPIC18 are as follows:

- **-processor** This option defines the processor which is being used. You can also add your own processors to the compiler. For more information about this, See “Processor Support” on page 67.

- **-A** An assembler file with an extension .opt will be produced if this option is used. This is useful when checking the optimized assembler produced using the -O option.

- **-C** A cross reference file will be produced when this option is used. This file, called srcfile.crf where srcfile is the base portion of the first source file name, will contain
raw cross reference information. The cross reference utility CREF must then be run to produce the formatted cross reference listing.

-Chipinfo Define the chipinfo file to use. This option is not normally required as the chipinfo file is normally not used. The chip in file is used by the code generator and is called pic18.ini and can be found in the LIB directory of the compiler distribution.

-E The default format for an error message is in the form:

    filename: line: message

where the error of type message occurred on line line of the file filename. The –E2 option will produce a less-readable format suitable for some development environments.

-Flength By default the listing format is pageless, i.e. the assembler output is continuous. The output may be formatted into pages of varying lengths. Each page will begin with a header and title, if specified. The –F option allows a different page length to be
specified. A zero value of length implies pageless output. Length is specified in a number of lines.

-H Particularly useful in conjunction with the -A option, this option specifies that output constants should be shown as hexadecimal values rather than decimal values.

-I This option forces listing of macro expansions and unassembled conditionals which would otherwise be suppressed by a NOLIST assembler control. The -L option is still necessary to produce a listing.

-L listfile This option requests the generation of an assembly listing file. If listfile is specified then the listing will be written to that file, otherwise it will be written to the standard output.

-O This requests the assembler to perform optimization on the assembly code. Note that the use of this option slows the assembly down, as the assembler must make an additional pass over the input code.

-O outfile By default the assembler determines the name of the object file to be created by stripping any suffix or extension (i.e. the portion after the last dot) from the first source filename and appending .obj. The -O option allows the user to override the default and specify and explicit filename for the object file.

-S If a byte-size memory location is initialized with a value which is too large to fit in 8 bits, then the assembler will generate a Size error message. Use of the -S option will suppress this type of message.

-U Undefined symbols encountered during assembly are treated as external, however an error message is issued for each undefined symbol unless the -U option is given. Use of this option suppresses the error messages only, it does not change the generated code.

-V This option will include in the object file produced by the assembler, line number and file name information for the use of a debugger. Note that the line numbers will be assembler code lines - when assembling a file produced by the compiler, there will be line and file directives inserted by the compiler so this option is not required.

-W width This option allows specification of the listfile paper width, in characters. Width should be a decimal number greater than 41. The default width is 80 characters.

-X The object file created by the assembler contains symbol information, including local symbols, i.e. symbols that are neither public or external. The -X option will prevent the local symbols from being included in the object file, thereby reducing the file size.
5.3 PICC-18 Assembly Language

The source language accepted by the HI-TECH Software PICC-18 Macro Assembler is described below. All opcode mnemonics and operand syntax are strictly PIC18 assembly language. Additional mnemonics are documented in this section.

5.3.1 Assembler Format Deviations

The HI-TECH PICC-18 assembler uses a slightly modified form of assembly language to that specified by Microchip. Certain PIC18 instructions used by Microchip assembler use the operands "0", "1" to specify the destination for the result of that operation. The HI-TECH PICC-18 assembler uses the more-readable operands "w", "f" to specify the destination register. The W register is selected as the destination when using the "w" operand, and the file register is selected when using the "f" operand or if no destination operand is specified. The case of the letter in the destination operand in not important. The Microchip numerical operands cannot be used with the HI-TECH PICC-18 assembler.

The PICC-18 assembler also uses the operands "b", "c" to indicate that a file register is banked or common. A common register is one that resides in the access bank. Instructions using this operand will have the RAM access bit in the instruction cleared by the assembler. A banked register does not reside in the access bank. Instructions using this operand will have the RAM access bit in the instruction set by the assembler. The BSR register must be correctly loaded prior to executing a banked instruction to select the appropriate bank. Identifiers that do not use either of these operands are assumed to be banked.

Reference to a symbol in PICC-18 assembler may be preceded by the characters "c:" to indicate that the symbol resides is common. A common object resides in the access bank.

An access bank indicator, such as "c" or "c:" is not required when an address used in an instruction is absolute and the value of the address is within the access bank. The assembler will determine from the address that this is the case. However, these indicators must be used with all unresolved identifiers. For example, the following instructions show the WREG being moved to first, an absolute location and then to an address represented by an identifier. The op codes for these instructions, assuming that the address assigned to _foo is 0516h, are shown.

\[
\begin{align*}
6E16 & \text{ movwf } _\text{foo},w \\
6E16 & \text{ movwf } _\text{foo},f \\
6F16 & \text{ movwf } _\text{foo},b \\
6F16 & \text{ movwf } _\text{foo}
\end{align*}
\]

Notice that first two instruction have the RAM access bit (bit 8 of the op-code) cleared, but that it is set in the last two instructions.
The **retfie** instruction may be followed by “f” to indicate that the shadow registers should be retrieved and copied to their corresponding registers on execution.

**Special Comment Strings**

Several comment strings are appended to assembler instructions by the code generator. These are typically used by the assembler optimizer.

The comment string **;volatile** is used to indicate that the memory location being accessed is associated with a variable that was declared as **volatile** in the C source code. Accesses to this location which appear to be redundant will not be removed by the assembler optimizer if this string is present.

The comment string **;wreg free** is placed on some CALL instructions. The string indicates that the WREG was not loaded with a function parameter, i.e. it is not in use. If this string is present, optimizations may be made to assembler instructions before the function call which load the WREG redundantly.

**5.3.2 Pre-defined Macros**

The file **sfr.h**, contained in the SOURCES directory contains useful definitions for assembler programming. In particular it contains an assembler macro called **loadfsr**, which can be used when you require any of the FSR registers to be loaded. The two arguments to this macro are the FSR register number and the value to be loaded. For example:

```
loadfsr 2,1FFh
```

which will load FSR2 with the value 1FFh. This macro should be used in preference to the **lsfr** instruction.

**5.3.3 Character Set**

The character set used is standard 7 bit ASCII. Alphabetic case is significant for identifiers, but not opcodes and reserved words. Tabs are treated as equivalent to spaces.

**5.3.4 Constants**

**5.3.4.1 Numeric Constants**

The assembler performs all arithmetic as signed 32 bit. Errors will be caused if a quantity is too large to fit in a memory location. The default radix for all numbers is 10. Other radices may be specified by a trailing base specifier as given in Table 5 - 2.

Hexadecimal numbers must have a leading digit (e.g. 0xffffh) to differentiate them from identifiers. Hexadecimal constants are accepted in either upper or lower case.

Note that a binary constant must have an upper case **B** following it, as a lower case **b** is used for temporary (numeric) label backward references.
In expressions, real numbers are accepted in the usual format, and are interpreted as IEEE 32-bit format. A real number may be converted into the truncated IEEE 24-bit format by using the `float24` pseudo-function. Here is an example of its use:

```
movlw low(float24(31.415926590000002))
```

### 5.3.4.2 Character Constants

A character constant is a single character enclosed in *single quotes* “’”. Multi-character constants may be specified using *double quotes* “””. See “Strings” on page 127.

### 5.3.5 Delimiters

All numbers and identifiers must be delimited by white space, non-alphanumeric characters or the end of a line.

### 5.3.6 Special Characters

There are a few characters that are special in certain contexts. Within a macro body, the character “&” is used for token concatenation. To use the bitwise & operator within a macro body, escape it by using `&` instead. In a macro argument list, the *angle brackets* `<` and `>` are used to quote macro arguments.

### 5.3.7 Identifiers

Identifiers are user-defined symbols representing memory locations or numbers. A symbol may contain any number of characters drawn from the alphabetics, numerics and the special characters `dollar`, “$”, *question mark*, “?” and *underscore*, “_”. The first character of an identifier may not be numeric. The case of alpahabetics is significant, e.g. `Fred` is not the same symbol as `fred`. Some examples of identifiers are shown here:

```
An_identifier
an_identifier
an_identifier1
$$$
?$$_{12345}
```
5.3.7.1 Significance of Identifiers
Users of other assemblers that attempt to implement forms of data typing for identifiers should note that this assembler attaches no significance to any symbol, and places no restrictions or expectations on the usage of a symbol. The names of psects (program sections) and ordinary symbols occupy separate, overlapping name spaces, but other than this, the assembler does not care whether a symbol is used to represent bytes, words or chicken sheds. No special syntax is needed or provided to define the addresses of bits or any other data type, nor will the assembler issue any warnings if a symbol is used in more than one context. The instruction and addressing mode syntax provide all the information necessary for the assembler to generate correct code.

5.3.7.2 Assembler-Generated Identifiers
Where a LOCAL directive is used in a macro block, the assembler will generate a unique symbol to replace each specified identifier in each expansion of that macro. These unique symbols will have the form ??nnnn where nnnn is a 4 digit number. The user should avoid defining symbols with the same form.

5.3.7.3 Location Counter
The current location within the active program section is accessible via the symbol “$”. This symbol expands to the address of the currently executing instruction. The PIC18 PC register will contain the address of the instruction following the currently executing instruction so $ will not be the same as the PC register. Thus:

    goto $ 

will represent code that will form an endless loop.

The address represented by $ is a byte address, the same as symbols used with the PIC18 compiler. When determining an offset for this address, you must take into account the size of each instruction as a byte quantity. Thus the goto in the following:

    goto  $+6 ; size 4 bytes (2 words)
    movlw 55h ; size 2 bytes (1 word)
    movwf _foo

will skip the movlw instruction.

5.3.7.4 Register Symbols
When the code generator compiles a C module, it includes a list of EQU directives for some of the more commonly used SFRs. Thus any assembler that is placed in-line into a C module can uses these register names. This list of registers can be seen when generating an assembler list file. If in-line assembler code uses registers other than those listed, then an EQU definition for those registers must be included.
If writing true assembler modules, these SFR definitions will not be present since the code generator does not process assembler files. Thus to use any SFRs, an EQU definition must be included into then module.

Another way of using the SFRs is to linker in symbols with the C definitions for the SFRs that are included in the chip specific header file. Whenever you include `<pic18.H>` into a C module, all the available SFRs are defined as absolute C variables. This file cannot be included into an assembler module, but assembler code can uses these definitions. To use a SFR in in-line assembler code from within the same C module that includes `<pic18.h>`, simply use the symbol with an underscore character prepended to the name. For example:

```c
#include <pic18.h>
void main(void)
{
    asm("movff wreg,_PORTC");
}
```

To use these definitions from an assembler module you need to place a `GLOBAL` directive for the symbols that are to be used in both the C module that includes `<pic18.h>` and in the assembler module that uses the symbols. For example, in the C module you will need:

```c
#include <pic18.h>
asm("GLOBAL _PORTC");
```

and in the assembler module you will need:

```assembly
GLOBAL _PORTC
GLOBAL wreg
psect text, class=CODE, reloc=2
movff wreg,_PORTC
```

Note that `wreg` is an exception. Although a `GLOBAL` declaration is required in the assembler module, the leading `underscore` character is not required and the declaration is not required at all in the C module.

It is not possible to equate a symbol to a register.

### 5.3.7.5 Symbolic Labels

A label is a name at the beginning of a statement which is assigned a value equal to the current offset within the current psect (program section). A label is not the same as a macro name, which also appears at the beginning of the line in a macro declaration.

A label may be any symbol followed by a `colon`, `:`. Here are two examples of legitimate labels:

```assembly
frank:

simon44:
```
5.3.8 Strings

A string is a sequence of characters not including carriage return or newline characters, enclosed within matching quotes. Either single quotes ‘’ or double quotes “” may be used, but the opening and closing quotes must be the same. A string used as an operand to a DB directive may be any length, but a string used as operand to an instruction must not exceed 1 or 2 characters, depending on the size of the operand required.

5.3.9 Expressions

Expressions are made up of numbers, symbols, strings and operators. Operators can be unary (one operand, e.g. not) or binary (two operands, e.g. +). The operators allowable in expressions are listed in Table 5 - 3. The usual rules governing the syntax of expressions apply.

The operators listed may all be freely combined in both constant and relocatable expressions. The HI-TECH linker permits relocation of complex expressions, so the results of expressions involving relocatable identifiers may not be resolved until link time.

5.3.10 Statement Format

Legal statement formats are shown in Table 5 - 4 on page 128. The second form is only legal with certain directives, such as MACRO, SET and EQU. The label field is optional and if present should contain one identifier. The name field is mandatory and should also contain one identifier. There is no limitation on what column or part of the line any part of the statement should appear in.

5.3.11 Program Sections

Program sections, or psects, are a way of grouping together parts of a program even though the source code may not be physically adjacent in the source file, or even where spread over several source files. Unless defined as abs (absolute), psects are relocatable.

A psect is identified by a name and has several attributes. The psect directive is used to define psects. It takes as arguments a name and an optional comma-separated list of flags. See the section PSECT on page 131 for full information. The assembler associates no significance to the name of a psect.

The following is an example showing some executable instructions being placed in the text0 psect, and some data being placed in the rbss psect.

```
processor 18C452

PSECT text0,class=CODE,local,delta=1
adjust:
call clear_fred
movf flag
btfss 3,2
```
Table 5 - 3 Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>*</td>
<td>Multiplication</td>
</tr>
<tr>
<td>+</td>
<td>Addition</td>
</tr>
<tr>
<td>-</td>
<td>Subtraction</td>
</tr>
<tr>
<td>/</td>
<td>Division</td>
</tr>
<tr>
<td>= or eq</td>
<td>Equality</td>
</tr>
<tr>
<td>&gt; or gt</td>
<td>Signed greater than</td>
</tr>
<tr>
<td>&gt;= or ge</td>
<td>Signed greater than or equal to</td>
</tr>
<tr>
<td>&lt; or lt</td>
<td>Signed less than</td>
</tr>
<tr>
<td>&lt;= or le</td>
<td>Signed less than or equal to</td>
</tr>
<tr>
<td>&lt;&gt; or ne</td>
<td>Signed not equal to</td>
</tr>
<tr>
<td>low</td>
<td>Low byte of operand</td>
</tr>
<tr>
<td>high</td>
<td>High byte of operand</td>
</tr>
<tr>
<td>highword</td>
<td>High 16 bits of operand</td>
</tr>
<tr>
<td>mod</td>
<td>Modulus</td>
</tr>
<tr>
<td>&amp;</td>
<td>Bitwise AND</td>
</tr>
<tr>
<td>^</td>
<td>Bitwise XOR (exclusive or)</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td>not</td>
<td>Bitwise complement</td>
</tr>
<tr>
<td>&lt;&lt; or shl</td>
<td>Shift left</td>
</tr>
<tr>
<td>&gt;&gt; or shr</td>
<td>Shift right</td>
</tr>
<tr>
<td>rol</td>
<td>Rotate left</td>
</tr>
<tr>
<td>ror</td>
<td>Rotate right</td>
</tr>
<tr>
<td>seg</td>
<td>Segment (bank number) of address</td>
</tr>
<tr>
<td>float24</td>
<td>24-bit version of real operand</td>
</tr>
<tr>
<td>nul</td>
<td>Tests if macro argument is null</td>
</tr>
</tbody>
</table>

Table 5 - 4 ASPIC18 Statement formats

<table>
<thead>
<tr>
<th>Format 1:</th>
<th>label:</th>
<th>opcode</th>
<th>operands</th>
<th>; comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>Format 2:</td>
<td>name</td>
<td>pseudo-op</td>
<td>operands</td>
<td>; comment</td>
</tr>
<tr>
<td>Format 3:</td>
<td>; comment only</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

| goto    | next   |
```assembly
incf fred
goto clear_fred
next: decf fred

PSECT rbss, class=RAM, space=1
flag:
   ds 1
fred:
   ds 1

PSECT text, class=CODE, local, reloc=2
clear_fred:
   clrf fred
   bcf status, 5
   return
```

Note that even though the two blocks of code in the text psect are separated by a block in the rbss psect, the two text psect blocks will be contiguous when loaded by the linker. In other words, the decf fred instruction will fall through to the label clear_fred: during execution. The actual location in memory of the two psects will be determined by the linker. See the linker manual for information on how psect addresses are determined.

A label defined in a psect is said to be relocatable, that is, its actual memory address is not determined at assembly time. Note that this does not apply if the label is in the default (unnamed) psect, or in a psect declared absolute (see the PSECT pseudo-op description below). Any labels declared in an absolute psect will be absolute, that is their address will be determined by the assembler.

Relocatable expressions may be combined freely in expressions.

### 5.3.12 Assembler Directives

Assembler directives, or pseudo-ops, are used in a similar way to opcodes, but either do not generate code, or generate non-executable code, i.e. data bytes. The directives are listed in Table 5 - 5 on page 130, and are detailed below.

#### 5.3.12.1 GLOBAL

GLOBAL declares a list of symbols which, if defined within the current module, are made public. If the symbols are not defined in the current module, it is a reference to symbols in external modules. Example:

```assembly
GLOBAL lab1, lab2, lab3
```
Table 5 - 5 ASPIC18 Directives (pseudo-ops)

<table>
<thead>
<tr>
<th>Directive</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>GLOBAL</td>
<td>Make symbols accessible to other modules or allow reference to other modules’ symbols</td>
</tr>
<tr>
<td>END</td>
<td>End assembly</td>
</tr>
<tr>
<td>PSECT</td>
<td>Declare or resume program section</td>
</tr>
<tr>
<td>ORG</td>
<td>Set location counter</td>
</tr>
<tr>
<td>EQU</td>
<td>Define symbol value</td>
</tr>
<tr>
<td>SET</td>
<td>Define or re-define symbol value</td>
</tr>
<tr>
<td>DB</td>
<td>Define constant byte(s)</td>
</tr>
<tr>
<td>DW</td>
<td>Define constant word(s)</td>
</tr>
<tr>
<td>DS</td>
<td>Reserve storage</td>
</tr>
<tr>
<td>IF</td>
<td>Conditional assembly</td>
</tr>
<tr>
<td>ELSIF</td>
<td>Alternate conditional assembly</td>
</tr>
<tr>
<td>ELSE</td>
<td>Alternate conditional assembly</td>
</tr>
<tr>
<td>ENDIF</td>
<td>End conditional assembly</td>
</tr>
<tr>
<td>FNADDR</td>
<td>Inform linker that a function may be indirectly called</td>
</tr>
<tr>
<td>FNARG</td>
<td>Inform linker that evaluation of arguments for one function requires calling another</td>
</tr>
<tr>
<td>FNBREAK</td>
<td>Break call graph links</td>
</tr>
<tr>
<td>FCALL</td>
<td>Inform linker that one function calls another</td>
</tr>
<tr>
<td>FNCONF</td>
<td>Supply call graph configuration info for linker</td>
</tr>
<tr>
<td>FNINDIR</td>
<td>Inform linker that all functions with a particular signature may be indirectly called</td>
</tr>
<tr>
<td>FNROOT</td>
<td>Inform linker that a function is the “root” of a call graph</td>
</tr>
<tr>
<td>FNSIZE</td>
<td>Inform linker of argument and local variable sizes for a function</td>
</tr>
<tr>
<td>MACRO</td>
<td>Macro definition</td>
</tr>
<tr>
<td>ENDM</td>
<td>End macro definition</td>
</tr>
<tr>
<td>LOCAL</td>
<td>Define local tabs</td>
</tr>
<tr>
<td>ALIGN</td>
<td>Align output to the specified boundary</td>
</tr>
<tr>
<td>PAGESEL</td>
<td>Generate set/reset instruction to set PCLATH for this page</td>
</tr>
<tr>
<td>PROCESSOR</td>
<td>Define the particular chip for which this file is to be assembled</td>
</tr>
<tr>
<td>REPT</td>
<td>Repeat a block of code n times</td>
</tr>
<tr>
<td>IRP</td>
<td>Repeat a block of code with a list</td>
</tr>
<tr>
<td>IRPC</td>
<td>Repeat a block of code with a character list</td>
</tr>
<tr>
<td>SIGNAT</td>
<td>Define function signature</td>
</tr>
</tbody>
</table>
5.3.12.2 END

END is optional, but if present should be at the very end of the program. It will terminate the assembly. If an expression is supplied as an argument, that expression will be used to define the start address of the program. Whether this is of any use will depend on the linker. Example:

    END  start_label

5.3.12.3 PSECT

The PSECT directive declares or resumes a program section. It takes as arguments a name and optionally a comma-separated list of flags. The allowed flags are listed in Table 5-6 below. Once a psect has been declared it may be resumed later by simply giving its name as an argument to another psect directive; the flags need not be repeated.

Table 5-6 PSECT flags

<table>
<thead>
<tr>
<th>Flag</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>abs</td>
<td>Psect is absolute</td>
</tr>
<tr>
<td>bit</td>
<td>Psect holds bit objects</td>
</tr>
<tr>
<td>class</td>
<td>Specify class name for psect</td>
</tr>
<tr>
<td>delta</td>
<td>Size of an addressing unit</td>
</tr>
<tr>
<td>global</td>
<td>Psect is global (default)</td>
</tr>
<tr>
<td>limit</td>
<td>Upper address limit of psect</td>
</tr>
<tr>
<td>local</td>
<td>Psect is not global</td>
</tr>
<tr>
<td>ovrl</td>
<td>Psect will overlap same psect in other modules</td>
</tr>
<tr>
<td>pure</td>
<td>Psect is to be read-only</td>
</tr>
<tr>
<td>reloc</td>
<td>Start psect on specified boundary</td>
</tr>
<tr>
<td>size</td>
<td>Maximum size of psect</td>
</tr>
<tr>
<td>space</td>
<td>Represents area in which psect will reside</td>
</tr>
<tr>
<td>with</td>
<td>Place psect in the same page as specified psect</td>
</tr>
</tbody>
</table>

- abs defines the current psect as being absolute, i.e. it is to start at location 0. This does not mean that this module’s contribution to the psect will start at 0, since other modules may contribute to the same psect.

- The bit flag specifies that a psect hold objects that are 1 bit long. Such psects have a scale value of 8 to indicate that there are 8 addressable units to each byte of storage.

- The class flag specifies a class name for this psect. Class names are used to allow local psects to be referred to by a class name at link time, since they cannot be referred to by their own name. Class names are also useful where the linker address range feature is to be used.
The **delta** flag defines the size of an addressing unit. In other words, the number of bytes covered for an increment in the address. Since the PIC18 devices have both byte-wide ROM and RAM memory spaces, the **delta** value associated with PIC18 psects is usually one (which is the default **delta** value).

A psect defined as **global** will be combined with other **global** psects of the same name from other modules at link time. **global** is the default.

The **limit** flag specifies a limit on the highest address to which a psect may extend.

A psect defined as **local** will not be combined with other **local** psects at link time, even if there are others with the same name. Where there are two **local** psect directives in the one module, they reference the same psect. A **local** psect may not have the same name as any **global** psect, even one in another module.

A psect defined as **ovrld** will have the contribution from each module overlaid, rather than concatenated at run time. **ovrld** in combination with **abs** defines a truly absolute psect, i.e. a psect within which any symbols defined are absolute.

The **pure** flag instructs the linker that this psect will not be modified at run time and may therefore, for example, be placed in ROM. This flag is of limited usefulness since it depends on the linker and target system enforcing it.

The **reloc** flag allows specification of a requirement for alignment of the psect on a particular boundary, e.g. **reloc=100h** would specify that this psect must start on an address that is a multiple of 100h. The PIC18 psects which hold code use a **reloc=2** flag to align the instructions with the 16-bit word boundary.

The **size** flag allows a maximum size to be specified for the psect, e.g. **size=100h**. This will be checked by the linker after psects have been combined from all modules.

The **space** flag is used to differentiate areas of memory which have overlapping addresses, but which are distinct. Psects which are positioned in ROM and RAM have a different **space** value to indicate that ROM address zero, for example, is a different location to RAM address zero. Objects in different RAM banks have the same **space** value as their full addresses (including bank information) are unique.

The **with** flag allows a psect to be placed in the same page with a specified psect. For example **with=text0** will specify that this psect should be placed in the same page as the **text0** psect.

Some examples of the use of the **PSECT** directive follow:

```
PSECT fred
PSECT bill,size=100h,global
PSECT joh,abs,ovrld,class=CODE,delta=2
```
5.3.12.4 ORG

ORG changes the value of the location counter within the current psect. This means that the addresses set with ORG are relative to the base of the psect, which is not determined until link time.

The argument to ORG must be either an absolute value, or a value referencing the current psect. In either case the current location counter is set to the value of the argument. For example:

```
ORG 100h
```

will move the location counter to the beginning of the current psect plus 100h. The actual location will not be known until link time. It is possible to move the location counter backwards.

In order to use the ORG directive to set the location counter to an absolute value, an absolute, overlaid psect must be used:

```
PSECT absdata,abs,ovrld
ORG addr
```

where addr is an absolute address.

5.3.12.5 EQU

This pseudo-op defines a symbol and equates its value to an expression. For example

```
thomas EQU 123h
```

The identifier thomas will be given the value 123h. EQU is legal only when the symbol has not previously been defined. See also SET on page 133

5.3.12.6 SET

This pseudo-op is equivalent to EQU except that allows a symbol to be re-defined. For example

```
thomas SET 0h
```

5.3.12.7 DB

DB is used to initialize storage as bytes. The argument is a list of expressions, each of which will be assembled into one byte. Each character of the string will be assembled into one memory location.

An error will occur if the value of an expression is too big to fit into the memory location, e.g. if the value 1020 is given as an argument to DB.

Examples:

```
 alabel DB 'X',1,2,3,4,
```
Note that because the size of an address unit in ROM is 2 bytes, the DB pseudo-op will initialise a word with the upper byte set to zero.

### 5.3.12.8 DW

**DW** operates in a similar fashion to **DB**, except that it assembles expressions into words. An error will occur if the value of an expression is too big to fit into a word.

Example:

```
DW -1, 3664h, 'A', 3777Q
```

### 5.3.12.9 DS

This directive reserves, but does not initialize, memory locations. The single argument is the number of bytes to be reserved. Examples:

```
alabel:DS 23 ;Reserve 23 bytes of memory
xlabel:DS 2+3 ;Reserve 5 bytes of memory
```

### 5.3.12.10 FNADDR

This directive tells the linker that a function has its address taken, and thus could be called indirectly through a function pointer. Such would be the case if a function pointer was assigned the address of a function. For example, if the function `func()` had its address assigned to a pointer, the following would be produced in the assembler code:

```
FNADDR _func1
```

which tells the linker that `func1()` has its address taken.

### 5.3.12.11 FNARG

The directive

```
FNARG fun1,fun2
```

tells the linker that evaluation of the arguments to function `fun1()` involves a call to `fun2()`, thus the memory argument memory allocated for the two functions should not overlap. For example, the C function calls

```
fred(var1, bill(), 2);
```

will generate the assembler directive

```
FNARG _fred,_bill
```

thereby telling the linker that `bill()` is called while evaluating the arguments for a call to `fred()`.
5.3.12.12 **FNBREAK**

This directive is used to break links in the call graph information. The form of this directive is as follows:

\[
\text{FNBREAK } \text{fun1,fun2}
\]

and is automatically generated when the `interrupt_level` pragma is used. It states that any calls to `fun1` in trees other than the one rooted at `fun2` should not be considered when checking for functions that appear in multiple call graphs. `Fun2()` is typically `intlevel0` or `intlevel1` in compiler-generated code when the `interrupt_level` pragma is used. Memory for the `auto`/parameter area for a `fun1` will only be assigned in the tree rooted at `fun2`.

5.3.12.13 **FNCALL**

This directive takes the form:

\[
\text{FNCALL } \text{fun1,fun2}
\]

`FNCALL` is usually used in compiler generated code. It tells the linker that function `fun1()` calls function `fun2()`. This information is used by the linker when performing call graph analysis. If you write assembler code which calls a C function, use the `FNCALL` directive to ensure that your assembler function is taken into account. For example, if you have an assembler routine called `_fred` which calls a C routine called `foo()`, in your assembler code you should write:

\[
\text{FNCALL } _\text{fred},_\text{foo}
\]

5.3.12.14 **FNCONF**

The `FNCONF` directive is used to supply the linker with configuration information for a call graph. `FNCONF` is written as follows:

\[
\text{FNCONF } \text{psect,auto,args}
\]

where `psect` is the psect containing the call graph, `auto` is the prefix on all `auto` variable symbol names and `args` is the prefix on all function argument symbol names. This directive normally appears in only one place: the runtime startoff code used by compiler-generated code. For the HI-TECH PICC-18 Compiler the `picrt18x.as` module should include the directive:

\[
\text{FNCONF } \text{bss},?,?
\]

telling the linker that the call graph is in the `rbss` psect, auto variable blocks start with `?a` and function argument blocks start with `?`.

5.3.12.15 **FNINDIR**

This directive tells the linker that a function performs an indirect call to another function with a particular signature (see the `SIGNAT` directive). The linker must assume worst case that the function could call any other function which has the same signature and has had its address taken (see the `FNADDR`
directive). For example, if a function called fred() performs an indirect call to a function with signature 8249, the compiler will produce the directive:

```
FNINDIR _fred,8249
```

5.3.12.16 FNSIZE

The FNSIZE directive informs the linker of the size of the local variable and argument area associated with a function. These values are used by the linker when building the call graph and assigning addresses to the variable and argument areas. This directive takes the form:

```
FNSIZE  func,local,args
```

The named function has a local variable area and argument area as specified, for example

```
FNSIZE  _fred,10,5
```

means the function fred() has 10 bytes of local variables and 5 bytes of arguments. The function name arguments to any of the call graph associated directives may be local or global. Local functions are of course defined in the current module, but most be used in the call graph construction in the same manner as global names.

5.3.12.17 FNROOT

This directive tells the assembler that a function is a “root function” and thus forms the root of a call graph. It could either be the C main() function or an interrupt function. For example, the C main module produces the directive:

```
FNROOT  _main
```

5.3.12.18 IF, ELSIF, ELSE and ENDIF

These directives implement conditional assembly. The argument to IF and ELSIF should be an absolute expression. If it is non-zero, then the code following it up to the next matching ELSE will be assembled. If the expression is zero then the code up to the next matching ELSE will be skipped.

At an ELSE the sense of the conditional compilation will be inverted, while an ENDIF will terminate the conditional assembly block. Example:

```
IF      ABC
   call  aardvark
ELSIF   DEF
   call  denver
ELSE
   call  grapes
ENDC
```
In this example, if ABC is non-zero, the first call instruction will be assembled but not the second or third. If ABS is zero and DEF is non-zero, the second call will be assembled but the first and third will not. If both ABS and DEF are zero, the third call will be assembled. Conditional assembly blocks may be nested.

### 5.3.12.19 MACRO and ENDM

These directives provide for the definition of macros. The MACRO directive should be preceded by the macro name and optionally followed by a comma-separated list of formal parameters. When the macro is used, the macro name should be used in the same manner as a machine opcode, followed by a list of arguments to be substituted for the formal parameters.

For example:

```assembly
;macro: swap
;args: arg1, arg2 - the NUMBERS of the variables to swap
; arg3 - the NAME of the variable to use for temp storage;
; descr: Swaps two specified variables, where the variables
; are named:
; var_x
; and x is a number.
; uses: Uses the w register.

swap MACRO arg1, arg2, arg3
    movf var_&arg1, w
    movwf arg3
    movf var_&arg2, w
    movwf var_&arg1
    movf arg3, w
    movwf var_&arg2
ENDM
```

When used, this macro will expand to the 3 instructions in the body of the macro, with the formal parameters substituted by the arguments. Thus:

```assembly
swap 2,4,tempvar
```

expands to:

```assembly
movf var_2, w
movwf tempvar
movf var_4, w
```
PICC-18 Macro Assembler

```
movwf   var_2
movf    tempvar,w
movwf   var_4
```

A point to note in the above example: the “&” character is used to permit the concatenation of macro parameters with other text, but is removed in the actual expansion. The `nul` operator may be used within a macro to test a macro argument, for example:

```
if     nul     arg3     ; argument was not supplied.
...    
else    ; argument was supplied
... 
endif
```

A comment may be suppressed within the expansion of a macro (thus saving space in the macro storage) by opening the comment with a double `semicolon `;;`.```

5.3.12.20 LOCAL

The `LOCAL` directive allows unique labels to be defined for each expansion of a given macro. Any symbols listed after the `LOCAL` directive will have a unique assembler generated symbol substituted for them when the macro is expanded. For example:

```
down   MACRO   count
      LOCAL   more
      movlw   count
      movwf   tempvar
      more:   decfsz  tempvar
              goto    more
      ENDM
```

when expanded will include a unique assembler generated label in place of `more`. For example:

```
down   4
```

expands to:

```
      movlw   4
      movwf   tempvar
??0001   decfsz  tempvar
      goto    ??0001
```
if invoked a second time, the label more would expand to ??0002.

5.3.12.21 ALIGN

The ALIGN directive aligns whatever is following, data storage or code etc., to the specified boundary in the psect in which the directive is found. The boundary is specified by a number following the directive and is a number of bytes. For example, to align output to a 2 byte (even) address within a psect, the following could be used.

ALIGN 2

Note, however, that what follows will only begin on an even absolute address if the psect begins on an even address. The ALIGN directive can also be used to ensure that a psect’s length is a multiple of a certain number. For example, if the above ALIGN directive was placed at the end of a psect, the psect would have a length that was always an even number of bytes long.

5.3.12.22 REPT

The REPT directive temporarily defines an unnamed macro then expands it a number of times as determined by its argument. For example:

```
REPT 3
  addwf fred, fred
  andwf fred, w
ENDM
```

will expand to

```
addwf fred, fred
andwf fred, w
addwf fred, fred
andwf fred, w
addwf fred, fred
andwf fred, w
```

5.3.12.23 IRP and IRPC

The IRP and IRPC directives operate similarly to REPT. However, instead of repeating the block a fixed number of times, it is repeated once for each member of an argument list. In the case of IRP the list is a conventional macro argument list, in the case or IRPC it is each character in one argument. For each repetition the argument is substituted for one formal parameter.

For example:
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PSECT idata_0
IRP number,4865h,6C6Ch,6F00h
DW number
ENDM

PSECT text0

would expand to:

PSECT idata_0

DW 4865h
DW 6C6Ch
DW 6F00h

PSECT text0

Note that you can use local labels and angle brackets in the same manner as with conventional macros.

The IRPC directive is similar, except it substitutes one character at a time from a string of non-space characters.

For example:

PSECT idata_0
IRPC char,ABC
DW ‘char’
ENDM

PSECT text0

will expand to:

PSECT idata_0

DW ‘A’
DW ‘B’
DW ‘C’

PSECT text0
5.3.12.24 PAGESEL
It's sometimes necessary to set the current PCLATH bits so that a modify-PC type instruction will jump to a location in the current page of ROM, e.g.

```
    PAGESEL $
```

5.3.12.25 PROCESSOR
The output of the assembler depends on which chip it is desired to assemble for. This can be set on the command line, or with this directive, e.g.

```
    PROCESSOR 18C452
```

5.3.12.26 SIGNAT
This directive is used to associate a 16-bit signature value with a label. At link time the linker checks that all signatures defined for a particular label are the same and produces an error if they are not. The `SIGNAT` directive is used by the HI-TECH C compiler to enforce link time checking of function prototypes and calling conventions.

Use the `SIGNAT` directive if you want to write assembly language routines which are called from C. For example:

```
    SIGNAT _fred,8192
```

will associate the signature value 8192 with symbol `_fred`. If a different signature value for `_fred` is present in any object file, the linker will report an error.

5.3.13 Macro Invocations
When invoking a macro, the argument list must be comma-separated. If it is desired to include a comma (or other delimiter such as a `space`) in an argument then `angle brackets` `<` and `>` may be used to quote the argument. In addition the `exclamation mark` `!` may be used to quote a single character. The character immediately following the `exclamation mark` will be passed into the macro argument even if it is normally a comment indicator.

If an argument is preceded by a percent sign `%`, that argument will be evaluated as an expression and passed as a decimal number, rather than as a string. This is useful if evaluation of the argument inside the macro body would yield a different result.

5.3.14 Assembler Controls
Assembler controls may be included in the assembler source to control such things as listing format. These keywords have no significance anywhere else in the program. The control is invoked by the directive `OPT` followed by the control name. Some keywords are followed by one or more parameters. For example:
A list of keywords is given in Table 5 - 7, and each is described further below.

### Table 5 - 7 ASPIC18 Assembler controls

<table>
<thead>
<tr>
<th>Control</th>
<th>Meaning</th>
<th>Format</th>
</tr>
</thead>
<tbody>
<tr>
<td>COND*</td>
<td>Include conditional code in the listing</td>
<td>COND</td>
</tr>
<tr>
<td>EXPAND</td>
<td>Expand macros in the listing output</td>
<td>EXPAND</td>
</tr>
<tr>
<td>INCLUDE</td>
<td>Textually include another source file</td>
<td>INCLUDE &lt;pathname&gt;</td>
</tr>
<tr>
<td>LIST*</td>
<td>Define options for listing output</td>
<td>LIST [&lt;listopt&gt;, ..., &lt;listopt&gt;]</td>
</tr>
<tr>
<td>NOCOND</td>
<td>Leave conditional code out of the listing</td>
<td>NOCOND</td>
</tr>
<tr>
<td>NOEXPAND*</td>
<td>Disable macro expansion</td>
<td>NOEXPAND</td>
</tr>
<tr>
<td>NOLIST</td>
<td>Disable listing output</td>
<td>NOLIST</td>
</tr>
<tr>
<td>PAGE</td>
<td>Start a new page in the listing output</td>
<td>PAGE</td>
</tr>
<tr>
<td>SUBTITLE</td>
<td>Specify the subtitle of the program</td>
<td>SUBTITLE &quot;&lt;subtitle&gt;&quot;</td>
</tr>
<tr>
<td>TITLE</td>
<td>Specify the title of the program</td>
<td>TITLE &quot;&lt;title&gt;&quot;</td>
</tr>
</tbody>
</table>

a. The default options are listed with an asterix (*).

#### 5.3.14.1 COND

Any conditional code will be included in the listing output. See also the NOCOND control.

#### 5.3.14.2 EXPAND

When EXPAND is in effect, the code generated by macro expansions will appear in the listing output. See also the NOEXPAND control.

#### 5.3.14.3 INCLUDE

This control causes the file specified by `pathname` to be textually included at that point in the assembly file. The INCLUDE control must be the last control keyword on the line.

#### 5.3.14.4 LIST

If the listing was previously turned off using the NOLIST control, the LIST control on its own will turn the listing on.

Alternatively, the LIST control may includes options to control the assembly and the listing. The options are listed in Table 5 - 8. See also the NOLIST control.
Table 5 - 8 LIST Control Options

<table>
<thead>
<tr>
<th>List Option</th>
<th>Default</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>c=nnn</td>
<td>80</td>
<td>Set the page (i.e. column) width.</td>
</tr>
<tr>
<td>n=nnn</td>
<td>59</td>
<td>Set the page length.</td>
</tr>
<tr>
<td>t=ON</td>
<td>OFF</td>
<td>OFF</td>
</tr>
<tr>
<td>p=processor&gt;</td>
<td>n/a</td>
<td>Set the processor type.</td>
</tr>
<tr>
<td>r=radix&gt;</td>
<td>hex</td>
<td>Set the default radix to hex, dec or oct.</td>
</tr>
<tr>
<td>x=ON</td>
<td>OFF</td>
<td>OFF</td>
</tr>
</tbody>
</table>

5.3.14.5 NOCOND

Any conditional code will not be included in the listing output. See also the COND control.

5.3.14.6 NOEXPAND

NOEXPAND disables macro expansion in the listing file. The macro call will be listed instead. See also the EXPAND control.

5.3.14.7 NOLIST

This control turns the listing output off from this point onwards. See also the LIST control.

5.3.14.8 NOXREF

NOXREF will disable generation of the raw cross reference file. See also the XREF control.

5.3.14.9 PAGE

PAGE causes a new page to be started in the listing output. A Control-L (form feed) character will also cause a new page when encountered in the source.

5.3.14.10 SPACE

The SPACE control will place the a number of blank lines in the listing output as specified by its parameter.

5.3.14.11 SUBTITLE

SUBTITLE defines a subtitle to appear at the top of every listing page, but under the title. The string should be enclosed in single or double quotes. See also the TITLE control.

5.3.14.12 TITLE

This control keyword defines a title to appear at the top of every listing page. The string should be enclosed in single or double quotes. See also the SUBTITLE control.
5.3.14.13 XREF

XREF is equivalent to the command line option -CR; it causes the assembler to produce a raw cross reference file. The utility CREF should be used to actually generate the formatted cross-reference listing.
6.1 Introduction

HI-TECH C incorporates a relocating assembler and linker to permit separate compilation of C source files. This means that a program may be divided into several source files, each of which may be kept to a manageable size for ease of editing and compilation, then each source file may be compiled separately and finally all the object files linked together into a single executable program.

This chapter describes the theory behind and the usage of the linker. Note however that in most instances it will not be necessary to use the linker directly, as the compiler drivers (HI-TIDE, HPD or command line) will automatically invoke the linker with all necessary arguments. Using the linker directly is not simple, and should be attempted only by those with a sound knowledge of the compiler and linking in general.

If it is absolutely necessary to use the linker directly, the best way to start is to copy the linker arguments constructed by the compiler driver, and modify them as appropriate. This will ensure that the necessary startup module and arguments are present.

Note also that the linker supplied with HI-TECH C is generic to a wide variety of compilers for several different processors. Not all features described in this chapter are applicable to all compilers.

6.2 Relocation and Psects

The fundamental task of the linker is to combine several relocatable object files into one. The object files are said to be relocatable since the files have sufficient information in them so that any references to program or data addresses (e.g. the address of a function) within the file may be adjusted according to where the file is ultimately located in memory after the linkage process. Thus the file is said to be relocatable. Relocation may take two basic forms; relocation by name, i.e. relocation by the ultimate value of a global symbol, or relocation by psect, i.e. relocation by the base address of a particular section of code, for example the section of code containing the actual executable instructions.

6.3 Program Sections

Any object file may contain bytes to be stored in memory in one or more program sections, which will be referred to as psects. These psects represent logical groupings of certain types of code bytes in the program. In general the compiler will produce code in three basic types of psects, although there will be several different types of each. The three basic kinds are text psects, containing
executable code, data psects, containing initialised data, and bss psects, containing uninitialised but reserved data.

The difference between the data and bss psects may be illustrated by considering two external variables; one is initialised to the value 1, and the other is not initialised. The first will be placed into the data psect, and the second in the bss psect. The bss psect is always cleared to zeros on startup of the program, thus the second variable will be initialised at run time to zero. The first will however occupy space in the program file, and will maintain its initialised value of 1 at startup. It is quite possible to modify the value of a variable in the data psect during execution, however it is better practice not to do so, since this leads to more consistent use of variables, and allows for restartable and romable programs.

For more information on the particular psects used in a specific compiler, refer to the appropriate machine-specific chapter.

6.4 Local Psects

Most psects are global, i.e. they are referred to by the same name in all modules, and any reference in any module to a global psect will refer to the same psect as any other reference. Some psects are local, which means that they are local to only one module, and will be considered as separate from any other psect even of the same name in another module. Local psects can only be referred to at link time by a class name, which is a name associated with one or more psects via the PSECT directive class= in assembler code. See The Macro Assembler chapter for more information on psect options.

6.5 Global Symbols

The linker handles only symbols which have been declared as GLOBAL to the assembler. The code generator generates these assembler directives whenever it encounters global C objects. At the C source level, this means all names which have storage class external and which are not declared as static. These symbols may be referred to by modules other than the one in which they are defined. It is the linker's job to match up the definition of a global symbol with the references to it. Other symbols (local symbols) are passed through the linker to the symbol file, but are not otherwise processed by the linker.

6.6 Link and load addresses

The linker deals with two kinds of addresses; link and load addresses. Generally speaking the link address of a psect is the address by which it will be accessed at run time. The load address, which may or may not be the same as the link address, is the address at which the psect will start within the output file (HEX or binary file etc.). In the case of the 8086 processor, the link address roughly corresponds to the offset within a segment, while the load address corresponds to the physical address of a segment. The segment address is the load address divided by 16.
Other examples of link and load addresses being different are; an initialised data psect that is copied from ROM to RAM at startup, so that it may be modified at run time; a banked text psect that is mapped from a physical (== load) address to a virtual (== link) address at run time.

The exact manner in which link and load addresses are used depends very much on the particular compiler and memory model being used.

### 6.7 Operation

A command to the linker takes the following form:

```
hlink options files ...
```

*Options* is zero or more linker options, each of which modifies the behaviour of the linker in some way. *Files* is one or more object files, and zero or more library names. The options recognised by the linker are listed in Table 6 - 1 on page 147 and discussed in the following paragraphs.

<table>
<thead>
<tr>
<th>Option</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>-A class=low-high,...</td>
<td>Specify address ranges for a class</td>
</tr>
<tr>
<td>-C x</td>
<td>Call graph options</td>
</tr>
<tr>
<td>-C psect=class</td>
<td>Specify a class name for a global psect</td>
</tr>
<tr>
<td>-C baseaddr</td>
<td>Produce binary output file based at <code>baseaddr</code></td>
</tr>
<tr>
<td>-D class=delta</td>
<td>Specify a class delta value</td>
</tr>
<tr>
<td>-D symfile</td>
<td>Produce old-style symbol file</td>
</tr>
<tr>
<td>-E errfile</td>
<td>Write error messages to <code>errfile</code></td>
</tr>
<tr>
<td>-F</td>
<td>Produce <code>.obj</code> file with only symbol records</td>
</tr>
<tr>
<td>-G spec</td>
<td>Specify calculation for segment selectors</td>
</tr>
<tr>
<td>-H symfile</td>
<td>Generate symbol file</td>
</tr>
<tr>
<td>-H+symfile</td>
<td>Generate enhanced symbol file</td>
</tr>
<tr>
<td>-I</td>
<td>Ignore undefined symbols</td>
</tr>
<tr>
<td>-J num</td>
<td>Set maximum number of errors before aborting</td>
</tr>
<tr>
<td>-K</td>
<td>Prevent overlaying function parameter and auto areas</td>
</tr>
<tr>
<td>-L</td>
<td>Preserve relocation items in <code>.obj</code> file</td>
</tr>
<tr>
<td>-LM</td>
<td>Preserve segment relocation items in <code>.obj</code> file</td>
</tr>
<tr>
<td>-N</td>
<td>Sort symbol table in map file by address order</td>
</tr>
</tbody>
</table>

1. In earlier versions of HI-TECH C the linker was called LINK.EXE
6.7.1 **Numbers in linker options**

Several linker options require memory addresses or sizes to be specified. The syntax for all these is similar. By default, the number will be interpreted as a decimal value. To force interpretation as a hex number, a trailing “H” should be added, e.g. 765FH will be treated as a hex number.

### 6.7.2 **-Aclass=low-high,**

Normally psects are linked according to the information given to a -P option (see below) but sometimes it is desired to have a class of psects linked into more than one non-contiguous address range. This option allows a number of address ranges to be specified for a class. For example:

```
-ACODE=1020h-7FFEh,8000h-BFFEh
```

specifies that the class CODE is to be linked into the given address ranges. Note that a contribution to a psect from one module cannot be split, but the linker will attempt to pack each block from each module into the address ranges, starting with the first specified.

Where there are a number of identical, contiguous address ranges, they may be specified with a repeat count, e.g.

```
-ACODE=0-FFFFhx16
```
specifies that there are 16 contiguous ranges each 64k bytes in size, starting from zero. Even though the ranges are contiguous, no code will straddle a 64k boundary. The repeat count is specified as the character “x” or “*” after a range, followed by a count.

6.7.3 -Cx

These options allow control over the call graph information which may be included in the map file produced by the linker. The -CN option removes the call graph information from the map file. The -CC option only include the critical paths of the call graph. A function call that is marked with a “*” in a full call graph is on a critical path and only these calls are included when the -CC option is used. A call graph is only produced for processors and memory models that use a compiled stack.

6.7.4 -Cpsect=class

This option will allow a psect to be associated with a specific class. Normally this is not required on the command line since classes are specified in object files.

6.7.5 -Dclass=delta

This option allows the delta value for psects that are members of the specified class to be defined. The delta value should be a number and represents the number of bytes per addressable unit of objects within the psects. Most psects do not need this option as they are defined with a delta value.

6.7.6 -Dsymfile

Use this option to produce an old-style symbol file. An old-style symbol file is an ASCII file, where each line has the link address of the symbol followed by the symbol name.

6.7.7 -Errfile

Error messages from the linker are written to standard error (file handle 2). Under DOS there is no convenient way to redirect this to a file (the compiler drivers will redirect standard error if standard output is redirected). This option will make the linker write all error messages to the specified file instead of the screen, which is the default standard error destination.

6.7.8 -F

Normally the linker will produce an object file that contains both program code and data bytes, and symbol information. Sometimes it is desired to produce a symbol-only object file that can be used again in a subsequent linker run to supply symbol values. The -F option will suppress data and code bytes from the output file, leaving only the symbol records.

This option can be used when producing more than one hex file for situations where the program is contained in different memory devices located at different addresses. The files for one device are
compiled using this linker option to produce a symbol-only object file; this is then linked with the files for the other device. The process can then be repeated for the other files and device.

### 6.7.9 -G spec

When linking programs using segmented, or bank-switched psects, there are two ways the linker can assign segment addresses, or selectors, to each segment. A segment is defined as a contiguous group of psects where each psect in sequence has both its link and load address concatenated with the previous psect in the group. The segment address or selector for the segment is the value derived when a segment type relocation is processed by the linker.

By default the segment selector will be generated by dividing the base load address of the segment by the relocation quantum of the segment, which is based on the `reloc=` directive value given to psects at the assembler level. This is appropriate for 8086 real mode code, but not for protected mode or some bank-switched arrangements. In this instance the `-G` option is used to specify a method for calculating the segment selector. The argument to `-G` is a string similar to:

```
A/10h-4h
```

where `A` represents the load address of the segment and `/` represents division. This means "Take the load address of the psect, divide by 10 hex, then subtract 4". This form can be modified by substituting `N` for `A`, `*` for `/` (to represent multiplication), and adding rather than subtracting a constant. The token `N` is replaced by the ordinal number of the segment, which is allocated by the linker. For example:

```
N*8+4
```

means "take the segment number, multiply by 8 then add 4". The result is the segment selector. This particular example would allocate segment selectors in the sequence 4, 12, 20, ... for the number of segments defined. This would be appropriate when compiling for 80286 protected mode, where these selectors would represent LDT entries.

### 6.7.10 -H symfile

This option will instruct the linker to generate a symbol file. The optional argument `symfile` specifies a file to receive the symbol file. The default file name is `l.sym`.

### 6.7.11 -H+symfile

This option will instruct the linker to generate an enhanced symbol file, which provides, in addition to the standard symbol file, class names associated with each symbol and a segments section which lists each class name and the range of memory it occupies. This format is recommended if the code is to be run in conjunction with a debugger. The optional argument `symfile` specifies a file to receive the symbol file. The default file name is `l.sym`. 
6.7.12 -Jerrcount

The linker will stop processing object files after a certain number of errors (other than warnings). The default number is 10, but the -J option allows this to be altered.

6.7.13 -K

For compilers that use a compiled stack, the linker will try and overlay function auto and parameter areas in an attempt to reduce the total amount of RAM required. For debugging purposes, this feature can be disabled with this option.

6.7.14 -I

Usually failure to resolve a reference to an undefined symbol is a fatal error. Use of this option will cause undefined symbols to be treated as warnings instead.

6.7.15 -L

When the linker produces an output file it does not usually preserve any relocation information, since the file is now absolute. In some circumstances a further "relocation" of the program will be done at load time, e.g. when running a .exe file under DOS or a .prg file under TOS. This requires that some information about what addresses require relocation is preserved in the object (and subsequently the executable) file. The -L option will generate in the output file one null relocation record for each relocation record in the input.

6.7.16 -LM

Similar to the above option, this preserves relocation records in the output file, but only segment relocations. This is used particularly for generating .exe files to run under DOS.

6.7.17 -M mapfile

This option causes the linker to generate a link map in the named file, or on the standard output if the file name is omitted. The format of the map file is illustrated in Section 6.9 on page 155.

6.7.18 -N, -Ns and -Nc

By default the symbol table in the link map will be sorted by name. The -N option will cause it to be sorted numerically, based on the value of the symbol. The -Ns and -Nc options work similarly except that the symbols are grouped by either their space value, or class.

6.7.19 -O outfile

This option allows specification of an output file name for the linker. The default output file name is l.obj. Use of this option will override the default.
6.7.20 \texttt{-Pspec}

Psects are linked together and assigned addresses based on information supplied to the linker via \texttt{-P} options. The argument to the \texttt{-P} option consists basically of \textit{comma}-separated sequences thus:

\begin{verbatim}
\texttt{-Ppsect=lnkaddr+min/ldaddr+min,psect=lnkaddr/ldaddr, ...}
\end{verbatim}

There are several variations, but essentially each psect is listed with its desired link and load addresses, and a minimum value. All values may be omitted, in which case a default will apply, depending on previous values.

The minimum value, \texttt{min}, is preceded by a \texttt{+} sign, if present. It sets a minimum value for the link or load address. The address will be calculated as described below, but if it is less than the minimum then it will be set equal to the minimum.

The link and load addresses are either numbers as described above, or the names of other psects or classes, or special tokens. If the link address is a negative number, the psect is linked in reverse order with the top of the psect appearing at the specified address minus one. Psects following a negative address will be placed before the first psect in memory. If a link address is omitted, the psect's link address will be derived from the top of the previous psect, e.g.

\begin{verbatim}
-P\texttt{text}=100h,\texttt{data},\texttt{bss}
\end{verbatim}

In this example the \texttt{text} psect is linked at 100 hex (its load address defaults to the same). The \texttt{data} psect will be linked (and loaded) at an address which is 100 hex plus the length of the \texttt{text} psect, rounded up as necessary if the \texttt{data} psect has a \texttt{reloc=} value associated with it. Similarly, the \texttt{bss} psect will concatenate with the \texttt{data} psect. Again:

\begin{verbatim}
-P\texttt{text}=-100h,\texttt{data},\texttt{bss}
\end{verbatim}

will link in assending order \texttt{bss},\texttt{data} then \texttt{text} with the top of \texttt{text} appearing at address \texttt{0ffh}.

If the load address is omitted entirely, it defaults to the same as the link address. If the \texttt{slash} “/” character is supplied, but no address is supplied after it, the load address will concatenate with the previous psect, e.g.

\begin{verbatim}
-P\texttt{text}=0,\texttt{data}=0/,\texttt{bss}
\end{verbatim}

will cause both \texttt{text} and \texttt{data} to have a link address of zero, \texttt{text} will have a load address of 0, and \texttt{data} will have a load address starting after the end of \texttt{text}. The \texttt{bss} psect will concatenate with \texttt{data} for both link and load addresses.

The load address may be replaced with a \texttt{dot} “.” character. This tells the linker to set the load address of this psect to the same as its link address. The link or load address may also be the name
of another (already linked) psect. This will explicitly concatenate the current psect with the previously specified psect, e.g.

```
-Ptext=0,data=8000h/,bss/. -Pnvram=bss,heap
```

This example shows text at zero, data linked at 8000h but loaded after text, bss is linked and loaded at 8000h plus the size of data, and nvram and heap are concatenated with bss. Note here the use of two -P options. Multiple -P options are processed in order.

If -A options have been used to specify address ranges for a class then this class name may be used in place of a link or load address, and space will be found in one of the address ranges. For example:

```
-ACODE=8000h-BFFEh,E000h-FFFEh
-Pdata=C000h/CODE
```

This will link data at C000h, but find space to load it in the address ranges associated with CODE. If no sufficiently large space is available, an error will result. Note that in this case the data psect will still be assembled into one contiguous block, whereas other psects in the class CODE will be distributed into the address ranges wherever they will fit. This means that if there are two or more psects in class CODE, they may be intermixed in the address ranges.

Any psects allocated by a -P option will have their load address range subtracted from any address ranges specified with the -A option. This allows a range to be specified with the -A option without knowing in advance how much of the lower part of the range, for example, will be required for other psects.

**6.7.21 -Qprocessor**

This option allows a processor type to be specified. This is purely for information placed in the map file. The argument to this option is a string describing the processor.

**6.7.22 -S**

This option prevents symbol information relating from being included in the symbol file produced by the linker. Segment information is still included.

**6.7.23 -Sclass=limit[, bound]**

A class of psects may have an upper address limit associated with it. The following example places a limit on the maximum address of the CODE class of psects to one less than 400h.

```
-SCODE=400h
```

Note that to set an upper limit to a psect, this must be set in assembler code (with a limit= flag on a PSECT directive).
If the *bound* (boundary) argument is used, the class of psects will start on a multiple of the bound address. This example places the **FARCODE** class of psects at a multiple of 1000h, but with an upper address limit of 6000h:

```
-SFARCODE=6000h,1000h
```

### 6.7.24 -Usymbol

This option will enter the specified symbol into the linker's symbol table as an undefined symbol. This is useful for linking entirely from libraries, or for linking a module from a library where the ordering has been arranged so that by default a later module will be linked.

### 6.7.25 -Vavmap

To produce an *Avocet* format symbol file, the linker needs to be given a map file to allow it to map psect names to *Avocet* memory identifiers. The avmap file will normally be supplied with the compiler, or created automatically by the compiler driver as required.

### 6.7.26 -Wnum

The `-W` option can be used to set the warning level, in the range -9 to 9, or the width of the map file, for values of `num >= 10`.

- `-W9` will suppress all warning messages. `-W0` is the default. Setting the warning level to -9 (`-W-9`) will give the most comprehensive warning messages.

### 6.7.27 -X

Local symbols can be suppressed from a symbol file with this option. Global symbols will always appear in the symbol file.

### 6.7.28 -Z

Some local symbols are compiler generated and not of interest in debugging. This option will suppress from the symbol file all local symbols that have the form of a single alphabetic character, followed by a digit string. The set of letters that can start a trivial symbol is currently "klfL$u". The `-Z` option will strip any local symbols starting with one of these letters, and followed by a digit string.

## 6.8 Invoking the Linker

The linker is called **HLINK**, and normally resides in the BIN subdirectory of the compiler installation directory. It may be invoked with no arguments, in which case it will prompt for input from standard input. If the standard input is a file, no prompts will be printed. This manner of invocation is generally useful if the number of arguments to **HLINK** is large. Even if the list of files is too long to
fit on one line, continuation lines may be included by leaving a backslash “\” at the end of the preceding line. In this fashion, HLINK commands of almost unlimited length may be issued. For example a link command file called x.lnk and containing the following text:

```
-Z -OX.OBJ -MX.MAP \ 
-Ptext=0, data=0/, bss, nvram=bss/ . \ 
X.OBJ Y.OBJ Z.OBJ C:\HT-Z80\LIB\Z80-SC.LIB
```

may be passed to the linker by one of the following:

```
  hlink @x.lnk
  hlink <x.lnk
```

### 6.9 Map Files

The map file contains information relating to the relocation of psects and the addresses assigned to symbols within those psects. The sections in the map file are as follows; first is a copy of the command line used to invoke the linker. This is followed by the version number of the object code in the first file linked, and the machine type. This is optionally followed by call graph information, depended on the processor and memory model selected. Then are listed all object files that were linked, along with their psect information. Libraries are listed, with each module within the library. The TOTALS section summarises the psects from the object files. The SEGMENTS section summarises major memory groupings. This will typically show RAM and ROM usage. The segment names are derived from the name of the first psect in the segment.

Lastly (not shown in the example) is a symbol table, where each global symbol is listed with its associated psect and link address.

```
Linker command line:

-z -Mmap -p vectors=00h, text, strings, const, im2vecs -pbaseram=00h \ 
-pramstart=08000h, data/im2vecs, bss/., stack=09000h - 
-pnvram=bss, heap \ 
-oC:\TEMP\l.obj C:\HT-Z80\LIB\rtz80-s.obj hello.obj \ 
C:\HT-Z80\LIB\z80-sc.lib
```

Object code version is 2.4
Machine type is Z80

<table>
<thead>
<tr>
<th>Name</th>
<th>Link</th>
<th>Load</th>
<th>Length</th>
<th>Selector</th>
</tr>
</thead>
<tbody>
<tr>
<td>C:\HT-Z80\LIB\rtz80-s.obj</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>vectors</td>
<td>0</td>
<td>0</td>
<td>71</td>
<td></td>
</tr>
</tbody>
</table>
### Call Graph Information

A call graph is produced for chip types and memory models that use a compiled stack, rather than a hardware stack, to facilitate parameter passing between functions and `auto` variables defined within a function. When a compiled stack is used, functions are not re-entrant since the function will use a fixed area of memory for its local objects (parameters/`auto` variables). A function called `foo()`, for example, will use symbols like `_foo` for parameters and `a_foo` for `auto` variables. Compilers such as the PIC, 6805 and V8 use compiled stacks. The 8051 compiler uses a compiled stack in small and medium memory models. The call graph shows information relating to the placement of function parameters and `auto` variables by the linker. A typical call graph may look something like:

Call graph:

```
* _main size 0,0 offset 0
  _init size 2,3 offset 0
    _ports size 2,2 offset 5
* _sprintf size 5,10 offset 0
```
The graph shows the functions called and the memory usage (RAM) of the functions for their own local objects. In the example above, the symbol \_main is associated with the function main(). It is shown at the far left of the call graph. This indicates that it is the root of a call tree. The run-time code has the FNROOT assembler directive that specifies this. The size field after the name indicates the number of parameters and auto variables, respectively. Here, main() takes no parameters and defines no auto variables. The offset field is the offset at which the function’s parameters and auto variables have been placed from the beginning of the area of memory used for this purpose. The runtime code contains a FNCONF directive which tells the compiler in which psect parameters and auto variables should reside. This memory will be shown in the map file under the name COMMON.

Main() calls a function called init(). This function uses a total of two bytes of parameters (it may be two objects of type char or one int; that is not important) and has three bytes of auto variables. These figures are the total of bytes of memory consumed by the function. If the function was passed a two-byte int, but that was done via a register, then the two bytes would not be included in this total. Since main() did not use any of the local object memory, the offset of init()’s memory is still at 0.

The function init() itself calls another function called ports(). This function uses two bytes of parameters and another two bytes of auto variables. Since ports() is called by init(), its local variables cannot be overlapped with those of init()’, so the offset is 5, which means that ports()’s local objects were placed immediately after those of init()’s.

The function main also calls sprintf(). Since the function sprintf is not active at the same time as init() or ports(), their local objects can be overlapped and the offset is hence set to 0. Sprintf() calls a function putch(), but this function uses no memory for parameters (the char passed as argument is apparently done so via a register) or locals, so the size and offset are zero and are not printed.

Main() also calls another function indirectly using a function pointer. This is indicated by the two INDIRECT entries in the graph. The number following is the signature value of functions that could potentially be called by the indirect call. This number is calculated from the parameters and return type of the functions the pointer can indirectly call. The names of any functions that have this signature value are listed underneath the INDIRECT entries. Their inclusion does not mean that they were called (there is no way to determine that), but that they could potentially be called.
The last line shows another function whose name is at the far left of the call graph. This implies that this is the root of another call graph tree. This is an interrupt function which is not called by any code, but which is automatically invoked when an enabled interrupt occurs. This interrupt routine calls the function incr(), which is shown shorthand in the graph by the “->” symbol followed by the called function’s name instead of having that function shown indented on the following line. This is done whenever the calling function does not takes parameters, nor defines any variables.

Those lines in the graph which are starred “*” are those functions which are on a critical path in terms of RAM usage. For example, in the above, (main() is a trivial example) consider the function sprintf(). This uses a large amount of local memory and if you could somehow rewrite it so that it used less local memory, it would reduce the entire program’s RAM usage. The functions init() and ports() have had their local memory overlapped with that of sprintf(), so reducing the size of these functions’ local memory will have no affect on the program’s RAM usage. Their memory usage could be increased, as long as the total size of the memory used by these two functions did not exceed that of sprintf(), with no additional memory used by the program. So if you have to reduce the amount of RAM used by the program, look at those functions that are starred.

If, when searching a call graph, you notice that a function’s parameter and auto areas have been overlapped (i.e. ?a_foo was placed at the same address as ?_foo, for example), then check to make sure that you have actually called the function in your program. If the linker has not seen a function actually called, then it overlaps these areas of memory since that are not needed. This is a consequence of the linker’s ability to overlap the local memory areas of functions which are not active at the same time. Once the function is called, unique addresses will be assigned to both the parameters and auto objects.

If you are writing a routine that calls C code from assembler, you will need to include the appropriate assembler directives to ensure that the linker sees the C function being called.

### 6.10 Librarian

The librarian program, LIBR, has the function of combining several object files into a single file known as a library. The purposes of combining several such object modules are several.

- fewer files to link
- faster access
- uses less disk space

In order to make the library concept useful, it is necessary for the linker to treat modules in a library differently from object files. If an object file is specified to the linker, it will be linked into the final linked module. A module in a library, however, will only be linked in if it defines one or more symbols previously known, but not defined, to the linker. Thus modules in a library will be linked only if required. Since the choice of modules to link is made on the first pass of the linker, and the
library is searched in a linear fashion, it is possible to order the modules in a library to produce special effects when linking. More will be said about this later.

### 6.10.1 The Library Format

The modules in a library are basically just concatenated, but at the beginning of a library is maintained a directory of the modules and symbols in the library. Since this directory is smaller than the sum of the modules, the linker is speeded up when searching a library since it need read only the directory and not all the modules on the first pass. On the second pass it need read only those modules which are required, seeking over the others. This all minimises disk I/O when linking.

It should be noted that the library format is geared exclusively toward object modules, and is not a general purpose archiving mechanism as is used by some other compiler systems. This has the advantage that the format may be optimized toward speeding up the linkage process.

### 6.10.2 Using the Librarian

The librarian program is called **LIBR**, and the format of commands to it is as follows:

```plaintext
libr options k file.lib file.obj ...
```

Interpreting this, **LIBR** is the name of the program, **options** is zero or more librarian options which affect the output of the program. **k** is a key letter denoting the function requested of the librarian (replacing, extracting or deleting modules, listing modules or symbols), **file.lib** is the name of the library file to be operated on, and **file.obj** is zero or more object file names.

The librarian options are listed in Table 6 - 2.

#### Table 6 - 2 Librarian Options

<table>
<thead>
<tr>
<th>Option</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>-Pwidth</td>
<td>specify page width</td>
</tr>
<tr>
<td>-W</td>
<td>suppress non-fatal errors</td>
</tr>
</tbody>
</table>

The key letters are listed in Table 6 - 3.

When replacing or extracting modules, the **file.obj** arguments are the names of the modules to be replaced or extracted. If no such arguments are supplied, all the modules in the library will be replaced or extracted respectively. Adding a file to a library is performed by requesting the librarian to replace it in the library. Since it is not present, the module will be appended to the library. If the **r** key is used and the library does not exist, it will be created.

Under the **d** key letter, the named object files will be deleted from the library. In this instance, it is an error not to give any object file names.
The \texttt{m} and \texttt{s} key letters will list the named modules and, in the case of the \texttt{s} keyletter, the symbols defined or referenced within (global symbols only are handled by the librarian). As with the \texttt{r} and \texttt{x} key letters, an empty list of modules means all the modules in the library.

### 6.10.3 Examples

Here are some examples of usage of the librarian. The following lists the global symbols in the modules \texttt{a.obj}, \texttt{b.obj} and \texttt{c.obj}:

\begin{verbatim}
  libr s file.lib a.obj b.obj c.obj
\end{verbatim}

This command deletes the object modules \texttt{a.obj}, \texttt{b.obj} and \texttt{2.obj} from the library \texttt{file.lib}:

\begin{verbatim}
  libr d file.lib a.obj b.obj 2.obj
\end{verbatim}

### 6.10.4 Supplying Arguments

Since it is often necessary to supply many object file arguments to LIBR, and command lines are restricted to 127 characters by CP/M and MS-DOS, LIBR will accept commands from standard input if no command line arguments are given. If the standard input is attached to the console, LIBR will prompt for input. Multiple line input may be given by using a \texttt{backslash} as a continuation character on the end of a line. If standard input is redirected from a file, LIBR will take input from the file, without prompting. For example:

\begin{verbatim}
  libr
  libr> r file.lib 1.obj 2.obj 3.obj \ 
  libr> 4.obj 5.obj 6.obj
\end{verbatim}

will perform much the same as if the object files had been typed on the command line. The libr> prompts were printed by LIBR itself, the remainder of the text was typed as input.

\begin{verbatim}
  libr <lib.cmd
\end{verbatim}

LIBR will read input from \texttt{lib.cmd}, and execute the command found therein. This allows a virtually unlimited length command to be given to LIBR.
6.10.5 Listing Format

A request to LIBR to list module names will simply produce a list of names, one per line, on standard output. The s keyletter will produce the same, with a list of symbols after each module name. Each symbol will be preceded by the letter D or U, representing a definition or reference to the symbol respectively. The -P option may be used to determine the width of the paper for this operation. For example:

```
LIBR -P80 s file.lib
```

will list all modules in file.lib with their global symbols, with the output formatted for an 80 column printer or display.

6.10.6 Ordering of Libraries

The librarian creates libraries with the modules in the order in which they were given on the command line. When updating a library the order of the modules is preserved. Any new modules added to a library after it has been created will be appended to the end.

The ordering of the modules in a library is significant to the linker. If a library contains a module which references a symbol defined in another module in the same library, the module defining the symbol should come after the module referencing the symbol.

6.10.7 Error Messages

LIBR issues various error messages, most of which represent a fatal error, while some represent a harmless occurrence which will nonetheless be reported unless the -W option was used. In this case all warning messages will be suppressed.

6.11 Objtohex

The HI-TECH linker is capable of producing simple binary files, or object files as output. Any other format required must be produced by running the utility program OBJTOHEX. This allows conversion of object files as produced by the linker into a variety of different formats, including various hex formats. The program is invoked thus:

```
objtohex options inputfile outputfile
```

All of the arguments are optional. If outputfile is omitted it defaults to l.hex or l.bin depending on whether the -b option is used. The inputfile defaults to l.obj.

The options for OBJTOHEX are listed in Table 6 - 4 on page 162. Where an address is required, the format is the same as for HLINK.
6.11.1 Checksum Specifications

The checksum specification allows automated checksum calculation. The checksum specification takes the form of several lines, each line describing one checksum. The syntax of a checksum line is:

```
addr1-addr2 where1-where2 +offset
```

All of `addr1`, `addr2`, `where1`, `where2` and `offset` are hex numbers, without the usual H suffix. Such a specification says that the bytes at `addr1` through to `addr2` inclusive should be summed and the sum placed in the locations `where1` through `where2` inclusive. For an 8 bit checksum these two addresses should be the same. For a checksum stored low byte first, `where1` should be less than `where2`, and vice versa. The `+offset` is optional, but if supplied, the value offset will be used to initialise the checksum. Otherwise it is initialised to zero. For example:

<table>
<thead>
<tr>
<th>Option</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-A</td>
<td>Produce an ATDOS .atx output file</td>
</tr>
<tr>
<td>-Bbase</td>
<td>Produce a binary file with offset of base. Default file name is l.obj</td>
</tr>
<tr>
<td>-Cckfile</td>
<td>Read a list of checksum specifications from ckfile or standard input</td>
</tr>
<tr>
<td>-D</td>
<td>Produce a COD file</td>
</tr>
<tr>
<td>-E</td>
<td>Produce an MS-DOS .exe file</td>
</tr>
<tr>
<td>-Ffill</td>
<td>Fill unused memory with words of value fill - default value is 0FFh</td>
</tr>
<tr>
<td>-I</td>
<td>Produce an Intel HEX file with linear addressed extended records.</td>
</tr>
<tr>
<td>-L</td>
<td>Pass relocation information into the output file (used with .exe files)</td>
</tr>
<tr>
<td>-M</td>
<td>Produce a Motorola HEX file (S19, S28 or S37 format)</td>
</tr>
<tr>
<td>-N</td>
<td>Produce an output file for Minix</td>
</tr>
<tr>
<td>-Pstk</td>
<td>Produce an output file for an Atari ST, with optional stack size</td>
</tr>
<tr>
<td>-R</td>
<td>Include relocation information in the output file</td>
</tr>
<tr>
<td>-Sfile</td>
<td>Write a symbol file into file</td>
</tr>
<tr>
<td>-T</td>
<td>Produce a Tektronix HEX file. -TE produces an extended TekHEX file.</td>
</tr>
<tr>
<td>-U</td>
<td>Produce a COFF output file</td>
</tr>
<tr>
<td>-UB</td>
<td>Produce a UBROF format file</td>
</tr>
<tr>
<td>-V</td>
<td>Reverse the order of words and long words in the output file</td>
</tr>
<tr>
<td>-x</td>
<td>Create an x.out format file</td>
</tr>
<tr>
<td>-n,m</td>
<td>Format either Motorola or Intel HEX file where n is the max number of bytes per record and m specifies the multiple to which the record size is rounded</td>
</tr>
</tbody>
</table>
0005-1FFF 3–4 +1FFF

This will sum the bytes in 5 through 1FFFH inclusive, then add 1FFFH to the sum. The 16 bit checksum will be placed in locations 3 and 4, low byte in 3. The checksum is initialised with 1FFFH to provide protection against an all zero ROM, or a ROM misplaced in memory. A run time check of this checksum would add the last address of the ROM being checksummed into the checksum. For the ROM in question, this should be 1FFFH. The initialization value may, however, be used in any desired fashion.

6.12 Cref

The cross reference list utility CREF is used to format raw cross-reference information produced by the compiler or the assembler into a sorted listing. A raw cross-reference file is produced with the –CR option to the compiler. The assembler will generate a raw cross-reference file with a –C option (most assemblers) or by using an OPT CREF directive (6800 series assemblers) or a XREF control line (PIC assembler). The general form of the CREF command is:

```
cref options files
```

where options is zero or more options as described below and files is one or more raw cross-reference files. CREF takes the options listed in Table 6 - 5 on page 163. Each option is described in more detail in the following paragraphs.

**Table 6 - 5 Cref Options**

<table>
<thead>
<tr>
<th>Option</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-Fprefix</td>
<td>Exclude symbols from files with a pathname or filename starting with prefix</td>
</tr>
<tr>
<td>-Hheading</td>
<td>Specify a heading for the listing file</td>
</tr>
<tr>
<td>-Llen</td>
<td>Specify the page length for the listing file</td>
</tr>
<tr>
<td>-Ooutfile</td>
<td>Specify the name of the listing file</td>
</tr>
<tr>
<td>-Pwidth</td>
<td>Set the listing width</td>
</tr>
<tr>
<td>-Sstoplist</td>
<td>Read file stoplist and ignore any symbols listed.</td>
</tr>
<tr>
<td>-Xprefix</td>
<td>Exclude any symbols starting with the given prefix</td>
</tr>
</tbody>
</table>

**6.12.1 -Fprefix**

It is often desired to exclude from the cross-reference listing any symbols defined in a system header file, e.g. `<stdio.h>`. The -F option allows specification of a path name prefix that will be used to exclude any symbols defined in a file whose path name begins with that prefix. For example, -F\ will exclude any symbols from all files with a path name starting with \. 
6.12.2 -Hheading

The -H option takes a string as an argument which will be used as a header in the listing. The default heading is the name of the first raw cross-ref information file specified.

6.12.3 -Llen

Specify the length of the paper on which the listing is to be produced, e.g. if the listing is to be printed on 55 line paper you would use a -L55 option. The default is 66 lines.

6.12.4 -Ooutfile

Allows specification of the output file name. By default the listing will be written to the standard output and may be redirected in the usual manner. Alternatively outfile may be specified as the output file name.

6.12.5 -Pwidth

This option allows the specification of the width to which the listing is to be formatted, e.g. -P132 will format the listing for a 132 column printer. The default is 80 columns.

6.12.6 -Sstoplist

The -S option should have as its argument the name of a file containing a list of symbols not to be listed in the cross-reference. Multiple stoplists may be supplied with multiple -S options.

6.12.7 -Xprefix

The -X option allows the exclusion of symbols from the listing, based on a prefix given as argument to -X. For example if it was desired to exclude all symbols starting with the character sequence xyz then the option -Xxyz would be used. If a digit appears in the character sequence then this will match any digit in the symbol, e.g. -XX0 would exclude any symbols starting with the letter X followed by a digit.

CREF will accept wildcard filenames and I/O redirection. Long command lines may be supplied by invoking CREF with no arguments and typing the command line in response to the cref> prompt. A backslash at the end of the line will be interpreted to mean that more command lines follow.

6.13 Cromwell

The CROMWELL utility converts code and symbol files into different formats. The formats available are shown in Table 6 - 6.

The general form of the CROMWELL command is:

```
cromwell options input_files -okey output_file
```
where *options* can be any of the options shown in Table 6 - 7. *Output_file* (optional) is the name of the output file. The *input_files* are typically the HEX and SYM file. CROMWELL automatically searches for the SDB files and reads those if they are found. The options are further described in the following paragraphs.

### 6.13.1 -Pname

The `-P` options takes a string which is the name of the processor used. CROMWELL may use this in the generation of the output format selected.

### 6.13.2 -D

The `-D` option is used to display to the screen details about the named input file in a readable format. The input file can be one of the file types as shown in Table 6 - 6.

### 6.13.3 -C

This option will attempt to identify if the specified input files are one of the formats as shown in Table 6 - 6. If the file is recognised, a confirmation of its type will be displayed.

### 6.13.4 -F

When generating a COD file, this option can be used to force all local symbols to be represented as global symbols. The may be useful where an emulator cannot read local symbol information from the COD file.

---

**Table 6 - 6 Format Types**

<table>
<thead>
<tr>
<th>Key</th>
<th>Format</th>
</tr>
</thead>
<tbody>
<tr>
<td>cod</td>
<td><em>Bytecraft</em> COD file</td>
</tr>
<tr>
<td>coff</td>
<td>COFF file format</td>
</tr>
<tr>
<td>elf</td>
<td>ELF/DWARF file</td>
</tr>
<tr>
<td>eomf51</td>
<td>Extended OMF-51 format</td>
</tr>
<tr>
<td>hitech</td>
<td>HI-TECH Software format</td>
</tr>
<tr>
<td>icoff</td>
<td>ICOFF file format</td>
</tr>
<tr>
<td>ihex</td>
<td><em>Intel</em> HEX file format</td>
</tr>
<tr>
<td>omf51</td>
<td>OMF-51 file format</td>
</tr>
<tr>
<td>pe</td>
<td>P&amp;E file format</td>
</tr>
<tr>
<td>s19</td>
<td><em>Motorola</em> HEX file format</td>
</tr>
</tbody>
</table>
6.13.5 -Okey

This option specifies the format of the output file. The `key` can be any of the types listed in Table 6-6.

6.13.6 -Ikey

This option can be used to specify the default input file format. The `key` can be any of the types listed in Table 6-6.

6.13.7 -L

Use this option to show what file format types are supported. A list similar to that given in Table 6-6 will be shown.

6.13.8 -E

Use this option to tell CROMWELL to ignore any filename extensions that were given. The default extension will be used instead.

6.13.9 -B

In formats that support different endian types, use this option to specify big-endian byte ordering.

6.13.10 -M

When generating COD files this option will remove the preceding underscore character from symbols.

---

**Table 6-7 Cromwell Options**

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-Pname</td>
<td>Processor name</td>
</tr>
<tr>
<td>-D</td>
<td>Dump input file</td>
</tr>
<tr>
<td>-C</td>
<td>Identify input files only</td>
</tr>
<tr>
<td>-F</td>
<td>Fake local symbols as globals</td>
</tr>
<tr>
<td>-Okey</td>
<td>Set the output format</td>
</tr>
<tr>
<td>-Ikey</td>
<td>Set the input format</td>
</tr>
<tr>
<td>-L</td>
<td>List the available formats</td>
</tr>
<tr>
<td>-E</td>
<td>Strip file extensions</td>
</tr>
<tr>
<td>-B</td>
<td>Specify big-endian byte ordering</td>
</tr>
<tr>
<td>-M</td>
<td>Strip underscore character</td>
</tr>
<tr>
<td>-V</td>
<td>Verbose mode</td>
</tr>
</tbody>
</table>
6.13.11 -V

Turns on verbose mode which will display information about operations CROMWELL is performing.

6.14 Memmap

MEMMAP has been individualized for each processor. The MEMMAP program that appears in your BIN directory will conform with the following criteria; XXmap.exe where XX stands for the processor type. From here on, we will be referring to this application as MEMMAP, as to cover all processors.

At the end of compilation and linking, HPD and the command line compiler produce a summary of memory usage. If, however, the compilation is performed in separate stages and the linker is invoked explicitly, this memory information is not displayed. The MEMMAP program reads the information stored in the map file and produces either a summary of psect address allocation or a memory map of program sections similar to that shown by HPD and the command line compiler.

6.14.1 Using MEMMAP

A command to the memory usage program takes the form:

```
memmap options file
```

Options is zero or more MEMMAP options which are listed in Table 6 - 8 on page 167. File is the name of a map file. Only one map file can be processed by MEMMAP.

### Table 6 - 8 Memmap options

<table>
<thead>
<tr>
<th>Option</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>-P</td>
<td>Print psect usage map</td>
</tr>
<tr>
<td>-W wid</td>
<td>Specifies width to which address are printed</td>
</tr>
</tbody>
</table>

6.14.1.1 -P

The default behaviour of MEMMAP is to produce a segment memory map. This output is similar to that printed by HPD and the command line compiler after compilation and linking. This behaviour can be changed by using the -P option. This forces a psect usage map to be printed. The output in this case will be similar to that shown by the HPD’s Memory Usage Map item under the Utility menu or if the -PSECTMAP option is used with the command line compiler.

6.14.1.2 -W wid

The width to which addresses are printed can be adjusted by using the -w option. The default width is determined in respect to the processor’s address range. Depending on the type of processor used,
determines the default width of the printed address, for example a processor with less than or equal to 64k will have a default width of 4. Whereas a processor with greater than 64k may have a default value of 6 digits.
Error Messages

This chapter lists all possible error messages from the HI-TECH C compiler, with an explanation of each one. The name of the applications that could have produced the error are listed in brackets opposite the error message. The tutorial chapter describes the function of each application.

'. ' expected after '..' (Parser)
The only context in which two successive dots may appear is as part of the ellipsis symbol, which must have 3 dots.

'case' not in switch (Parser)
A case label has been encountered but there is no enclosing switch statement. A case label may only appear inside the body of a switch statement.

'default' not in switch (Parser)
A label has been encountered called "default" but it is not enclosed by a switch statement. The label "default" is only legal inside the body of a switch statement.

'with=' flags are cyclic (Assembler)
If Psect A is to be placed 'with' Psect B, and Psect B is to be placed 'with' Psect A, there is no hierarchy. Remove a 'with' flag from one of the psect declarations.

( expected (Parser)
An opening parenthesis was expected here. This must be the first token after a while, for, if, do or asm keyword.

) expected (Parser)
A closing parenthesis was expected here. This may indicate you have left out a parenthesis in an expression, or you have some other syntax error.

*: no match (Preprocessor, Parser)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

, expected (Parser)
A comma was expected here. This probably means you have left out the comma between two identifiers in a declaration list. It may also mean that the immediately preceding type name is misspelled, and has thus been interpreted as an identifier.
**Error Messages**

-s, too few values specified in *  
*Preprocessor*

The list of values to the preprocessor -S option is incomplete. This should not happen if the preprocessor is being invoked by the compiler driver or HPD.

-s, too many values, * unused  
*Preprocessor*

There were too many values supplied to a -S preprocessor option.

... illegal in non-prototype arg list  
*Parser*

The ellipsis symbol may only appear as the last item in a prototyped argument list. It may not appear on its own, nor may it appear after argument names that do not have types.

: expected  
*Parser*

A colon is missing in a case label, or after the keyword "default". This often occurs when a semicolon is accidentally typed instead of a colon.

; expected  
*Parser*

A semicolon is missing here. The semicolon is used as a terminator in many kinds of statements, e.g. do .. while, return etc.

= expected  
*Code Generator, Assembler*

An equal sign was expected here.

#define syntax error  
*Preprocessor*

A macro definition has a syntax error. This could be due to a macro or formal parameter name that does not start with a letter or a missing closing parenthesis (’)’).

#elif may not follow #else  
*Preprocessor*

If a #else has been used after #if, you cannot then use a #elif in the same conditional block.

#elif must be in an #if  
*Preprocessor*

#elif must be preceded by a matching #if line. If there is an apparently corresponding #if line, check for things like extra #endif’s, or improperly terminated comments.

#else may not follow #else  
*Preprocessor*

There can be only one #else corresponding to each #if.

#else must be in an #if  
*Preprocessor*

#else can only be used after a matching #if.

#endif must be in an #if  
*Preprocessor*

There must be a matching #if for each #endif. Check for the correct number of #ifs.

#error: *  
*Preprocessor*

This is a programmer generated error; there is a directive causing a deliberate error. This is normally used to check compile time defines etc.
#if ... sizeof() syntax error
(Preprocessor)
The preprocessor found a syntax error in the argument to sizeof, in a #if expression. Probable causes are mismatched parentheses and similar things.

#if ... sizeof: bug, unknown type code *
(Preprocessor)
The preprocessor has made an internal error in evaluating a sizeof() expression. Check for a malformed type specifier.

#if ... sizeof: illegal type combination
(Preprocessor)
The preprocessor found an illegal type combination in the argument to sizeof() in a #if expression. Illegal combinations include such things as "short long int".

#if bug, operand = *
(Preprocessor)
The preprocessor has tried to evaluate an expression with an operator it does not understand. This is an internal error.

#if sizeof() error, no type specified
(Preprocessor)
Sizeof() was used in a preprocessor #if expression, but no type was specified. The argument to sizeof() in a preprocessor expression must be a valid simple type, or pointer to a simple type.

#if sizeof, unknown type *
(Preprocessor)
An unknown type was used in a preprocessor sizeof(). The preprocessor can only evaluate sizeof() with basic types, or pointers to basic types.

#if value stack overflow
(Preprocessor)
The preprocessor filled up its expression evaluation stack in a #if expression. Simplify the expression - it probably contains too many parenthesized subexpressions.

#if, ifdef, or ifndef without an argument
(Preprocessor)
The preprocessor directives #if, ifdef and ifndef must have an argument. The argument to #if should be an expression, while the argument to ifdef or ifndef should be a single name.

#include syntax error
(Preprocessor)
The syntax of the filename argument to #include is invalid. The argument to #include must be a valid file name, either enclosed in double quotes (""") or angle brackets (< >). For example:

#include "afile.h"
#include <otherfile.h>

Spaces should not be included, and the closing quote or bracket must be present. There should be nothing else on the line.

#include file * was converted to lower case
(Preprocessor)
The #include file name had to be converted to lowercase before it could be opened.
Error Messages

] expected (Parser)
A closing square bracket was expected in an array declaration or an expression using an array index.

{ expected (Parser)
An opening brace was expected here.

} expected (Parser)
A closing brace was expected here.

a macro name cannot also be a label (Assembler)
A label has been found with the same name as a macro. This is not allowed.

a maximum of * reserved areas are allowed. remainder of -RES* ignored (Driver)
Too many address ranges were specified with either the -RESROM or -RESRAM option.

A maximum of * ROM banks are allowed. Remainder of -ROM option ignored (Driver)
Too many ranges were specified with the -ROM option.

a parameter may not be a function (Parser)
A function parameter may not be a function. It may be a pointer to a function, so perhaps a "*" has been omitted from the declaration.

a psect may only be in one class (Assembler)
You cannot assign a psect to more than one class. The psect was defined differently at this point than when it was defined elsewhere.

a psect may only have one 'with' option (Assembler)
A psect can only be placed 'with' one other psect.

absolute expression required (Assembler)
An absolute expression is required in this context.

add_reloc - bad size (Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

ambiguous chip type * -> * or * (Driver)
The chip type specified on the command line is not complete and could refer to more than one chip. Specify the full name of the chip type.

ambiguous format name '*’ (Cromwell)
The output format specified to Cromwell is ambiguous.

argument * conflicts with prototype (Parser)
The argument specified (argument 1 is the left most argument) of this function declaration does not agree with a previous prototype for this function.
argument -w* ignored
The argument to the linker option -w is out of range. For warning levels, the range is -9 to 9. For the map file width, the range is greater than or equal to 10.

argument list conflicts with prototype
The argument list in a function definition is not the same as a previous prototype for that function. Check that the number and types of the arguments are all the same.

argument redeclared: *
The specified argument is declared more than once in the same argument list.

argument too long
This is an internal compiler error. Contact HI-TECH Software technical support with details.

arithmetic overflow in constant expression
A constant expression has been evaluated by the code generator that has resulted in a value that is too big for the type of the expression, e.g. trying to store the value 256 in a "char".

array dimension on * ignored
An array dimension on a function parameter has been ignored because the argument is actually converted to a pointer when passed. Thus arrays of any size may be passed.

array dimension redeclared
An array dimension has been declared as a different non-zero value from its previous declaration. It is acceptable to redeclare the size of an array that was previously declared with a zero dimension, but not otherwise.

array index out of bounds
An array is being indexed with a constant value that is less than zero, or greater than or equal to the number of elements in the array.

assertion
This is an internal compiler error. Contact HI-TECH Software technical support with details.

assertion failed: *
This is an internal compiler error. Contact HI-TECH Software technical support with details.

attempt to modify const object
Objects declared "const" may not be assigned to or modified in any other way.

auto variable * should not be qualified
An auto variable should not have qualifiers such as "near" or "far" associated with it. Its storage class is implicitly defined by the stack organization.
**Error Messages**

**bad #if ... defined() syntax**  
(Preprocessor)  
The defined() pseudo-function in a preprocessor expression requires its argument to be a single name. The name must start with a letter. It should be enclosed in parentheses.

**bad '-p' format**  
(Linker)  
The "-P" option given to the linker is malformed.

**bad -A option: ***  
(Driver)  
The format of a -A option to shift the ROM image was not correct. The -A should be immediately followed by a valid hex number.

**bad -a spec: ***  
(Linker)  
The format of a -A specification, giving address ranges to the linker, is invalid. The correct format is:

```
-Aclass=low-high
```

where class is the name of a psect class, and low and high are hex numbers.

**bad -m option: ***  
(Code Generator)  
The code generator has been passed a -M option that it does not understand. This should not happen if it is being invoked by a standard compiler driver.

**bad -q option ***  
(Parser)  
The first pass of the compiler has been invoked with a -Q option, to specify a type qualifier name, that is badly formed.

**bad -RES* arguments**  
(Driver)  
The address ranges specified to either the -RESROM or -RESRAM option are invalid.

**Bad -ROM arguments ***  
(Driver)  
The arguments to -ROM were either not present or badly formed.

**bad arg * to tysize**  
(Parser)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad arg to e: ***  
(Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**Bad arg to extraspecial ?**  
(Code Generator)  
This is an internal code generator error. Contact HI-TECH technical support with full details of the code that caused this error.

**bad arg to im**  
(Assembler)  
The opcode "IM" only takes the constants 0, 1 or 2 as an argument.
bad bconfloat - *  (Code Generator)
This is an internal code generator error. Contact HI-TECH technical support with full details of the code that caused this error.

bad bit number  (Assembler, Optimiser)
A bit number must be an absolute expression in the range 0-7.

bad bitfield type  (Parser)
A bitfield may only have a type of int.

bad character const  (Parser, Assembler, Optimiser)
This character constant is badly formed.

bad character constant in expression  (Assembler)
The character constant was expected to consist of only one character, but was found to be greater than one character.

bad character in extended tekhex line *  (Objtohex)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad checksum specification  (Linker)
A checksum list supplied to the linker is syntactically incorrect.

bad combination of flags  (Objtohex)
The combination of options supplied to objtohex is invalid.

bad common spec in -p option  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad complex range check  (Linker)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad complex relocation  (Linker)
The linker has been asked to perform complex relocation that is not syntactically correct. Probably means a corrupted object file.

bad confloat - *  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad conval - *  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad dimensions  (Code Generator)
The code generator has been passed a declaration that results in an array having a zero dimension.
**Error Messages**

**bad dp/nargs in openpar: c = ***  
*(Preprocessor)*  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad element count expr**  
*(Code Generator)*  
There is an error in the intermediate code. Try re-installing the compiler from the distribution disks, as this could be caused by a corrupted file.

**bad extraspecial ***  
*(Code Generator)*  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad fixup value**  
*(Optimiser)*  
The assembler file passed to the optimizer is invalid.

**bad float operand size**  
*(Assembler)*  
The maximum size of a float is 4 bytes.

**bad format for -p option**  
*(Code Generator)*  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad gn**  
*(Code Generator)*  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad high address in -a spec**  
*(Linker)*  
The high address given in a -A specification is invalid: it should be a valid number, in decimal, octal or hexadecimal radix. The radix is specified by a trailing O (for octal) or H for hex. Decimal is default.

**bad int. code**  
*(Code Generator)*  
The code generator has been passed input that is not syntactically correct.

**bad list record type ***

**bad load address in -a spec**  
*(Linker)*  
The load address given in a -A specification is invalid: it should be a valid number, in decimal, octal or hexadecimal radix. The radix is specified by a trailing O (for octal) or H for hex. Decimal is default.

**bad low address in -a spec**  
*(Linker)*  
The low address given in a -A specification is invalid: it should be a valid number, in decimal, octal or hexadecimal radix. The radix is specified by a trailing O (for octal) or H for hex. Decimal is default.

**bad min (+) format in spec**  
*(Linker)*  
The minimum address specification in the linker’s -p option is badly formatted.

**bad mod ’+’ for how = ***  
*(Code Generator)*  
This is an internal compiler error. Contact HI-TECH Software technical support with details.
**bad non-zero node in call graph**  
* (Linker)  
The linker has encountered a top level node in the call graph that is referenced from lower down in the call graph. This probably means the program has indirect recursion, which is not allowed when using a compiled stack.

**bad object code format**  
* (Linker)  
The object code format of this object file is invalid. This probably means it is either truncated, corrupted, or not a HI-TECH object file.

**bad op * to revlog**  
* (Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad op * to swaplog**  
* (Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad op: "***"**  
* (Code Generator)  
This is caused by an error in the intermediate code file. You may have run out of disk space for temporary files.

**bad operand**  
* (Optimiser)  
This operand is invalid. Check the syntax.

**bad origin format in spec**  
* (Linker)  
The origin format in a -p option is not a validly formed decimal, octal or hex number. A hex number must have a trailing H.

**bad overrun address in -a spec**  
* (Linker)  
The overrun address given in a -A specification is invalid: it should be a valid number, in decimal, octal or hexadecimal radix. The radix is specified by a trailing O (for octal) or H for hex. Decimal is default.

**bad popreg: * **  
* (Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad pragma * **  
* (Code Generator)  
The code generator has been passed a "pragma" directive that it does not understand.

**bad pushreg: * **  
* (Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad putwsize**  
* (Code Generator)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad record type * **  
* (Linker)  
This indicates that the object file is not a valid HI-TECH object file.
**Error Messages**

**bad relocation type**  
*(Assembler)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad repeat count in -a spec**  
*(Linker)*
The repeat count given in a -A specification is invalid: it should be a valid decimal number.

**bad ret_mask**  
*(Code Generator)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad segment fixups**  
*(Objtohex)*
This is an obscure message from objtohex that is not likely to occur in practice.

**bad segspec ***  
*(Linker)*
The segspec option (-G) to the linker is invalid. The correct form of a segspec option is along the following lines:

-Gnxc+o

where n stands for the segment number, x is a multiplier symbol, c is a constant (multiplier) and o is a constant offset. For example the option

-Gnx4+16

would assign segment selectors starting from 16, and incrementing by 4 for each segment, i.e. in the order 16, 20, 24 etc.

**bad size in -s option**  
*(Linker)*
The size part of a -S option is not a validly formed number. The number must be a decimal, octal or hex number. A hex number needs a trailing H, and an octal number a trailing O. All others are assumed to be decimal.

**bad size in index_type**  
*(Parser)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**bad size list**  
*(Parser)*
The first pass of the compiler has been invoked with a -Z option, specifying sizes of types, that is badly formed.

**bad storage class**  
*(Code Generator)*
The storage class "auto" may only be used inside a function. A function parameter may not have any storage class specifier other than "register". If this error is issued by the code generator, it could mean that the intermediate code file is invalid. This could be caused by running out of disk space.
bad string * in psect pragma  (Code Generator)
The code generator has been passed a "pragma psect" directive that has a badly formed string. "Pragma psect" should be followed by something of the form "oldname=newname".

bad switch size *  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad sx  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad u usage  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad uconval - *  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

bad variable syntax  (Code Generator)
There is an error in the intermediate code file. This could be caused by running out of disk space for temporary files.

bad which * after i  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

banked/common conflict  (Assembler)
The assembler has found conflicting information that suggests that a symbol is located in the access bank, but also in the banked RAM area.

binary digit expected  (Parser)
A binary digit was expected. The format for a binary number is 0Bxxx where xxx is a string containing zeroes and/or ones, e.g.

0B0110

bit address overflow  (Assembler)
The bit address supplied is outside the XA bit space. The bit address should be in the bit space which has a range 0h to 3FFh.

bit field too large (* bits)  (Code Generator)
The maximum number of bits in a bit field is the same as the number of bits in an "int".

bit range check failed *  (Linker)
The bit addressing was out of range.
Error Messages

**bit variables must be global or static**  
(Code Generator)  
A bit variable cannot be of type auto. If you require a bit variable with scope local to a block of code or function, qualify it static.

**bitfield comparison out of range**  
(Code Generator)  
This is the result of comparing a bitfield with a value when the value is out of range of the bitfield. For example, comparing a 2-bit bitfield to the value 5 will never be true as a 2-bit bitfield has a range from 0 to 3,

**bug: illegal __ macro**  
(Preprocessor)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**c= must specify a positive constant**  
(Assembler)  
The parameter to the LIST assembler control’s ’C’ option (which sets the column width of the listing output) must be a positive constant number.

**call depth exceeded by**  
(Linker)  
The call graph shows that functions are nested to a depth greater than specified.

**can’t allocate memory for arguments**  
(Preprocessor, Parser)  
The compiler could not allocate any more memory. Try increasing the size of available memory.

**can’t be both far and near**  
(Parser)  
It is illegal to qualify a type as both far and near.

**can’t be long**  
(Parser)  
Only "int" and "float" can be qualified with "long". Thus combinations like "long char" are illegal.

**can’t be register**  
(Parser)  
Only function parameters or auto (local) variables may be declared "register".

**can’t be short**  
(Parser)  
Only "int" can be modified with short. Thus combinations like "short float" are illegal.

**can’t be unsigned**  
(Parser)  
There is no such thing as an unsigned floating point number.

**can’t call an interrupt function**  
(Parser)  
A function qualified "interrupt" can’t be called from other functions. It can only be called by a hardware (or software) interrupt. This is because an interrupt function has special function entry and exit code that is appropriate only for calling from an interrupt. An "interrupt" function can call other non-interrupt functions.

**can’t create**  
(Code Generator, Assembler, Linker, Optimiser)  
The named file could not be created. Check that all directories in the path are present.
can’t create cross reference file *  
(Assembler)
The cross reference file could not be created. Check that all directories are present. This can also be caused by the assembler running out of memory.

can’t create temp file  
(Linker)
The compiler was unable to create a temporary file. Check the DOS Environment variable TEMP (and TMP) and verify it points to a directory that exists, and that there is space available on that drive. For example, AUTOEXEC.BAT should have something like:

SET TEMP=C:\TEMP

where the directory C:\TEMP exists.

can’t create temp file *  
(Code Generator)
The compiler could not create the temporary file named. Check that all the directories in the file path exist.

can’t enter abs psect  
(Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

can’t find op  
(Assembler, Optimiser)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

can’t find space for psect * in segment *  
(Linker)
The named psect cannot be placed in the specified segment. This either means that the memory associated with the segment has been filled, or that the psect cannot be positioned in any of the available gaps in the memory. Split large functions (for CODE segments) in several smaller functions and ensure that the optimizers are being used.

can’t generate code for this expression  
(Code Generator)
This expression is too difficult for the code generator to handle. Try simplifying the expression, e.g. using a temporary variable to hold an intermediate result.

can’t have 'port' variable: *  
(Code Generator)
The qualifier "port" can be used only with pointers or absolute variables. You cannot define a port variable as the compiler does not allocate space for port variables. You can declare an external port variable.

can’t have 'signed' and 'unsigned' together  
(Parser)
The type modifiers signed and unsigned cannot be used together in the same declaration, as they have opposite meaning.

can’t have an array of bits or a pointer to bit  
(Parser)
It is not legal to have an array of bits, or a pointer to bit.
**Error Messages**

**can’t have array of functions** *(Parser)*

You can’t have an array of functions. You can however have an array of pointers to functions. The correct syntax for an array of pointers to functions is "int (* arrayname[])(");. Note that parentheses are used to associate the star (*) with the array name before the parentheses denoting a function.

**can’t initialize arg** *(Parser)*

A function argument can’t have an initialiser. The initialisation of the argument happens when the function is called and a value is provided for the argument by the calling function.

**can’t initialize bit type** *(Code Generator)*

Variables of type bit cannot be initialised.

**Can’t initialize far variables** *(Code Generator)*

Variables declared as ‘far’ cannot be initialised.

**can’t mix proto and non-proto args** *(Parser)*

A function declaration can only have all prototyped arguments (i.e. with types inside the parentheses) or all K&R style args (i.e. only names inside the parentheses and the argument types in a declaration list before the start of the function body).

**can’t open** *(Linker)*

A file can’t be opened - check spelling.

**can’t open** *(Code Generator, Assembler, Optimiser, Cromwell)*

The named file could not be opened. Check the spelling and the directory path. This can also be caused by running out of memory.

**can’t open * for input** *(Cref)*

Cref cannot open the specified file.

**can’t open * for output** *(Cref)*

Cref cannot open the specified file.

**can’t open avmap file** *(Linker)*

A file required for producing Avocet format symbol files is missing. Try re-installing the compiler.

**can’t open checksum file** *(Linker)*

The checksum file specified to objtohex could not be opened. Check spelling etc.

**can’t open chip info file** *(Assembler)*

The chipinfo file (libpicinfo.ini by default) could not be opened. It may have been incorrectly specified.

**can’t open command file** *(Preprocessor, Linker)*

The command file specified could not be opened for reading. Check spelling!
**can’t open error file** *(Linker)*
The error file specified using the -e option could not be opened.

**can’t open include file** *(Assembler)*
The named include file could not be opened. Check spelling. This can also be caused by running out of memory, or running out of file handles.

**can’t open input file** *(Preprocessor, Assembler)*
The specified input file could not be opened. Check the spelling of the file name.

**can’t open output file** *(Preprocessor, Assembler)*
The specified output file could not be created. This could be because a directory in the path name does not exist.

**can’t reopen** *(Parser)*
The compiler could not reopen a temporary file it had just created.

**can’t seek in** *(Linker)*
The linker can’t seek in the specified file. Make sure the output file is a valid filename.

**can’t take address of register variable** *(Parser)*
A variable declared "register" may not have storage allocated for it in memory, and thus it is illegal to attempt to take the address of it by applying the "&" operator.

**can’t take sizeof func** *(Parser)*
Functions don’t have sizes, so you can’t take use the "sizeof" operator on a function.

**can’t take sizeof(bit)** *(Parser)*
You can’t take sizeof a bit value, since it is smaller than a byte.

**can’t take this address** *(Parser)*
The expression which was the object of the "&" operator is not one that denotes memory storage ("an lvalue") and therefore its address can not be defined.

**can’t use a string in an #if** *(Preprocessor)*
The preprocessor does not allow the use of strings in #if expressions.

**cannot get memory** *(Linker)*
The linker is out of memory! This is unlikely to happen, but removing TSR’s etc. is the cure.

**cannot open** *(Linker)*
A file cannot be opened - check spelling.
Error Messages

cannot open include file *

(Preprocessor)
The named include file could not be opened for reading by the preprocessor. Check the spelling of the filename. If it is a standard header file, not in the current directory, then the name should be enclosed in angle brackets (<>), not quotes.

cast type must be scalar or void

(Parser)
A typecast (an abstract type declarator enclosed in parentheses) must denote a type which is either scalar (i.e. not an array or a structure) or the type "void".

char const too long

(Parser)
A character constant enclosed in single quotes may not contain more than one character.

character not valid at this point in format specifier

(Parser)
The printf() style format specifier has an illegal character.

checksum error in intel hex file *, line *

(Cromwell)
A checksum error was found at the specified line in the specified Intel hex file. The file may have been corrupted.

chip name * not found in chipinfo file

(Driver)
The chip type specified on the command line was not found in the chipinfo INI file. The compiler doesn’t know how to compile for this chip. If this is a device not yet supported by the compiler, you might be able to add the memory specifications to the chipinfo file and try again.

circular indirect definition of symbol *

(Assembler)
The specified symbol has been equated to an external symbol which, in turn, has been equated to the first symbol.

class * memory space redefined: */*

(Assembler)
A class has been defined in two different memory spaces. Either rename one of the classes or, if they are the same class, place them in the same memory space.

close error (disk space?)

(Parser)
When the compiler closed a temporary file, an error was reported. The most likely cause of this is that there was insufficient space on disk for the file.

common symbol may not be in absolute psect

(Assembler)
If a symbol is defined as common, you cannot place it in a psect which is absolute.

common symbol psect conflict: *

(Assembler)
A common symbol has been defined to be in more than one psect.

compiler not installed properly - reinstall and try again

(Driver)
This is a message from the compiler’s security system. Firstly, to move the compiler from one drive to another, or even from one directory to another, you must reinstall. You cannot copy the installed
compiler (even backing up and restoring will not work unless you simply restore over the existing files). If you have reinstalled, then it is possible that you are running an older version of the same compiler still installed on your machine. Check your PATH environment variable to make sure you’re running what you think you are, i.e. make sure your PATH specifies the newly installed compiler.

**complex relocation not supported for -r or -l options yet** *(Linker)*
The linker was given a -R or -L option with file that contain complex relocation. This is not yet supported.

**conflicting fnconf records** *(Linker)*
This is probably caused by multiple run-time startoff module. Check the linker arguments, or "Object Files..." in HPD.

**constant conditional branch** *(Code Generator)*
A conditional branch (generated by an "if" statement etc.) always follows the same path. This may indicate an expression with missing or badly placed parentheses, causing the evaluation to yield a value different to what you expected, or it may be because you have written something like "while(1)". To produce an infinite loop, use "for(;;)".

**constant conditional branch: possible use of = instead of ==** *(Code Generator)*
There is an expression inside an if or other conditional construct, where a constant is being assigned to a variable. This may mean you have inadvertently used an assignment (=) instead of a compare (==).

**constant expression required** *(Parser)*
In this context an expression is required that can be evaluated to a constant at compile time.

**constant left operand to ?** *(Code Generator)*
The left operand to a conditional operator (?) is constant, thus the result of the tertiary operator ?: will always be the same.

**constant operand to || or &&** *(Code Generator)*
One operand to the logical operators || or && is a constant. Check the expression for missing or badly placed parentheses.

**constant relational expression** *(Code Generator)*
There is a relational expression that will always be true or false. This may be because e.g. you are comparing an unsigned number with a negative value, or comparing a variable with a value greater than the largest number it can represent.

**control line * within macro expansion** *(Preprocessor)*
A preprocessor control line (one starting with a #) has been encountered while expanding a macro. This should not happen.
### Error Messages

**conversion to shorter data type**
*(Code Generator)*
Truncation may occur in this expression as the lvalue is of shorter type than the rvalue.

**copyexpr: can’t handle v_rtype = ***
*(Assembler)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**couldn’t create error file: ***
*(Driver)*
The error file specified after the -Efile or -E+file options could not be opened. Check to ensure that the file or directory is not specified read only.

**declaration of * hides outer declaration**
*(Parser)*
An object has been declared that has the same name as an outer declaration (i.e. one outside and preceding the current function or block). This is legal, but can lead to accidental use of one variable when the outer one was intended.

**declarator too complex**
*(Parser)*
This declarator is too complex for the compiler to handle. Examine the declaration and find a way to simplify it. If the compiler finds it too complex, so will anybody maintaining the code.

**default case redefined**
*(Parser)*
There is only allowed to be one "default" label in a switch statement. You have more than one.

**degenerate signed comparison**
*(Code Generator)*
There is a comparision of a signed value with the most negative value possible for this type, such that the comparision will always be true or false. E.g. char c;

if(c >= -128)

will always be true, because an 8 bit signed char has a maximum negative value of -128.

**degenerate unsigned comparison**
*(Code Generator)*
There is a comparision of an unsigned value with zero, which will always be true or false. E.g.

unsigned char c;
if(c >= 0)

will always be true, because an unsigned value can never be less than zero.

**delete what ?**
*(Libr)*
The librarian requires one or more modules to be listed for deletion when using the ’d’ key.

**delta= must specify a positive constant**
*(Assembler)*
The parameter to the PSECT assembler directive’s ’DELTA’ option must be a positive constant number.
did not recognize format of input file  
(Cromwell)
The input file to Cromwell is required to be COD, Intel HEX, Motorola HEX, COFF, OMF51, P&E or HI-TECH.

digit out of range  
(Parser, Assembler, Optimiser)
A digit in this number is out of range of the radix for the number, e.g. using the digit 8 in an octal number, or hex digits A-F in a decimal number. An octal number is denoted by the digit string commencing with a zero, while a hex number starts with "0X" or "0x".

dimension required  
(Parser)
Only the most significant (i.e. the first) dimension in a multi-dimension array may not be assigned a value. All succeeding dimensions must be present.

direct range check failed *  
(Linker)
The direct addressing was out of range.

divide by zero in #if, zero result assumed  
(Preprocessor)
Inside a #if expression, there is a division by zero which has been treated as yielding zero.

division by zero  
(Code Generator)
A constant expression that was being evaluated involved a division by zero.

double float argument required  
(Parser)
The printf format specifier corresponding to this argument is %f or similar, and requires a floating point expression. Check for missing or extra format specifiers or arguments to printf.

ds argument must be a positive constant  
(Assembler)
The argument to the DS assembler directive must be a positive constant.

duplicate * for * in chipinfo file at line *  
(Assembler, Driver)
The chipinfo file has a processor section with multiple values for a field. Only one value is allowed per chip.

duplicate -d or -h flag  
(Linker)
The symbol file name has been specified to the linker for a second time.

duplicate -m flag  
(Linker)
The linker only likes to see one -m flag, unless one of them does not specify a file name. Two map file names are more than it can handle!

duplicate arch for * in chipinfo file at line *  
(Assembler, Driver)
The chipinfo file has a processor section with multiple ARCH values. Only one ARCH value is allowed.

duplicate case label  
(Code Generator)
There are two case labels with the same value in this switch statement.
**Error Messages**

**duplicate case label** *(Code Generator)*
There is more than one case label with this value in a switch statement.

**duplicate fnconf directive** *(Assembler)*

**duplicate label** *(Parser)*
The same name is used for a label more than once in this function. Note that the scope of labels is the entire function, not just the block that encloses a label.

**duplicate lib for in chipinfo file at line** *(Assembler)*
The chipinfo file (libpicinfo.ini by default) has a processor section with multiple LIB values. Only one LIB value is allowed.

**duplicate qualifier** *(Parser)*
There are two occurrences of the same qualifier in this type specification. This can occur either directly or through the use of a typedef. Remove the redundant qualifier.

**duplicate qualifier key** *(Parser)*
This qualifier key (given via a -Q option) has been used twice.

**duplicate qualifier name** *(Parser)*
A duplicate qualifier name has been specified to P1 via a -Q option. This should not occur if the standard compiler drivers are used.

**duplicate romsize for in chipinfo file at line** *(Assembler)*
The chipinfo file (libpicinfo.ini by default) has a processor section with multiple ROMSIZE values. Only one ROMSIZE value is allowed.

**duplicate sparebit for in chipinfo file at line** *(Assembler)*
The chipinfo file (libpicinfo.ini by default) has a processor section with multiple SPAREBIT values. Only one SPAREBIT value is allowed.

**duplicate zeroreg for in chipinfo file at line** *(Assembler)*
The chipinfo file (libpicinfo.ini by default) has a processor section with multiple ZEROREG values. Only one ZEROREG value is allowed.

**empty chip info file** *(Assembler)*
The chipinfo file (libpicinfo.ini by default) contains no data.

**end of file within macro argument from line** *(Preprocessor)*
A macro argument has not been terminated. This probably means the closing parenthesis has been omitted from a macro invocation. The line number given is the line where the macro argument started.
end of string in format specifier  
(PARSER)  
The format specifier for the printf() style function is malformed.

end statement inside include file or macro  
(ASSEMBLER)  
An END statement was found inside an include file or a macro.

entry point multiply defined  
(LINKER)  
There is more than one entry point defined in the object files given the linker.

enum tag or { expected  
(PARSER)  
After the keyword "enum" must come either an identifier that is or will be defined as an enum tag, or an opening brace.

eof in #asm  
(PREPROCESSOR)  
An end of file has been encountered inside a #asm block. This probably means the #endasm is missing or misspelt.

eof in comment  
(PREPROCESSOR)  
End of file was encountered inside a comment. Check for a missing closing comment flag.

eof inside conditional  
(ASSEMBLER)  
END-of-FILE was encountered while scanning for an "endif" to match a previous "if".

eof inside macro def’n  
(ASSEMBLER)  
End-of-file was encountered while processing a macro definition. This means there is a missing "endm" directive.

eof on string file  
(PARSER)  
P1 has encountered an unexpected end-of-file while re-reading its file used to store constant strings before sorting and merging. This is most probably due to running out of disk space. Check free disk space.

error closing output file  
(CODE GENERATOR, OPTIMISER)  
The compiler detected an error when closing a file. This most probably means there is insufficient disk space.

error dumping *  
(CROMWELL)  
Either the input file to Cromwell is of an unsupported type or that file cannot be dumped to the screen.

error in format string  
(PARSER)  
There is an error in the format string here. The string has been interpreted as a printf() style format string, and it is not syntactically correct. If not corrected, this will cause unexpected behaviour at run time.

evaluation period has expired  
(DRIVER)  
The evaluation period for this compiler has expired. Contact HI-TECH to purchase a full licence.
Error Messages

expand - bad how
(Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

expand - bad which
(Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

expected ’-’ in -a spec
(Linker)
There should be a minus sign (-) between the high and low addresses in a -A spec, e.g.

-AROM=1000h-1FFFh

exponent expected
(Parser)
A floating point constant must have at least one digit after the "e" or "E".

expression error
(Code Generator, Assembler, Optimiser)
There is a syntax error in this expression, OR there is an error in the intermediate code file. This could be caused by running out of disk space.

expression generates no code
(Code Generator)
This expression generates no code. Check for things like leaving off the parentheses in a function call.

expression stack overflow at op *
(Preprocessor)
Expressions in #if lines are evaluated using a stack with a size of 128. It is possible for very complex expressions to overflow this. Simplify the expression.

expression syntax
(Parser)
This expression is badly formed and cannot be parsed by the compiler.

expression too complex
(Parser)
This expression has caused overflow of the compiler’s internal stack and should be re-arranged or split into two expressions.

external declaration inside function
(Parser)
A function contains an "extern" declaration. This is legal but is invariably A Bad Thing as it restricts the scope of the function declaration to the function body. This means that if the compiler encounters another declaration, use or definition of the extern object later in the same file, it will no longer have the earlier declaration and thus will be unable to check that the declarations are consistent. This can lead to strange behaviour of your program or signature errors at link time. It will also hide any previous declarations of the same thing, again subverting the compiler’s type checking. As a general rule, always declare "extern" variables and functions outside any other functions.

field width not valid at this point
(Parser)
A field width may not appear at this point in a printf() type format specifier.
file locking not enabled on network drive  (Driver)
The driver has attempted to modify the lock file located in the LIB directory but was unable to do so. This has probably resulted from the network drive used to hold the compiler being read only.

file name index out of range in line no. record  (Cromwell)
The .COD file has an invalid format in the specified record.

filename work buffer overflow  (Preprocessor)
A filename constructed while looking for an include file has exceeded the length of an internal buffer. Since this buffer is 4096 bytes long, this is unlikely to happen.

fixup overflow in expression *  (Linker)
The linker was asked to relocate (fixup) an item that would not fit back into the space after relocation. For example this will occur if a byte size object is initialized with an address that is bigger than 255. This error occurred in a complex expression.

fixup overflow referencing *  (Linker)
The linker was asked to relocate (fixup) an item that would not fit back into the space after relocation. For example this will occur if a byte size object is initialized with an address that is bigger than 255.

float param coerced to double  (Parser)
Where a non-prototyped function has a parameter declared as "float", the compiler converts this into a "double float". This is because the default C type conversion conventions provide that when a floating point number is passed to a non-prototyped function, it will be converted to double. It is important that the function declaration be consistent with this convention.

form length must be >= 15  (Assembler)
The form length specified using the -Flength option must be at least 15 lines. Setting this length to zero turns off paging. Default value is zero.

formal parameter expected after #  (Preprocessor)
The stringization operator # (not to be confused with the leading # used for preprocessor control lines) must be followed by a formal macro parameter. If you need to stringize a token, you will need to define a special macro to do it, e.g.

#define __mkstr__(x) #x

then use __mkstr__(token) wherever you need to convert a token into a string.

function * appears in multiple call graphs: rooted at *  (Linker)
This function can be called from both main line code and interrupt code. Use the reentrant keyword, if this compiler supports it, or recode to avoid using local variables or parameters, or duplicate the function.
function * argument evaluation overlapped  (Linker)
A function call involves arguments which overlap between two functions. This could occur with a call like:

```c
void fn1(void) { fn3( 7, fn2(3), fn2(9)); /* Offending call */ } char fn2(char fred) { return fred + fn3(5,1,0); } char fn3(char one, char two, char three) { return one+two+three; }
```

where fn1 is calling fn3, and two arguments are evaluated by calling fn2, which in turn calls fn3. The structure should be modified to prevent this.

function * is never called  (Linker)
This function is never called. This may not represent a problem, but space could be saved by removing it. If you believe this function should be called, check your source code.

function body expected  (Parser)
Where a function declaration is encountered with K&R style arguments (i.e. argument names but no types inside the parentheses) a function body is expected to follow.

function declared implicit int  (Parser)
Where the compiler encounters a function call of a function whose name is presently undefined, the compiler will automatically declare the function to be of type "int", with unspecified (K&R style) parameters. If a definition of the function is subsequently encountered, it is possible that its type and arguments will be different from the earlier implicit declaration, causing a compiler error. The solution is to ensure that all functions are defined or at least declared before use, preferably with prototyped parameters. If it is necessary to make a forward declaration of a function, it should be preceded with the keywords "extern" or "static" as appropriate.

function does not take arguments  (Parser, Code Generator)
This function has no parameters, but it is called here with one or more arguments.

function is already 'extern'; can’t be 'static'  (Parser)
This function was already declared extern, possibly through an implicit declaration. It has now been redeclared static, but this redeclaration is invalid. If the problem has arisen because of use before definition, either move the definition earlier in the file, or place a static forward definition earlier in the file, e.g. static int fred(void);

function or function pointer required  (Parser)
Only a function or function pointer can be the subject of a function call. This error can be produced when an expression has a syntax error resulting in a variable or expression being followed by an opening parenthesis ("(") which denotes a function call.

functions can’t return arrays  (Parser)
A function can return only a scalar (simple) type or a structure. It cannot return an array.
functions can't return functions  (Parser)
A function cannot return a function. It can return a function pointer. A function returning a pointer to a function could be declared like this: int (* (name()))(). Note the many parentheses that are necessary to make the parts of the declaration bind correctly.

functions nested too deep  (CodeGen)
This error is unlikely to happen with C code, since C cannot have nested functions!

hex digit expected  (Parser)
After "0x" should follow at least one of the hex digits 0-9 and A-F or a-f.

I/O error reading symbol table  (Cromwell)
Cromwell could not read the symbol table. This could be because the file was truncated or there was some other problem reading the file.

ident records do not match  (Linker)
The object files passed to the linker do not have matching ident records. This means they are for different processor types.

identifier expected  (Parser)
Inside the braces of an "enum" declaration should be a comma-separated list of identifiers.

identifier redefined: *  (Parser)
This identifier has already been defined. It cannot be defined again.

identifier redefined: * (from line *)  (Parser)
This identifier has been defined twice. The 'from line' value is the line number of the first declaration.

illegal # command *  (Preprocessor)
The preprocessor has encountered a line starting with #, but which is not followed by a recognized control keyword. This probably means the keyword has been misspelt. Legal control keywords are: assert, asm, define, elif, else, endasm, endif, error, if, ifdef, ifndef, include, line, pragma, undef.

illegal #if line  (Preprocessor)
There is a syntax error in the expression following #if. Check the expression to ensure it is properly constructed.

illegal #undef argument  (Preprocessor)
The argument to #undef must be a valid name. It must start with a letter.

illegal '#' directive  (Preprocessor, Parser)
The compiler does not understand the "#" directive. It is probably a misspelling of a pre-processor "#" directive.
Error Messages

illegal character (* decimal) in #if  
* (Preprocessor)
The #if expression had an illegal character. Check the line for correct syntax.

illegal character *  
* (Parser)
This character is illegal.

illegal character * in #if  
* (Preprocessor)
There is a character in a #if expression that has no business being there. Valid characters are the letters, digits and those comprising the acceptable operators.

illegal conversion  
* (Parser)
This expression implies a conversion between incompatible types, e.g. a conversion of a structure type into an integer.

illegal conversion between pointer types  
* (Parser)
A pointer of one type (i.e. pointing to a particular kind of object) has been converted into a pointer of a different type. This will usually mean you have used the wrong variable, but if this is genuinely what you want to do, use a typecast to inform the compiler that you want the conversion and the warning will be suppressed.

illegal conversion of integer to pointer  
* (Parser)
An integer has been assigned to or otherwise converted to a pointer type. This will usually mean you have used the wrong variable, but if this is genuinely what you want to do, use a typecast to inform the compiler that you want the conversion and the warning will be suppressed.

illegal conversion of pointer to integer  
* (Parser)
A pointer has been assigned to or otherwise converted to an integral type. This will usually mean you have used the wrong variable, but if this is genuinely what you want to do, use a typecast to inform the compiler that you want the conversion and the warning will be suppressed.

illegal flag *  
* (Linker)
This flag is unrecognized.

illegal function qualifier(s)  
* (Parser)
A qualifier such as "const" or "volatile" has been applied to a function. These qualifiers only make sense when used with an lvalue (i.e. an expression denoting memory storage). Perhaps you left out a star ("\*\") indicating that the function should return a pointer to a qualified object.

illegal initialisation  
* (Parser)
You can’t initialise a "typedef" declaration, because it does not reserve any storage that could be initialised.

illegal instruction for this processor  
* (Assembler)
The instruction is not supported by this processor.
illegal operation on a bit variable  
* (Parser)
Not all operations on bit variables are supported. This operation is one of those.

illegal operator in #if  
* (Preprocessor)
A #if expression has an illegal operator. Check for correct syntax.

illegal or too many -g flags  
* (Linker)
There has been more than one -g option, or the -g option did not have any arguments following. The arguments specify how the segment addresses are calculated.

illegal or too many -o flags  
* (Linker)
This -o flag is illegal, or another -o option has been encountered. A -o option to the linker must have a filename. There should be no space between the filename and the -o, e.g. -ofile.obj

illegal or too many -p flags  
* (Linker)
There have been too many -p options passed to the linker, or a -p option was not followed by any arguments. The arguments of separate -p options may be combined and separated by commas.

illegal record type  
* (Linker)
There is an error in an object file. This is either an invalid object file, or an internal error in the linker. Try recreating the object file.

illegal register indirection
The register does not exist.

illegal relocation size: *
* (Linker)
There is an error in the object code format read by the linker. This either means you are using a linker that is out of date, or that there is an internal error in the assembler or linker.

illegal relocation type: *
* (Linker)
An object file contained a relocation record with an illegal relocation type. This probably means the file is corrupted or not an object file.

illegal switch *
* (Code Generator, Assembler, Optimiser)
This command line option was not understood.

illegal type for array dimension  
* (Parser)
An array dimension must be either an integral type or an enumerated value.

illegal type for index expression  
* (Parser)
An index expression must be either integral or an enumerated value.

illegal type for switch expression  
* (Parser)
A "switch" operation must have an expression that is either an integral type or an enumerated value.
illegal use of void expression  (Parser)
A void expression has no value and therefore you can’t use it anywhere an expression with a value is
required, e.g. as an operand to an arithmetic operator.

image too big  (Objtohex)
The program image being constructed by objtohex is too big for its virtual memory system.

implicit conversion of float to integer  (Parser)
A floating point value has been assigned or otherwise converted to an integral type. This could result in
truncation of the floating point value. A typecast will make this warning go away.

implicit return at end of non-void function  (Parser)
A function which has been declared to return a value has an execution path that will allow it to reach the
end of the function body, thus returning without a value. Either insert a return statement with a value, or
if the function is not to return a value, declare it "void".

implicit signed to unsigned conversion  (Parser)
A signed number is being assigned or otherwise converted to a larger unsigned type. Under the ANSI
"value preserving" rules, this will result in the signed value being first sign-extended to a signed number
the size of the target type, then converted to unsigned (which involves no change in bit pattern). Thus
an unexpected sign extension can occur. To ensure this does not happen, first convert the signed value
to an unsigned equivalent, e.g. if you want to assign a signed char to an unsigned int, first typecast the
char value to "unsigned char".

inappropriate 'else'  (Parser)
An "else" keyword has been encountered that cannot be associated with an "if" statement. This may
mean there is a missing brace or other syntactic error.

inappropriate break/continue  (Parser)
A "break" or "continue" statement has been found that is not enclosed in an appropriate control structure.
"continue" can only be used inside a "while", "for" or "do while" loop, while "break" can only be used
inside those loops or a "switch" statement.

include files nested too deep  (Assembler)
Macro expansions and include file handling have filled up the assembler’s internal stack. The maximum
number of open macros and include files is 30.

included file * was converted to lower case  (Preprocessor)
The file specified to be included was not found, but a file with a lowercase version of the name of the
file specified was found and used instead.

incompatible intermediate code version; should be *  (Code Generator)
The intermediate code file produced by P1 is not the correct version for use with this code generator.
This is either that incompatible versions of one or more compilers have been installed in the same
directory, or a temporary file error has occurred leading to corruption of a temporary file. Check the setting of the TEMP environment variable. If it refers to a long path name, change it to something shorter.

incomplete * record body: length = *  
**(Linker)**
An object file contained a record with an illegal size. This probably means the file is truncated or not an object file.

incomplete ident record  
**(Libr)**
The IDENT record in the object file was incomplete.

incomplete record  
**(Objtohex, Libr)**
The object file passed to objtohex or the librarian is corrupted.

incomplete record: *  
**(Linker)**
An object code record is incomplete. This is probably due to a corrupted or invalid object module. Re-compile the source file, watching for out of disk space errors etc.

incomplete record: type = * length = *  
This message is produced by the DUMP or XSTRIP utilities and indicates that the object file is not a valid HI-TECH object file, or that it has been truncated, possibly due to running out of disk or RAMdisk space.

incomplete symbol record  
**(Libr)**
The SYM record in the object file was incomplete.

inconsistent lineno tables  
**(Cromwell)**
This is an internal compiler error. Contact HI-TECH Software technical support with details.

inconsistent storage class  
**(Parser)**
A declaration has conflicting storage classes. Only one storage class should appear in a declaration.

inconsistent symbol tables  
**(Cromwell)**
This is an internal compiler error. Contact HI-TECH Software technical support with details.

inconsistent type  
**(Parser)**
Only one basic type may appear in a declaration, thus combinations like "int float" are illegal.

initialisation syntax  
**(Parser)**
The initialisation of this object is syntactically incorrect. Check for the correct placement and number of braces and commas.

initializer in 'extern' declaration  
**(Parser)**
A declaration containing the keyword "extern" has an initialiser. This overrides the "extern" storage class, since to initialise an object it is necessary to define (i.e. allocate storage for ) it.
**Error Messages**

**insufficient memory for macro def'n**  
(Assembler)  
There is not sufficient memory to store a macro definition.

**integer constant expected**  
(Parser)  
A colon appearing after a member name in a structure declaration indicates that the member is a bitfield. An integral constant must appear after the colon to define the number of bits in the bitfield.

**integer expression required**  
(Parser)  
In an "enum" declaration, values may be assigned to the members, but the expression must evaluate to a constant of type "int".

**integral argument required**  
(Parser)  
An integral argument is required for this format specifier. Check the number and order of format specifiers and corresponding arguments.

**integral type required**  
(Parser)  
This operator requires operands that are of integral type only.

**interrupt function * may only have one interrupt level**  
(Code Generator)  
Only one interrupt level may be associated with an interrupt function. Check to ensure that only one interrupt_level pragma has been used with the function specified.

**interrupt function requires an address**  
(Code Generator)  
The Highend PIC devices support multiple interrupts. An "@ address" is required with the interrupt definition to indicate with which vector this routine is associated.

**interrupt_level should be 0 to 7**  
(Parser)  
The pragma 'interrupt_level’ must have an argument from 0 to 7.

**invalid * limits in chipinfo file at line ***  
(Driver)  
The ranges of addresses for the ram banks or common memory supplied in the chipinfo INI file is not valid for architecture specified.

**invalid address after 'end' directive**  
(Assembler)  
The start address of the program which is specified after the assembler 'end' directive must be a label in the current file.

**invalid argument to float24**  
(Assembler)  
An argument to the float24 directive must be a number or a symbol which has been equated to a number.

**invalid character ("*") in number**  
(Assembler)  
A number contained a character that was not part of the range 0-9 or 0-F.

**invalid disable: ***  
(Preprocessor)  
This is an internal compiler error. Contact HI-TECH Software technical support with details.
invalid format specifier or type modifier  
( Parser)  
The format specifier or modifier in the printf() style string is illegal for this particular format.

invalid hex file: *, line *  
(Cromwell)  
The specified Hex file contains an invalid line.

invalid number syntax  
(Assembler, Optimiser)  
The syntax of a number is invalid. This can be, e.g. use of 8 or 9 in an octal number, or other malformed numbers.

invalid size for fnsize directive  
(Assembler)  
The assembler FNSIZE assembler directive arguments must be positive constants.

inverted common bank in chipinfo file at line *  
(Assembler, Driver)  
The second hex number specified in the COMMON field in the chipinfo file (libpicinfo.ini by default) must be greater in value than the first.

inverted ICD ROM address in chipinfo file at line *  
(Driver)  
The second hex number specified in the ICD ROM address field in the chipinfo file (libpicinfo.ini by default) must be greater in value than the first.

inverted ram bank in chipinfo file at line *  
(Assembler, Driver)  
The second hex number specified in the RAM field in the chipinfo file (libpicinfo.ini by default) must be greater in value than the first.

label identifier expected  
(Parser)  
An identifier denoting a label must appear after "goto".

lexical error  
(Assembler, Optimiser)  
An unrecognized character or token has been seen in the input.

library * is badly ordered  
(Linker)  
This library is badly ordered. It will still link correctly, but it will link faster if better ordered.

library file names should have .lib extension: *  
(Libr)  
Use the .lib extension when specifying a library.

line does not have a newline on the end  
(Parser)  
The last line in the file is missing the newline (linefeed, hex 0A) from the end. Some editors will create such files, which can cause problems for include files. The ANSI C standard requires all source files to consist of complete lines only.

line too long  
(Optimiser)  
This line is too long. It will not fit into the compiler’s internal buffers. It would require a line over 1000 characters long to do this, so it would normally only occur as a result of macro expansion.
**Error Messages**

- **local illegal outside macros** *(Assembler)*
The "LOCAL" directive is only legal inside macros. It defines local labels that will be unique for each invocation of the macro.

- **local psect '*' conflicts with global psect of same name** *(Linker)*
A local psect may not have the same name as a global psect.

- **logical type required** *(Parser)*
The expression used as an operand to "if", "while" statements or to boolean operators like ! and && must be a scalar integral type.

- **long argument required** *(Parser)*
A long argument is required for this format specifier. Check the number and order of format specifiers and corresponding arguments.

- **macro * wasn’t defined** *(Preprocessor)*
A macro name specified in a -U option to the preprocessor was not initially defined, and thus cannot be undefined.

- **macro argument after * must be absolute** *(Assembler)*
The argument after * in a macro call must be absolute, as it must be evaluated at macro call time.

- **macro argument may not appear after local** *(Assembler)*
The list of labels after the directive "LOCAL" may not include any of the formal parameters to the macro.

- **macro expansions nested too deep** *(Assembler)*
Macro expansions in the assembler are nested too deep. The limit is 30 macros and include files nested at one time.

- **macro work area overflow** *(Preprocessor)*
The total length of a macro expansion has exceeded the size of an internal table. This table is normally 8192 bytes long. Thus any macro expansion must not expand into a total of more than 8K bytes.

- **member * redefined** *(Parser)*
This name of this member of the struct or union has already been used in this struct or union.

- **members cannot be functions** *(Parser)*
A member of a structure or a union may not be a function. It may be a pointer to a function. The correct syntax for a function pointer requires the use of parentheses to bind the star ("*" ) to the pointer name, e.g. "int (*name)();".

- **metaregister * can’t be used directly** *(Code Generator)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.
mismatched comparison  (Code Generator)
A comparison is being made between a variable or expression and a constant value which is not in the range of possible values for that expression, e.g. if you compare an unsigned character to the constant value 300, the result will always be false (not equal) since an unsigned character can NEVER equal 300. As an 8 bit value it can represent only 0-255.

misplaced ‘?’ or ‘:’, previous operator is *  (Preprocessor)
A colon operator has been encountered in a #if expression that does not match up with a corresponding ? operator. Check parentheses etc.

misplaced constant in #if  (Preprocessor)
A constant in a #if expression should only occur in syntactically correct places. This error is most probably caused by omission of an operator.

missing ‘)’  (Parser)
A closing parenthesis was missing from this expression.

missing ‘=’ in class spec  (Linker)
A class spec needs an = sign, e.g. -Ctext=ROM

missing ‘]’  (Parser)
A closing square bracket was missing from this expression.

missing arch specification for * in chipinfo file  (Assembler)
The chipinfo file (libpicinfo.ini by default) has a processor section without an ARCH values. The architecture of the processor must be specified.

missing arg to -a  (Parser)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

missing arg to -e  (Linker)
The error file name must be specified following the -e linker option.

missing arg to -i  (Parser)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

missing arg to -j  (Linker)
The maximum number of errors before aborting must be specified following the -j linker option.

missing arg to -q  (Linker)
The -Q linker option requires the machine type for an argument.

missing arg to -u  (Linker)
The -U (undefine) option needs an argument, e.g. -U_symbol
**Error Messages**

**missing arg to -w**  
(Linker)  
The -W option (listing width) needs a numeric argument.

**missing argument to 'pragma psect'**  
(Parser)  
The pragma 'psect' requires an argument of the form oldname=newname where oldname is an existing psect name known to the compiler, and newname is the desired new name. Example: #pragma psect bss=battery

**missing argument to 'pragma switch'**  
(Parser)  
The pragma 'switch' requires an argument of auto, direct or simple.

**missing basic type: int assumed**  
(Parser)  
This declaration does not include a basic type, so int has been assumed. This declaration is not illegal, but it is preferable to include a basic type to make it clear what is intended.

**missing key in avmap file**  
(Linker)  
A file required for producing Avocet format symbol files is corrupted. Try re-installing the compiler.

**missing memory key in avmap file**  
(Linker)  
A file required for producing Avocet format symbol files is corrupted. Try re-installing the compiler.

**missing name after pragma 'inline'**  
(Parser)  
The 'inline' pragma has the syntax:

```
#pragma inline func_name
```

where func_name is the name of a function which is to be expanded to inline code. This pragma has no effect except on functions specially recognized by the code generator.

**missing name after pragma 'printf_check'**  
(Parser)  
The pragma 'printf_check', which enable printf style format string checking for a function, requires a function name, e.g.

```
#pragma printf_check sprintf
```

**missing newline**  
(Preprocessor)  
A new line is missing at the end of the line. Each line, including the last line, must have a new line at the end. This problem is normally introduced by editors.

**missing number after % in -p option**  
(Linker)  
The % operator in a -p option (for rounding boundaries) must have a number after it.

**missing number after pragma 'pack'**  
(Parser)  
The pragma 'pack' requires a decimal number as argument. For example

```
#pragma pack(1)
```
will prevent the compiler aligning structure members onto anything other than one byte boundaries. Use this with caution as some processors enforce alignment and will not operate correctly if word fetches are made on odd boundaries (e.g. 68000, 8096).

**missing number after pragma interrupt_level** *(Parser)*
Pragma ‘interrupt_level’ requires an argument from 0 to 7.

**missing processor spec after -p** *(Cromwell)*
The -p option to cromwell must specify a processor.

**mod by zero in #if, zero result assumed** *(Preprocessor)*
A modulus operation in a #if expression has a zero divisor. The result has been assumed to be zero.

**module * defines no symbols** *(Libr)*
No symbols were found in the module’s object file.

**module has code below file base of *** *(Linker)*
This module has code below the address given, but the -C option has been used to specify that a binary output file is to be created that is mapped to this address. This would mean code from this module would have to be placed before the beginning of the file! Check for missing psect directives in assembler files.

**multi-byte constant * isn’t portable** *(Preprocessor)*
Multi-byte constants are not portable, and in fact will be rejected by later passes of the compiler.

**multiple free: *** *(Code Generator)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**multiply defined symbol *** *(Assembler, Linker)*
This symbol has been defined in more than one place in this module.

**n= must specify a positive constant** *(Assembler)*
The parameter to the LIST assembler control’s ‘N’ option (which sets the page length for the listing output) must be a positive constant number.

**nested #asm directive** *(Preprocessor)*
It is not legal to nest #asm directives. Check for a missing or misspelt #endasm directive.

**nested comments** *(Preprocessor)*
This warning is issued when nested comments are found. A nested comment may indicate that a previous closing comment marker is missing or malformed.

**no #asm before #endasm** *(Preprocessor)*
A #endasm operator has been encountered, but there was no previous matching #asm.
Error Messages

**no addresses specified with -ROM option** *(Driver)*
No addresses ranges were specified with the -ROM option.

**No addresses specified with -RES* option** *(Driver)*
No address ranges were specified to either the -RESROM or -RESRAM option.

**no case labels** *(Code Generator)*
There are no case labels in this switch statement.

**no end record** *(Linker)*
This object file has no end record. This probably means it is not an object file.

**no end record found** *(Linker)*
An object file did not contain an end record. This probably means the file is corrupted or not an object file.

**no file arguments** *(Assembler)*
The assembler has been invoked without any file arguments. It cannot assemble anything.

**no identifier in declaration** *(Parser)*
The identifier is missing in this declaration. This error can also occur where the compiler has been confused by such things as missing closing braces.

**no input files specified** *(Cromwell)*
Cromwell must have an input file to convert.

**no memory for string buffer** *(Parser)*
P1 was unable to allocate memory for the longest string encountered, as it attempts to sort and merge strings. Try reducing the number or length of strings in this module.

**no output file format specified** *(Cromwell)*
The output format must be specified to Cromwell.

**no psect specified for function variable/argument allocation** *(Linker)*
This is probably caused by omission of correct run-time startoff module. Check the linker arguments, or "Object Files..." in HPD.

**no reserved * areas defined** *(Parser)*
No address ranges were specified with the -RESRAM or -RESOM option.

**no ROM banks defined** *(Driver)*
The -ROM options was invoked but no valid bank address ranges were present.
no ROM range covering address 0 encountered (Driver)
None of the on-chip memory or memory specified with -ROM was found to include address 0. This may have been deliberate.

no room for arguments (Preprocessor, Parser, Code Generator, Linker, Objtohex)
The code generator could not allocate any more memory. Try increasing the size of available memory.

no space for macro def’n (Assembler)
The assembler has run out of memory.

no start record: entry point defaults to zero (Linker)
None of the object files passed to the linker contained a start record. The start address of the program has been set to zero. This may be harmless, but it is recommended that you define a start address in your startup module by using the "END" directive.

no valid entries in chipinfo file (Assembler)
The chipinfo file (libpicinfo.ini by default) contains no valid processor descriptions.

no. of arguments redeclared (Parser)
The number of arguments in this function declaration does not agree with a previous declaration of the same function.

nodecount = * (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

non-constant case label (Code Generator)
A case label in this switch statement has a value which is not a constant.

non-prototyped function declaration: * (Parser)
A function has been declared using old-style (K&R) arguments. It is preferable to use prototype declarations for all functions. If the function has no arguments, declare it as e.g. "int func(void)".

non-scalar types can’t be converted (Parser)
You can’t convert a structure, union or array to anything else. You can convert a pointer to one of those things, so perhaps you left out an ampersand ("&").

non-void function returns no value (Parser)
A function that is declared as returning a value has a "return" statement that does not specify a return value.

not a member of the struct/union * (Parser)
This identifier is not a member of the structure or union type with which it used here.

not a variable identifier: * (Parser)
This identifier is not a variable; it may be some other kind of object, e.g. a label.
Error Messages

not an argument: * (Parser)
This identifier that has appeared in a K&R style argument declarator is not listed inside the parentheses after the function name. Check spelling.

null format name (Cromwell)
The -I or -O option to Cromwell must specify a file format.

object code version is greater than * (Linker)
The object code version of an object module is higher than the highest version the linker is known to work with. Check that you are using the correct linker.

object file is not absolute (Objtohex)
The object file passed to objtohex has relocation items in it. This may indicate it is the wrong object file, or that the linker or objtohex have been given invalid options.

only functions may be qualified interrupt (Parser)
The qualifier "interrupt" may not be applied to anything except a function.

only functions may be void (Parser)
A variable may not be "void". Only a function can be "void".

only Ivalues may be assigned to or modified (Parser)
Only an lvalue (i.e. an identifier or expression directly denoting addressable storage) can be assigned to or otherwise modified. A typecast does not yield an lvalue. To store a value of different type into a variable, take the address of the variable, convert it to a pointer to the desired type, then dereference that pointer, e.g. "*(int *)&x = 1" is legal whereas "(int)x = 1" is not.

only modifier I valid with this format (Parser)
The only modifier that is legal with this format is l (for long).

only modifiers h and I valid with this format (Parser)
Only modifiers h (short) and l (long) are legal with this printf() format specifier.

only register storage class allowed (Parser)
The only storage class allowed for a function parameter is "register".

operand error (Assembler, Optimiser)
The operand to this opcode is invalid. Check you assembler reference manual for the proper form of operands for this instruction.

operands of * not same pointer type (Parser)
The operands of this operator are of different pointer types. This probably means you have used the wrong pointer, but if the code is actually what you intended, use a typecast to suppress the error message.
operands of * not same type  
(Parser)
The operands of this operator are of different pointer. This probably means you have used the wrong variable, but if the code is actually what you intended, use a typecast to suppress the error message.

operator * in incorrect context  
(Preprocessor)
An operator has been encountered in a #if expression that is incorrectly placed, e.g. two binary operators are not separated by a value.

org argument must be a positive constant  
(Assembler)
An argument to the ORG assembler directive must be a positive constant or a symbol which has been equated to a positive constant.

out of far memory  
(Code Generator)
The compiler has run out of far memory. Try removing TSR’s etc. If your system supports EMS memory, the compiler will be able to use up to 64K of this, so if it is not enable, try enabling EMS.

out of memory  
(Code Generator, Assembler, Optimiser)
The compiler has run out of memory. If you have unnecessary TSRs loaded, remove them. If you are running the compiler from inside another program, try running it directly from the command prompt. Similarly, if you are using HPD, try using the command line compiler driver instead.

out of memory allocating * blocks of *  
(Linker)
Memory was required to extend an array but was unavailable.

out of near memory  
(Code Generator)
The compiler has run out of near memory. This is probably due to too many symbol names. Try splitting the program up, or reducing the number of unused symbols in header files etc.

out of space in macro * arg expansion  
(Preprocessor)
A macro argument has exceeded the length of an internal buffer. This buffer is normally 4096 bytes long.

out-of-range case label *  
(Code Generator)
This case label is not a value that the controlling expression can yield, and thus this label will never be selected.

output file cannot be also an input file  
(Linker)
The linker has detected an attempt to write its output file over one of its input files. This cannot be done, because it needs to simultaneously read and write input and output files.

overfreed  
(Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

page width must be >= *  
(Assembler)
The listing page width must be at least * characters. Any less will not allow a properly formatted listing to be produced.
**Error Messages**

**phase error** *(Assembler)*
The assembler has calculated a different value for a symbol on two different passes. This is probably due to bizarre use of macros or conditional assembly.

**pointer required** *(Parser)*
A pointer is required here. This often means you have used "->" with a structure rather than a structure pointer.

**pointer to * argument required** *(Parser)*
A pointer argument is required for this format specifier. Check the number and order of format specifiers and corresponding arguments.

**pointer to non-static object returned** *(Parser)*
This function returns a pointer to a non-static (e.g. automatic) variable. This is likely to be an error, since the storage associated with automatic variables becomes invalid when the function returns.

**popreg: bad reg (*)** *(Code Generator)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**portion of expression has no effect** *(Code Generator)*
Part of this expression has no side effects, and no effect on the value of the expression.

**possible pointer truncation** *(Parser)*
A pointer qualified "far" has been assigned to a default pointer or a pointer qualified "near", or a default pointer has been assigned to a pointer qualified "near". This may result in truncation of the pointer and loss of information, depending on the memory model in use.

**preprocessor assertion failure** *(Preprocessor)*
The argument to a preprocessor #assert directive has evaluated to zero. This is a programmer induced error.

**probable missing ‘}’ in previous block** *(Parser)*
The compiler has encountered what looks like a function or other declaration, but the preceding function has not been ended with a closing brace. This probably means that a closing brace has been omitted from somewhere in the previous function, although it may well not be the last one.

**processor type not defined** *(Assembler)*
The processor must be defined either from the command line (eg. -16c84), via the PROCESSOR assembler directive, or via the LIST assembler directive.

**psect * cannot be in classes *”** *(Linker)*
A psect cannot be in more than one class. This is either due to assembler modules with conflicting class= options, or use of the -C option to the linker.
psect * memory delta redefined: */*  (Linker)
A global psect has been defined with two different deltas.

psect * memory space redefined: */*  (Linker)
A global psect has been defined in two different memory spaces. Either rename one of the psects or, if they are the same psect, place them in the same memory space using the SPACE psect flag.

psect * not loaded on * boundary  (Linker)
This psect has a relocatability requirement that is not met by the load address given in a -P option. For example if a psect must be on a 4K byte boundary, you could not start it at 100H.

psect * not relocated on * boundary  (Linker)
This psect is not relocated on the required boundary. Check the relocatability of the psect and correct the -p option. if necessary.

psect * not specified in -p option  (Linker)
This psect was not specified in a -P or -A option to the linker. It has been linked at the end of the program, which is probably not where you wanted it.

psect * re-orged  (Linker)
This psect has had its start address specified more than once.

psect * selector value redefined  (Linker)
The selector value for this psect has been defined more than once.

psect * type redefined: *  (Linker)
This psect has had its type defined differently by different modules. This probably means you are trying to link incompatible object modules, e.g. linking 386 flat model code with 8086 real mode code.

psect alignment redefined  (Assembler)
The psect alignment has already been defined using the psect ALIGN flag elsewhere.

psect delta redefined  (Assembler)
The DELTA parameter to the PSECT assembler directive’s is different from a previous PSECT directive.

psect exceeds address limit: *  (Linker)
The maximum address of the psect exceeds the limit placed on it using the LIMIT psect flag.

psect exceeds max size: *  (Linker)
The psect has more bytes in it than the maximum allowed as specified using the SIZE psect flag.

psect is absolute: *  (Linker)
This psect is absolute and should not have an address specified in a -P option.
Error Messages

psect limit redefined  (Assembler)
The psect limit has already been defined using the psect LIMIT flag elsewhere.

psect may not be local and global  (Assembler)
A psect may not be declared to be local if it has already been declared to be (default) global.

psect origin multiply defined: *  (Linker)
The origin of this psect is defined more than once.

psect property redefined  (Assembler)
A property of a psect has been defined in more than place to be different.

psect relocability redefined  (Assembler)
The RELOC parameter to the PSECT assembler directive’s is different from a previous PSECT directive.

psect selector redefined  (Linker)
The selector associated with this psect has been defined differently in two or more places.

psect size redefined  (Assembler)
The maximum size of this psect has been defined differently in two or more places.

psect space redefined  (Assembler)
The psect space has already been defined using the psect SPACE flag elsewhere.

pushreg: bad reg (*)  (Code Generator)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

qualifiers redeclared  (Parser)
This function has different qualifiers in different declarations.

radix must be from 2 - 16  (Assembler)
The radix specified using the RADIX or LIST assembler directive must be in the range from 2 (binary) to 16 (hexadecimal).

range check too complex  (Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

read error on *  (Linker)
The linker encountered an error trying to read this file.

record too long  (Objtohex)
This indicates that the object file is not a valid HI-TECH object file.
record too long: *
An object file contained a record with an illegal size. This probably means the file is corrupted or not an object file.

recursive function calls:  
These functions (or function) call each other recursively. One or more of these functions has statically allocated local variables (compiled stack). Either use the reentrant keyword (if supported with this compiler) or recode to avoid recursion.

recursive macro definition of *
The named macro has been defined in such a manner that expanding it causes a recursive expansion of itself!

redefining macro *
The macro specified is being redefined, to something different to the original definition. If you want to deliberately redefine a macro, use #undef first to remove the original definition.

redundant & applied to array
The address operator "&" has been applied to an array. Since using the name of an array gives its address anyway, this is unnecessary and has been ignored.

refc == 0
This is an internal compiler error. Contact HI-TECH Software technical support with details.

regused - bad arg to g
This is an internal compiler error. Contact HI-TECH Software technical support with details.

reloc= must specify a positive constant
The parameter to the PSECT assembler directive’s ’RELOC’ option must be a positive constant number.

relocation error
It is not possible to add together two relocatable quantities. A constant may be added to a relocatable value, and two relocatable addresses in the same psect may be subtracted. An absolute value must be used in various places where the assembler must know a value at assembly time.

relocation offset * out of range *
An object file contained a relocation record with a relocation offset outside the range of the preceding text record. This means the object file is probably corrupted.

relocation too complex
The complex relocation in this expression is too big to be inserted into the object file.

remsym error
This is an internal compiler error. Contact HI-TECH Software technical support with details.
replace what? (Libr)
The librarian requires one or more modules to be listed for replacement when using the ‘r’ key.

rept argument must be >= 0 (Assembler)
The argument to a "REPT" directive must be greater than zero.

reserved * area * - * and * - * could be merged (Driver)
Two address ranges are contiguous. These could have been merged into one reserved range.

reserved * area * - * low bound greater than high bound (Driver)
This structure member is never used. Maybe it isn’t needed at all.

reserved * area * - * overlaps reserved * area * (Driver)
Two address ranges were specified with either then -RESROM or -RESRAM option that overlap.

Reserved * area and reserved ICD * range overlap in region * (Driver)
The -ICD option has been used which reserves memory locations for the debugger. Additional memory
areas have been reserved with the -RESROM or -RESRAM option and these address ranges overlap
those required by the ICD.

ROM bank * low bound greater than high bound (Driver)
An additional memory bank has been defined which has a lower address bound greater than the high
address bound.

ROM bank * overlaps ROM bank * (Driver)
The -ROM options was invoked but no valid bank address ranges were present.

seek error: * (Linker)
The linker could not seek when writing an output file.

segment * overlaps segment * (Linker)
The named segments have overlapping code or data. Check the addresses being assigned by the "-P"
opinion.

set_fact_bit on pic17! (Assembler)
Thi

signatures do not match: * (Linker)
The specified function has different signatures in different modules. This means it has been declared
differently, e.g. it may have been prototyped in one module and not another. Check what declarations
for the function are visible in the two modules specified and make sure they are compatible.

signed bitfields not supported (Parser)
Only unsigned bitfields are supported. If a bitfield is declared to be type "int", the compiler still treats it
as unsigned.
simple integer expression required  
(Parser)  
A simple integral expression is required after the operator "@", used to associate an absolute address with a variable.

simple type required for *  
(Parser)  
A simple type (i.e. not an array or structure) is required as an operand to this operator.

size= must specify a positive constant  
(Assembler)  
The parameter to the PSECT assembler directive’s ’SIZE’ option must be a positive constant number.

sizeof external array * is zero  
(Parser)  
The sizeof an external array evaluates to zero. This is probably due to the array not having an explicit dimension in the extern declaration.

sizeof yields 0  
(Code Generator)  
The code generator has taken the size of an object and found it to be zero. This almost certainly indicates an error in your declaration of a pointer, e.g. you may have declared a pointer to a zero length array. In general, pointers to arrays are of little use. If you require a pointer to an array of objects of unknown length, you only need a pointer to a single object that can then be indexed or incremented.

sizer required after dot  
(Assembler)  
The size of the operand is required. For example, MOV.w indicates that data of word size is to be moved. ’w’ is the ’sizer’.

space= must specify a positive constant  
(Assembler)  
The parameter to the PSECT assembler directive’s ’SPACE’ option must be a positive constant number.

static object has zero size: *  
(Code Generator)  
A static object has been declared, but has a size of zero.

storage class illegal  
(Parser)  
A structure or union member may not be given a storage class. Its storage class is determined by the storage class of the structure.

storage class redeclared  
(Parser)  
A variable or function has been re-declared with a different storage class. This can occur where there are two conflicting declarations, or where an implicit declaration is followed by an actual declaration.

strange character * after ##  
(Preprocessor)  
A character has been seen after the token catenation operator ## that is neither a letter nor a digit. Since the result of this operator must be a legal token, the operands must be tokens containing only letters and digits.

strange character after # *  
(Preprocessor)  
There is an unexpected character after #.
string concatenation across lines
Strings on two lines will be concatenated. Check that this is the desired result.

string expected
The operand to an "asm" statement must be a string enclosed in parentheses.

string lookup failed in coff:get_string()
This is an internal compiler error. Contact HI-TECH Software technical support with details.

struct/union member expected
A structure or union member name must follow a dot ("."), or arrow ("->").

struct/union redefined: *
A structure or union has been defined more than once.

struct/union required
A structure or union identifier is required before a dot (".").

struct/union tag or '{' expected
An identifier denoting a structure or union, or an opening brace must follow a "struct" or "union" keyword.

symbol * cannot be global
There is an error in an object file, where a local symbol has been declared global. This is either an invalid object file, or an internal error in the linker. Try recreating the object file.

symbol * has erroneous psect: *
There is an error in an object file, where a symbol has an invalid psect. This is either an invalid object file, or an internal error in the linker. Try recreating the object file.

symbol * is not external
A symbol has been declared as EXTRN but is also defined in the current module.

symbol * not defined in #undef
The symbol supplied as argument to #undef was not already defined. This is a warning only, but could be avoided by including the #undef in a #ifdef ... #endif block.

symbol cannot be both extern and public
If the symbol is declared as extern, it is to be imported. If it is declared as public, it is to be exported from the current module. It is not possible for a symbol to be both.

symbol has been declared extern
A symbol has been declared in the current module, but has previously been declared extern. A symbol cannot be both local and extern.
syntax error  
(Assembler, Optimiser)  
A syntax error has been detected. This could be caused a number of things.

syntax error in -a spec  
(Linker)  
The -A spec is invalid. A valid -A spec should be something like:

-AROM=1000h-1FFFh

syntax error in checksum list  
(Linker)  
There is a syntax error in a checksum list read by the linker. The checksum list is read from standard input by the linker, in response to an option. Re-read the manual on checksum list.

syntax error in chipinfo file at line *  
(Assembler)  
The chipinfo file contains non-standard syntax at the specified line.

syntax error in local argument  
(Assembler)  
There is a syntax error in a local argument.

text does not start at 0  
(Linker)  
Code in some things must start at zero. Here it doesn’t.

text offset too low  
(Linker)  
You aren’t likely to see this error. Rhubarb!

text record has bad length: *  
(Linker)  
There is an error in an object file. This is either an invalid object file, or an internal error in the linker. Try recreating the object file.

text record has length too small: *  
(Linker)  
This indicates that the object file is not a valid HI-TECH object file.

this function too large - try reducing level of optimization  
(Code Generator)  
A large function has been encountered when using a -Og (global optimization) switch. Try re-compiling without the global optimization, or reduce the size of the function.

this is a struct  
(Parser)  
This identifier following a "union" or "enum" keyword is already the tag for a structure, and thus should only follow the keyword "struct".

this is a union  
(Parser)  
This identifier following a "struct" or "enum" keyword is already the tag for a union, and thus should only follow the keyword "union".
Error Messages

this is an enum  (Parser)
This identifier following a "struct" or "union" keyword is already the tag for an enumerated type, and thus should only follow the keyword "enum".

too few arguments  (Parser)
This function requires more arguments than are provided in this call.

too few arguments for format string  (Parser)
There are too few arguments for this format string. This would result in a garbage value being printed or converted at run time.

too many (*) enumeration constants  (Parser)
There are too many enumeration constants in an enumerated type. The maximum number of enumerated constants allowed in an enumerated type is 512.

too many (*) structure members  (Parser)
There are too many members in a structure or union. The maximum number of members allowed in one structure or union is 512.

too many address spaces - space * ignored  (Linker)
The limit to the number of address spaces is currently 16.

too many arguments  (Parser)
This function does not accept as many arguments as there are here.

too many arguments for format string  (Parser)
There are too many arguments for this format string. This is harmless, but may represent an incorrect format string.

too many arguments for macro  (Preprocessor)
A macro may only have up to 31 parameters, as per the C Standard.

too many arguments in macro expansion  (Preprocessor)
There were too many arguments supplied in a macro invocation. The maximum number allowed is 31.

too many cases in switch  (Code Generator)
There are too many case labels in this switch statement. The maximum allowable number of case labels in any one switch statement is 511.

too many common lines in chipinfo file for *  (Assembler, Driver)
The chipinfo file (libpicinfo.ini by default) contains a processor section with too many COMMON fields. Only one COMMON field is allowed.
too many errors  
*Preprocessor, Parser, Code Generator, Assembler, Linker*
There were so many errors that the compiler has given up. Correct the first few errors and many of the later ones will probably go away.

too many file arguments. usage: cpp [input [output]]  
*Preprocessor*
CPP should be invoked with at most two file arguments.

too many files in coff file  
*Cromwell*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

too many include directories  
*Preprocessor*
A maximum of 7 directories may be specified for the preprocessor to search for include files.

too many initializers  
*Parser*
There are too many initializers for this object. Check the number of initializers against the object definition (array or structure).

too many input files  
*Cromwell*
Too many input files have been specified to be converted by Cromwell.

too many macro parameters  
*Assembler*
There are too many macro parameters on this macro definition.

too many nested #* statements  
*Preprocessor*
#if, #ifdef etc. blocks may only be nested to a maximum of 32.

too many nested #if statements  
*Preprocessor*
#if, #ifdef etc. blocks may only be nested to a maximum of 32.

too many object files  
*Driver*
A maximum of 128 object files may be passed to the linker. The driver exceeded this amount when generating the command line for the linker.

too many output files  
*Cromwell*
Too many output file formats have been specified to Cromwell.

too many psect class specifications  
*Linker*
There are too many psect class specifications (-C options)

too many psect pragmas  
*Code Generator*
Too many "pragma psect" directives have been used.

too many psects  
*Assembler*
There are too many psects! Boy, what a program!
Error Messages

**too many qualifier names** *(Parser)*
There are too many qualifier names specified.

**too many rambank lines in chipinfo file for *** *(Assembler, Driver)*
The chipinfo file (libpicinfo.ini by default) contains a processor section with too many RAMBANK fields. Reduce the number of values.

**too many references to *** *(Cref)*
This is an internal compiler error. Contact HI-TECH Software technical support with details.

**too many relocation items** *(Objtohex)*
Objtohex filled up a table. This program is just way too complex!

**too many segment fixups** *(Objtohex)*
There are too many segment fixups in the object file given to objtohex.

**too many segments** *(Objtohex)*
There are too many segments in the object file given to objtohex.

**too many symbols** *(Assembler)*
There are too many symbols for the assemblers symbol table. Reduce the number of symbols in your program. If it is the linker producing this error, suggest changing some global to local symbols.

**too many symbols (*)** *(Linker)*
There are too many symbols in the symbol table, which has a limit of * symbols. Change some global symbols to local symbols to reduce the number of symbols.

**too many temporary labels** *(Assembler)*
There are too many temporary labels in this assembler file. The assembler allows a maximum of 2000 temporary labels.

**too much indirection** *(Parser)*
A pointer declaration may only have 16 levels of indirection.

**too much pushback** *(Preprocessor)*
This error should not occur, and represents an internal error in the preprocessor.

**type conflict** *(Parser)*
The operands of this operator are of incompatible types.

**type modifier already specified** *(Parser)*
This type modifier has already be specified in this type.

**type modifiers not valid with this format** *(Parser)*
Type modifiers may not be used with this format.
type redeclared (Parser)
The type of this function or object has been redeclared. This can occur because of two incompatible
declarations, or because an implicit declaration is followed by an incompatible declaration.

type specifier reqd. for proto arg (Parser)
A type specifier is required for a prototyped argument. It is not acceptable to just have an identifier.

unable to open list file * (Linker)
The named list file could not be opened.

unbalanced paren's, op is * (Preprocessor)
The evaluation of a #if expression found mismatched parentheses. Check the expression for correct
parenthesisation.

undefined *: * (Parser)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

undefined enum tag: * (Parser)
This enum tag has not been defined.

undefined identifier: * (Parser)
This symbol has been used in the program, but has not been defined or declared. Check for spelling
errors.

undefined shift (* bits) (Code Generator)
An attempt has been made to shift a value by a number of bits equal to or greater than the number of bits
in the data type, e.g. shifting a long by 32 bits. This will produce an undefined result on many processors.
This is non-portable code and is flagged as having undefined results by the C Standard.

undefined struct/union (Parser)
This structure or union tag is undefined. Check spelling etc.

undefined struct/union: * (Parser)
The specified structure or union tag is undefined. Check spelling etc.

undefined symbol * (Assembler)
The named symbol is not defined, and has not been specified "GLOBAL".

undefined symbol * in #if, 0 used (Preprocessor)
A symbol on a #if expression was not a defined preprocessor macro. For the purposes of this expression,
its value has been taken as zero.

undefined symbol in fnaddr record: * (Linker)
The linker has found an undefined symbol in the fnaddr record for a non-reentrant function.
Error Messages

**undefined symbol in fnbreak record: ** (Linker)
The linker has found an undefined symbol in the fnbreak record for a non-reentrant function.

**undefined symbol in fncall record: ** (Linker)
The linker has found an undefined symbol in the fncall record for a non-reentrant function.

**undefined symbol in fnindir record: ** (Linker)
The linker has found an undefined symbol in the fnindir record for a non-reentrant function.

**undefined symbol in fnroot record: ** (Linker)
The linker has found an undefined symbol in the fnroot record for a non-reentrant function.

**undefined symbol in fnsize record: ** (Linker)
The linker has found an undefined symbol in the fnsize record for a non-reentrant function.

**undefined symbol:** (Assembler, Linker)
The symbol following is undefined at link time. This could be due to spelling error, or failure to link an appropriate module.

**undefined symbols:** (Linker)
A list of symbols follows that were undefined at link time.

**undefined temporary label** (Assembler)
A temporary label has been referenced that is not defined. Note that a temporary label must have a number \(\geq 0\).

**undefined variable:** (Parser)
This variable has been used but not defined at this point.

**unexpected end of file** (Linker)
This probably means an object file has been truncated because of a lack of disk space.

**unexpected eof** (Parser)
An end-of-file was encountered unexpectedly. Check syntax.

**unexpected text in #control line ignored** (Preprocessor)
This warning occurs when extra characters appear on the end of a control line, e.g.

```
#endif something
```

The "something" will be ignored, but a warning is issued. It is preferable (and in accordance with Standard C) to enclose the "something" as a comment, e.g.

```
#endif /* something */
```
unexpected \ in #if
The backslash is incorrect in the #if statement.

unknown 'with' psect referenced by psect *
(Linker)
The specified psect has been placed with a psect using the psect 'with' flag. The psect it has been placed with does not exist.

unknown addressing mode *
(Assembler, Optimiser)
An unknown addressing mode was used in the assembly file.

unknown architecture in chipinfo file at line *
(Assembler, Driver)
An chip architecture (family) that is unknown was encountered when reading the chip INI file. Valid architectures are: PIC12, PIC14 and PIC16, representing baseline, midrange and highend devices, respectively.

unknown argument to 'pragma switch': *
(Code Generator)
The '#pragma switch' directive has been used with an invalid switch code generation method. Possible arguments are: auto, simple and direct.

unknown complex operator *
(Linker)
There is an error in an object file. This is either an invalid object file, or an internal error in the linker. Try recreating the object file.

unknown fnrec type *
(Linker)
This indicates that the object file is not a valid HI-TECH object file.

unknown format name '*'
(Cromwell)
The output format specified to Cromwell is unknown.

unknown op * in emobj
(Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown op * in size_psect
(Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown op in emasm(): *
(Assembler)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown option *
(Preprocessor)
This option to the preprocessor is not recognized.

unknown pragma *
(Parser)
An unknown pragma directive was encountered.
Error Messages

unknown predicate *  
Internal error - Contact HI-TECH.

unknown predicate *  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown psect  
The assembler file read by the optimizer has an unknown psect.

unknown psect: *  
This psect has been listed in a -P option, but is not defined in any module within the program.

unknown qualifier '*' given to -a  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown qualifier '*' given to -i  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown record type: *  
An invalid object module has been read by the linker. It is either corrupted or not an object file.

unknown register name *  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unknown symbol type *  
The symbol type encountered is unknown to this linker. Check that the correct linker is being used.

unreachable code  
This section of code will never be executed, because there is no execution path by which it could be reached. Look for missing "break" statements inside a control structure like "while" or "for".

unreasonable matching depth  
This is an internal compiler error. Contact HI-TECH Software technical support with details.

unrecognised line in chipinfo file at line *  
The chipinfo file contains a processor section with an unrecognised line. Look in the chipinfo file for the possibilities.

unrecognized option to -z: *  
The code generator has been passed a -Z option it does not understand. This should not happen if it is invoked with the standard driver.

unrecognized qualifier name after ’strings’  
The pragma ’strings’ requires a list of valid qualifier names. For example

#pragma strings const code
would add const and code to the current string qualifiers. If no qualifiers are specified, all qualification will be removed from subsequent strings. The qualifier names must be recognized by the compiler.

**unterminated \#if[n][def] block from line** *(Preprocessor)*
A \#if or similar block was not terminated with a matching \#endif. The line number is the line on which the \#if block began.

**unterminated comment in included file** *(Preprocessor)*
Comments begun inside an included file must end inside the included file.

**unterminated macro arg** *(Assembler)*
An argument to a macro is not terminated. Note that angle brackets ("< >") are used to quote macro arguments.

**unterminated string** *(Assembler, Optimiser)*
A string constant appears not to have a closing quote missing.

**unterminated string in macro body** *(Preprocessor, Assembler)*
A macro definition contains a string that lacks a closing quote.

**unused constant:** *(Parser)*
This enumerated constant is never used. Maybe it isn’t needed at all.

**unused enum:** *(Parser)*
This enumerated type is never used. Maybe it isn’t needed at all.

**unused label:** *(Parser)*
This label is never used. Maybe it isn’t needed at all.

**unused member:** *(Parser)*
This structure member is never used. Maybe it isn’t needed at all.

**unused structure:** *(Parser)*
This structure tag is never used. Maybe it isn’t needed at all.

**unused typedef:** *(Parser)*
This typedef is never used. Maybe it isn’t needed at all.

**unused union:** *(Parser)*
This union type is never used. Maybe it isn’t needed at all.

**unused variable declaration:** *(Parser)*
This variable is never used. Maybe it isn’t needed at all.

**unused variable definition:** *(Parser)*
This variable is never used. Maybe it isn’t needed at all.
upper case #include files are non-portable  
(Preprocessor)
When using DOS, the case of an #include file does not matter. In other operating systems the case is significant.

variable * must be qualified 'const' to be initialised  
(Parser)
Any initialised variable must be declared 'const', as all initialised variables are placed in ROM, with no copy placed in RAM.

variable may be used before set: *  
(Code Generator)
This variable may be used before it has been assigned a value. Since it is an auto variable, this will result in it having a random value.

void function cannot return value  
(Parser)
A void function cannot return a value. Any "return" statement should not be followed by an expression.

while expected  
(Parser)
The keyword "while" is expected at the end of a "do" statement.

work buffer overflow doing * ##  
(Preprocessor)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

work buffer overflow: *  
(Preprocessor)
This is an internal compiler error. Contact HI-TECH Software technical support with details.

write error (out of disk space?) *  
(Linker)
Probably means that the hard disk is full.

write error on *  
(Assembler, Linker, Cromwell)
A write error occurred on the named file. This probably means you have run out of disk space.

write error on object file  
(Assembler)
An error was reported when the assembler was attempting to write an object file. This probably means there is not enough disk space.

wrong number of macro arguments for * - * instead of *  
(Preprocessor)
A macro has been invoked with the wrong number of arguments.

zero size ROM bank * defined  
(Driver)
An additional memory bank has been defined which has a size of zero.
Library Functions

The functions within the PICC-18 C compiler library are listed in this chapter. Each entry begins with the name of the function. This is followed by information analysed into the following headings.

Synopsis
This is the C definition of the function, and the header file in which it is declared.

Description
This is a narrative description of the function and its purpose.

Example
This is an example of the use of the function. It is usually a complete small program that illustrates the function.

Data types
If any special data types (structures etc.) are defined for use with the function, they are listed here with their C definition. These data types will be defined in the header file given under heading - Synopsis.

See also
This refers you to any allied functions.

Return value
The type and nature of the return value of the function, if any, is given. Information on error returns is also included

Only those headings which are relevant to each function are used.
Library Functions

ABS

Synopsis

#include <stdlib.h>

int abs (int j)

Description

The abs() function returns the absolute value of j.

Example

#include <stdio.h>
#include <stdlib.h>

void main (void)
{
    int a = -5;

    printf("The absolute value of %d is %d\n", a, abs(a));
}

Return Value

The absolute value of j.
ACOS

Synopsis

#include <math.h>

double acos (double f)

Description

The acos() function implements the converse of cos(), i.e. it is passed a value in the range -1 to +1, and returns an angle in radians whose cosine is equal to that value.

Example

#include <math.h>
#include <stdio.h>

/* Print acos() values for -1 to 1 in degrees. */

void
main (void)
{
    float i, a;

    for(i = -1.0; i < 1.0 ; i += 0.1) {
        a = acos(i)*180.0/3.141592;
        printf("acos(%f) = %f degrees\n", i, a);
    }
}

See Also

sin(), cos(), tan(), asin(), atan(), atan2()

Return Value

An angle in radians, in the range 0 to \pi. Where the argument value is outside the domain -1 to 1, the return value will be zero.
ASCTIME

Synopsis

```c
#include <time.h>

char * asctime (struct tm * t)
```

Description

The `asctime()` function takes the time broken down into the `struct tm` structure, pointed to by its argument, and returns a 26 character string describing the current date and time in the format:

```
Sun Sep 16 01:03:52 1973
```

Note the `newline` at the end of the string. The width of each field in the string is fixed. The example gets the current time, converts it to a `struct tm` pointer with `localtime()`, it then converts this to ASCII and prints it. The `time()` function will need to be provided by the user (see `time()` for details).

Example

```c
#include <stdio.h>
#include <time.h>

void
main (void)
{
    time_t clock;
    struct tm * tp;

    time(&clock);
    tp = localtime(&clock);
    printf("%s", asctime(tp));
}
```

See Also

cftime(), gmtime(), localtime(), time()

Return Value

A pointer to the string.

Note

The example will require the user to provide the `time()` routine as it cannot be supplied with the compiler. See `time()` for more details.
Data Types

struct tm {
    int tm_sec;
    int tm_min;
    int tm_hour;
    int tm_mday;
    int tm_mon;
    int tm_year;
    int tm_wday;
    int tm_yday;
    int tm_isdst;
};
**ASIN**

**Synopsis**

```c
#include <math.h>

double asin (double f)
```

**Description**

The `asin()` function implements the converse of `sin()`, i.e. it is passed a value in the range -1 to +1, and returns an angle in radians whose sine is equal to that value.

**Example**

```c
#include <math.h>
#include <stdio.h>

void
main (void)
{
  float i, a;

  for(i = -1.0; i < 1.0 ; i += 0.1) {
    a = asin(i)*180.0/3.141592;
    printf("asin(%f) = %f degrees\n", i, a);
  }
}
```

**See Also**

`sin()`, `cos()`, `tan()`, `acos()`, `atan()`, `atan2()`

**Return Value**

An angle in radians, in the range \(-\pi/2\) to \(+\pi/2\). Where the argument value is outside the domain -1 to 1, the return value will be zero.
ATAN

Synopsis

   #include <math.h>
   
   double atan (double x)

Description

This function returns the arc tangent of its argument, i.e. it returns an angle \( \theta \) in the range \( -\pi/2 \) to \( \pi/2 \) such that \( \tan(\theta) = x \).

Example

   #include <stdio.h>
   #include <math.h>
   
   void
   main (void)
   {
       printf("\%f\n", atan(1.5));
   }

See Also

   sin(), cos(), tan(), asin(), acos(), atan2()

Return Value

   The arc tangent of its argument.
ATAN2

Synopsis

```
#include <math.h>

double atan2 (double y, double x)
```

Description

This function returns the arc tangent of \( \frac{y}{x} \), using the sign of both arguments to determine the quadrant of the return value.

Example

```
#include <stdio.h>
#include <math.h>

void main (void)
{
    printf("%f\n", atan2(1.5, 1));
}
```

See Also

sin(), cos(), tan(), asin(), acos(), atan()

Return Value

The arc tangent of \( \frac{y}{x} \) in the range \(-\pi\) to \(+\pi\) radians. If both \( y \) and \( x \) are zero, a domain error occurs and zero is returned.
**ATOF**

**Synopsis**

```c
#include <stdlib.h>

double atof (const char * s)
```

**Description**

The `atof()` function scans the character string passed to it, skipping leading blanks. It then converts an ASCII representation of a number to a double. The number may be in decimal, normal floating point or scientific notation.

**Example**

```c
#include <stdlib.h>
#include <stdio.h>

void main (void)
{
    char buf[80];
    double i;

    gets(buf);
    i = atof(buf);
    printf("Read %s: converted to %f\n", buf, i);
}
```

**See Also**

`atoi()`, `atol()`

**Return Value**

A double precision floating point number. If no number is found in the string, 0.0 will be returned.
**ATOI**

**Synopsis**

```c
#include <stdlib.h>

int atoi (const char * s)
```

**Description**

The `atoi()` function scans the character string passed to it, skipping leading blanks and reading an optional sign. It then converts an ASCII representation of a decimal number to an integer.

**Example**

```c
#include <stdlib.h>
#include <stdio.h>

void main (void)
{
    char buf[80];
    int i;

    gets(buf);
    i = atoi(buf);
    printf("Read %s: converted to %d\n", buf, i);
}
```

**See Also**

`xtoi()`, `atof()`, `atol()`

**Return Value**

A signed integer. If no number is found in the string, 0 will be returned.
ATOL

Synopsis

    #include <stdlib.h>
    long atol (const char * s)

Description
The atol() function scans the character string passed to it, skipping leading blanks. It then converts an ASCII representation of a decimal number to a long integer.

Example

#include <stdlib.h>
#include <stdio.h>

void main (void)
{
    char buf[80];
    long i;
    gets(buf);
    i = atol(buf);
    printf("Read %s: converted to %ld\n", buf, i);
}

See Also
atoi(), atof()

Return Value
A long integer. If no number is found in the string, 0 will be returned.
**CEIL**

Synopsis

```
#include <math.h>

double ceil (double f)
```

Description

This routine returns the smallest whole number not less than \( f \).

Example

```
#include <stdio.h>
#include <math.h>

void main (void)
{
    double j;

    scanf("%lf", &j);
    printf("The ceiling of %lf is %lf\n", j, ceil(j));
}
```
CLRWDT

Synopsis

```
#include <pic18.h>
CLRWDT();
```

Description
This macro is used to clear the device’s internal watchdog timer.

Example
```
#include <pic18.h>

void main (void)
{
    WDTCON=1; /* enable the WDT */
    CLRWDT();
}
```
COS

Synopsis

#include <math.h>

double cos (double f)

Description

This function yields the cosine of its argument, which is an angle in radians. The cosine is calculated by expansion of a polynomial series approximation.

Example

#include <math.h>
#include <stdio.h>

#define C 3.141592/180.0

void
main (void)
{
    double i;

    for(i = 0 ; i <= 180.0 ; i += 10)
        printf("sin(%3.0f) = %f, cos = %f\n", i, sin(i*C), cos(i*C));
}

See Also

sin(), tan(), asin(), acos(), atan(), atan2()

Return Value

A double in the range -1 to +1.
COSH, SINH, TANH

Synopsis

```
#include <math.h>

double cosh (double f)
double sinh (double f)
double tanh (double f)
```

Description
These functions are the hyperbolic implementations of the trigonometric functions; \( \cos() \), \( \sin() \) and \( \tan() \).

Example
```
#include <stdio.h>
#include <math.h>

void main (void)
{
    printf("%f\n", cosh(1.5));
    printf("%f\n", sinh(1.5));
    printf("%f\n", tanh(1.5));
}
```

Return Value
The function \( \cosh() \) returns the hyperbolic cosine value.
The function \( \sinh() \) returns the hyperbolic sine value.
The function \( \tanh() \) returns the hyperbolic tangent value.
CTIME

Synopsis

#include <time.h>

char * ctime (time_t * t)

Description

The `ctime()` function converts the time in seconds pointed to by its argument to a string of the same form as described for `asctime()`. Thus the example program prints the current time and date.

Example

```c
#include <stdio.h>
#include <time.h>

void main (void)
{
    time_t clock;
    
    time(&clock);
    printf("%s", ctime(&clock));
}
```

See Also

gmtime(), localtime(), asctime(), time()

Return Value

A pointer to the string.

Note

*The example will require the user to provide the `time()` routine as one cannot be supplied with the compiler. See `time()` for more detail.*

Data Types

typedef long time_t;
**DI, EI**

**Synopsis**

```c
#include <pic18.h>

void ei(void)
void di(void)
```

**Description**

ei and di enable and disable interrupts respectively. These are implemented as macros defined in pic.h. They will expand to an in-line assembler instruction that sets or clears the interrupt enable bit.

The example shows the use of ei and di around access to a long variable that is modified during an interrupt. If this was not done, it would be possible to return an incorrect value, if the interrupt occurred between accesses to successive words of the count value.

**Example**

```c
#include <pic18.h>

long count;

void interrupt tick(void)
{
    count++;
}

long getticks(void)
{
    long val;    /* Disable interrupts around access to count, to ensure consistency.*/
    di();
    val = count;
    ei();
    return val;
}
```
**DIV**

**Synopsis**

```c
#include <stdlib.h>

div_t div (int numer, int demon)
```

**Description**

The `div()` function computes the quotient and remainder of the numerator divided by the denominator.

**Example**

```c
#include <stdlib.h>
#include <stdio.h>

void main (void)
{
    div_t x;

    x = div(12345, 66);
    printf("quotient = %d, remainder = %d\n", x.quot, x.rem);
}
```

**Return Value**

Returns the quotient and remainder into the `div_t` structure.

**Data Types**

```c
typedef struct
{
    int quot;
    int rem;
} div_t;
```
EEPROM_READ, EEPROM_WRITE

Synopsis

#include <pic18.h>

unsigned char eeprom_read (unsigned int address);
void eeprom_write (unsigned int address, unsigned char value);

Description
These functions allow access to the on-board eeprom (when present). The eeprom is not in the directly-accessible memory space and a special byte sequence is loaded to the eeprom control registers to access this memory. Writing a value to the eeprom is a slow process and the eeprom_write() function polls the appropriate registers to ensure that any previous writes have completed before writing the next datum.

Reading data is completed in the one cycle and no polling is necessary to check for a read completion.

Example
#include <pic18.h>

void main (void)
{
    unsigned char data;
    unsigned int address = 0x0010;

    data=eeprom_read(address);
    data=EEPROM_READ(address);
    eeprom_write(address, data);
}

See Also
FLASH_READ, FLASH_WRITE, FLASH_ERASE

Note
The high and low priority interrupt are disabled during sensitive sequences required to access EEPROM. Interrupts are restored after the sequence has completed. The EEIF interrupt flag is reset by these functions once the operation is complete.
Both eeprom_read() and eeprom_write() are available in a macro form.
EVAL_POLY

Synopsis

#include <math.h>

double eval_poly (double x, const double * d, int n)

Description

The eval_poly() function evaluates a polynomial, whose coefficients are contained in the array d, at x, for example:

\[ y = x^2 d_2 + x d_1 + d_0. \]

The order of the polynomial is passed in n.

Example

#include <stdio.h>
#include <math.h>

void main (void)
{
    double x, y;
    double d[3] = {1.1, 3.5, 2.7};

    x = 2.2;
    y = eval_poly(x, d, 2);
    printf("The polynomial evaluated at %f is %f\n", x, y);
}

Return Value

A double value, being the polynomial evaluated at x.
### EXP

**Synopsis**

```c
#include <math.h>

double exp (double f)
```

**Description**

The `exp()` routine returns the exponential function of its argument, i.e. $e$ to the power of $f$.

**Example**

```c
#include <math.h>
#include <stdio.h>

void
main (void)
{
    double f;

    for(f = 0.0 ; f <= 5 ; f += 1.0)
        printf("e to %1.0f = %f\n", f, exp(f));
}
```

**See Also**

`log()`, `log10()`, `pow()`
FABS

Synopsis

#include <math.h>

double fabs (double f)

Description
This routine returns the absolute value of its double argument.

Example
#include <stdio.h>
#include <math.h>

void
main (void)
{
    printf("%f %f\n", fabs(1.5), fabs(-1.5));
}

See Also
abs()
**flash_erase, flash_read, flash_write**

**Synopsis**

```c
#include <pic18.h>

void flash_erase (unsigned long addr);
unsigned char flash_read (unsigned long addr);
void flash_write(far unsigned char * source, unsigned char length,
                far unsigned char * dest_addr);
```

**Description**

These functions allow access to the flash memory of the microcontroller (when present).

Reading from the flash memory can be done one byte at a time with use of the `flash_read()` function. `flash_read()` returns the data value found at the specified address in flash memory.

Entire sectors of 64 bytes can be restored to an unprogrammed state (value=`FF`) with use of the `flash_erase()` function. Specifying an address to the `flash_erase()` function, will erase all 64 bytes of the sector that contains the given address.

`flash_write()` copies blocks of data/code from RAM/flash to a new destination in flash memory. `flash_read()` requires a pointer to the data that will be copied, the length of data to copy (in bytes) and a pointer to the destination address in flash memory. This function can be used to update values of variables declared as const. Lengths of data up to 256 bytes may be copied at a time.

**Example**

```c
#include <pic18.h>

void
main (void)
{
    const unsigned char old_text[]="insert text here";
    unsigned char new_text[]="HI-TECH Software";
    
    far unsigned char * source = &new_text[0];
    far unsigned char * dest = &old_text[0];
    unsigned char length = 16;
    
    unsigned char data;
    unsigned int address=0x1000;
    
    data = flash_read(address);
    flash_erase(address);
```
flash_write(source, length, destination);
}

Return Value

*flash_read()* returns the data found at the given address, as an unsigned char.

Note

The *flash_write()* function can be used to update anywhere from 1 to 256 bytes of data at a time, however it is optimal to write in data lengths that are multiples of 64.
Synopsis

```c
#include <math.h>

double floor (double f)
```

Description

This routine returns the largest whole number not greater than \( f \).

Example

```c
#include <stdio.h>
#include <math.h>

void main (void)
{
    printf("%f\n", floor( 1.5 ));
    printf("%f\n", floor( -1.5 ));
}
```
**FREXP**

**Synopsis**

```c
#include <math.h>

double frexp (double f, int * p)
```

**Description**
The `frexp()` function breaks a floating point number into a normalized fraction and an integral power of 2. The integer is stored into the `int` object pointed to by `p`. Its return value `x` is in the interval (0.5, 1.0) or zero, and `f` equals `x` times 2 raised to the power stored in `*p`. If `f` is zero, both parts of the result are zero.

**Example**

```c
#include <math.h>
#include <stdio.h>

void main (void)
{
   double f;
   int i;

   f = frexp(23456.34, &i);
   printf("23456.34 = %f * 2^%d\n", f, i);
}
```

**See Also**

`ldexp()`
GMTIME

Synopsis

```c
#include <time.h>

struct tm * gmtime (time_t * t)
```

Description

This function converts the time pointed to by `t` which is in seconds since 00:00:00 on Jan 1, 1970, into a broken down time stored in a structure as defined in `time.h`. The structure is defined in the 'Data Types' section.

Example

```c
#include <stdio.h>
#include <time.h>

void
main (void)
{
    time_t clock;
    struct tm * tp;

    time(&clock);
    tp = gmtime(&clock);
    printf("It’s %d in London\n", tp->tm_year+1900);
}
```

See Also

`ctime()`, `asctime()`, `time()`, `localtime()`

Return Value

Returns a structure of type `tm`.

Note

The example will require the user to provide the `time()` routine as one cannot be supplied with the compiler. See `time()` for more detail.
**Data Types**

typedef long time_t;

struct tm {
    int   tm_sec;
    int   tm_min;
    int   tm_hour;
    int   tm_mday;
    int   tm_mon;
    int   tm_year;
    int   tm_wday;
    int   tm_yday;
    int   tm_isdst;
};
**ISALNUM, ISALPHA, ISDIGIT, ISLOWER et. al.**

**Synopsis**

```c
#include <ctype.h>

int isalnum (char c)
int isalpha (char c)
int isascii (char c)
int iscntrl (char c)
int isdigit (char c)
int islower (char c)
int isprint (char c)
int isgraph (char c)
int ispunct (char c)
int isspace (char c)
int isupper (char c)
int isxdigit (char c)
```

**Description**

These macros, defined in `ctype.h`, test the supplied character for membership in one of several overlapping groups of characters. Note that all except `isascii()` are defined for `c`, if `isascii(c)` is true or if `c = EOF`.

- `isalnum (c)`: `c` is in 0-9 or a-z or A-Z
- `isalpha (c)`: `c` is in A-Z or a-z
- `isascii (c)`: `c` is a 7 bit ascii character
- `iscntrl (c)`: `c` is a control character
- `isdigit (c)`: `c` is a decimal digit
- `islower (c)`: `c` is in a-z
- `isprint (c)`: `c` is a printing char
- `isgraph (c)`: `c` is a non-space printable character
- `ispunct (c)`: `c` is not alphanumeric
- `isspace (c)`: `c` is a space, tab or newline
- `isupper (c)`: `c` is in A-Z
- `isxdigit (c)`: `c` is in 0-9 or a-f or A-F

**Example**

```c
#include <ctype.h>
#include <stdio.h>
```
void
main (void)
{
    char buf[80];
    int i;

    gets(buf);
    i = 0;
    while(isalnum(buf[i]))
        i++;
    buf[i] = 0;
    printf("'%s' is the word\n", buf);
}

See Also
toupper(), tolower(), toascii()
KBHIT

Synopsis

```c
#include <conio.h>

bit kbhit (void)
```

Description
This function returns 1 if a character has been pressed on the console keyboard, 0 otherwise. Normally
the character would then be read via `getch()`.

Example

```c
#include <conio.h>

void
main (void)
{
  int i;

  while(!kbhit()) {
    cputs("I’m waiting..");
    for(i = 0 ; i != 1000 ; i++)
      continue;
  }
}
```

See Also
`getch()`, `getche()`

Return Value
Returns one if a character has been pressed on the console keyboard, zero otherwise. Note: the return
value is a bit.

Note
*The body of the routine will need to be implemented by the user. The skeleton function will be found in
the sources directory.*
Library Functions

LDEXP

Synopsis

#include <math.h>

double ldexp (double f, int i)

Description

The ldexp() function performs the inverse of frexp() operation; the integer \( i \) is added to the exponent of the floating point \( f \) and the resultant returned.

Example

#include <math.h>
#include <stdio.h>

void main (void)
{
  double f;
  f = ldexp(1.0, 10);
  printf("1.0 * 2^10 = %f\n", f);
}

See Also

frexp()

Return Value

The return value is the integer \( i \) added to the exponent of the floating point value \( f \).
**LDIV**

Synopsis

```
#include <stdlib.h>

ldiv_t ldiv (long number, long denom)
```

Description

The `ldiv()` routine divides the numerator by the denominator, computing the quotient and the remainder. The sign of the quotient is the same as that of the mathematical quotient. Its absolute value is the largest integer which is less than the absolute value of the mathematical quotient.

The `ldiv()` function is similar to the `div()` function, the difference being that the arguments and the members of the returned structure are all of type `long int`.

Example

```c
#include <stdlib.h>
#include <stdio.h>

void
main (void)
{
    ldiv_t lt;

    lt = ldiv(1234567, 12345);
    printf("Quotient = %ld, remainder = %ld\n", lt.quot, lt.rem);
}
```

See Also

`div()`

Return Value

Returns a structure of type `ldiv_t`

Data Types

```c
typedef struct {
    long    quot; /* quotient */
    long    rem;  /* remainder */
} ldiv_t;
```
Library Functions

LOCALTIME

Synopsis

```
#include <time.h>

struct tm * localtime (time_t * t)
```

Description

The `localtime()` function converts the time pointed to by `t` which is in seconds since 00:00:00 on Jan 1, 1970, into a broken down time stored in a structure as defined in `time.h`. The routine `localtime()` takes into account the contents of the global integer `time_zone`. This should contain the number of minutes that the local time zone is westward of Greenwich. Since there is no way under MS-DOS of actually predetermining this value, by default `localtime()` will return the same result as `gmtime()`.

Example

```
#include <stdio.h>
#include <time.h>

char * wday[] = {
   "Sunday", "Monday", "Tuesday", "Wednesday",
   "Thursday", "Friday", "Saturday"
};

void
main (void)
{
   time_t clock;
   struct tm * tp;

   time(&clock);
   tp = localtime(&clock);
   printf("Today is %s\n", wday[tp->tm_wday]);
}
```

See Also

ctime(), asctime(), time()

Return Value

Returns a structure of type `tm`.
Note
The example will require the user to provide the time() routine as one cannot be supplied with the compiler. See time() for more detail.

Data Types

typedef long time_t;
struct tm {
    int     tm_sec;
    int     tm_min;
    int     tm_hour;
    int     tm_mday;
    int     tm_mon;
    int     tm_year;
    int     tm_wday;
    int     tm_yday;
    int     tm_isdst;
};
Library Functions

LOG, LOG10

Synopsis

```
#include <math.h>

double log (double f)
double log10 (double f)
```

Description
The \texttt{log()} function returns the natural logarithm of \texttt{f}. The function \texttt{log10()} returns the logarithm to base 10 of \texttt{f}.

Example

```
#include <math.h>
#include <stdio.h>

void main (void)
{
    double f;

    for(f = 1.0 ; f <= 10.0 ; f += 1.0)
        printf("log(%1.0f) = %f\n", f, log(f));
}
```

See Also

\texttt{exp()}, \texttt{pow()}

Return Value
Zero if the argument is negative.
MEMCHR

Synopsis

#include <string.h>

void * memchr (const void * block, int val, size_t length)

Description

The memchr() function is similar to strchr() except that instead of searching null terminated strings, it searches a block of memory specified by length for a particular byte. Its arguments are a pointer to the memory to be searched, the value of the byte to be searched for, and the length of the block. A pointer to the first occurrence of that byte in the block is returned.

Example

#include <string.h>
#include <stdio.h>

unsigned int ary[] = {1, 5, 0x6789, 0x23};

void main (void)
{
    char * cp;

    cp = memchr(ary, 0x89, sizeof ary);
    if(!cp)
        printf("not found\n");
    else
        printf("Found at offset %u\n", cp - (char *)ary);
}

See Also

strchr()

Return Value

A pointer to the first byte matching the argument if one exists; NULL otherwise.
MEMCMP

Synopsis

#include <string.h>

int memcmp (const void * s1, const void * s2, size_t n)

Description

The `memcmp()` function compares two blocks of memory, of length n, and returns a signed value similar to `strncpy()`. Unlike `strncpy()` the comparison does not stop on a null character. The ASCII collating sequence is used for the comparison, but the effect of including non-ASCII characters in the memory blocks on the sense of the return value is indeterminate. Testing for equality is always reliable.

Example

#include <stdio.h>
#include <string.h>

void main (void)
{
    int buf[10], cow[10], i;

    buf[0] = 1;
    buf[2] = 4;
    cow[0] = 1;
    cow[2] = 5;
    buf[1] = 3;
    cow[1] = 3;
    i = memcmp(buf, cow, 3*sizeof(int));
    if(i < 0)
        printf("less than\n");
    else if(i > 0)
        printf("Greater than\n");
    else
        printf("Equal\n");
}

See Also

`strncpy()`, `strncmp()`, `strchr()`, `memset()`, `memchr()`
Return Value
Returns negative one, zero or one, depending on whether \textit{s1} points to string which is less than, equal to or greater than the string pointed to by \textit{s2} in the collating sequence.
MEMCPY

Synopsis

#include <string.h>

void * memcpy (void * d, const void * s, size_t n)

Description

The `memcpy()` function copies `n` bytes of memory starting from the location pointed to by `s` to the block of memory pointed to by `d`. The result of copying overlapping blocks is undefined. The `memcpy()` function differs from `strcpy()` in that it copies a specified number of bytes, rather than all bytes up to a null terminator.

Example

#include <string.h>
#include <stdio.h>

void
main (void)
{
    char buf[80];

    memset(buf, 0, sizeof buf);
    memcpy(buf, "a partial string", 10);
    printf("buf = '%s'\n", buf);
}

See Also

`strncpy()`, `strncmp()`, `strchr()`, `memset()`

Return Value

The `memcpy()` routine returns its first argument.
**MEMMOVE**

Synopsis

```
#include <string.h>

void * memmove (void * s1, const void * s2, size_t n)
```

Description

The `memmove()` function is similar to the function `memcpy()` except copying of overlapping blocks is handled correctly. That is, it will copy forwards or backwards as appropriate to correctly copy one block to another that overlaps it.

See Also

`strncpy()`, `strncmp()`, `strchr()`, `memcpy()`

Return Value

The function `memmove()` returns its first argument.
MEMSET

Synopsis

```c
#include <string.h>

void * memset (void * s, int c, size_t n)
```

Description

The `memset()` function fills `n` bytes of memory starting at the location pointed to by `s` with the byte `c`.

Example

```c
#include <string.h>
#include <stdio.h>

void
main (void)
{
    char abuf[20];

    strcpy(abuf, "This is a string");
    memset(abuf, 'x', 5);
    printf("buf = '%s'\n", abuf);
}
```

See Also

`strncpy()`, `strncmp()`, `strchr()`, `memcpy()`, `memchr()`
**MODF**

**Synopsis**

```c
#include <math.h>

double modf (double value, double * iptr)
```

**Description**

The `modf()` function splits the argument `value` into integral and fractional parts, each having the same sign as `value`. For example, -3.17 would be split into the integral part (-3) and the fractional part (-0.17).

The integral part is stored as a double in the object pointed to by `intptr`.

**Example**

```c
#include <math.h>
#include <stdio.h>

void main (void)
{
    double i_val, f_val;
    
    f_val = modf( -3.17, &i_val);
}
```

**Return Value**

The signed fractional part of `value`. 
Library Functions

PERSIST_CHECK, PERSIST_VALIDATE

Synopsis

#include <sys.h>

int persist_check (int flag)
void persist_validate (void)

Description

The persist_check() function is used with non-volatile RAM variables, declared with the persistent qualifier. It tests the nvram area, using a magic number stored in a hidden variable by a previous call to persist_validate() and a checksum also calculated by persist_validate(). If the magic number and checksum are correct, it returns true (non-zero). If either are incorrect, it returns zero. In this case it will optionally zero out and re-validate the non-volatile RAM area (by calling persist_validate()). This is done if the flag argument is true.

The persist_validate() routine should be called after each change to a persistent variable. It will set up the magic number and recalculate the checksum.

Example

#include <sys.h>
#include <stdio.h>

persistent long reset_count;

void
main (void)
{
    if(!persist_check(1))
        printf("Reset count invalid - zeroed\n");
    else
        printf("Reset number %ld\n", reset_count);
    reset_count++;
    /* update count */
    persist_validate();
    /* and checksum */
    for(;;)
        continue; /* sleep until next reset */
}

Return Value

FALSE (zero) if the NV-RAM area is invalid; TRUE (non-zero) if the NVRAM area is valid.
**POW**

Synopsis

```c
#include <math.h>

double pow (double f, double p)
```

Description

The `pow()` function raises its first argument, `f`, to the power `p`.

Example

```c
#include <math.h>
#include <stdio.h>

void main (void)
{
    double f;

    for(f = 1.0 ; f <= 10.0 ; f += 1.0)
        printf("pow(2, %1.0f) = %f\n", f, pow(2, f));
}
```

See Also

`log()`, `log10()`, `exp()`

Return Value

`f` to the power of `p`. 
**PRINTF**

**Synopsis**

```c
#include <stdio.h>

unsigned char printf (const char * fmt, ...)
```

**Description**

The `printf()` function is a formatted output routine, operating on `stdout`. There are corresponding routines operating into a string buffer (`sprintf()`). The `printf()` routine is passed a format string, followed by a list of zero or more arguments. In the format string are conversion specifications, each of which is used to print out one of the argument list values.

Each conversion specification is of the form `%m.nc` where the percent symbol `%` introduces a conversion, followed by an optional width specification `m`. The `n` specification is an optional precision specification (introduced by the dot) and `c` is a letter specifying the type of the conversion.

If the character `*` is used in place of a decimal constant, e.g. in the format `%*d`, then one integer argument will be taken from the list to provide that value. The types of conversion are:

- **o** Integer conversion - in radices 8, 16, 16, 10 and 10 respectively. The conversion is signed in the case of `d`, unsigned otherwise. The precision value is the total number of digits to print, and may be used to force leading zeroes. E.g. `%8.4x` will print at least 4 hex digits in an 8 wide field. The letter `X` prints out hexadecimal numbers using the upper case letters A-F rather than a-f as would be printed when using `x`. When the alternate format is specified, a leading zero will be supplied for the octal format, and a leading 0x or 0X for the hex format.

- **s** Print a string - the value argument is assumed to be a character pointer. At most `n` characters from the string will be printed, in a field `m` characters wide.

- **c** The argument is assumed to be a single character and is printed literally.

Any other characters used as conversion specifications will be printed. Thus `%%` will produce a single percent sign.

- **l** Long integer conversion - Preceding the integer conversion key letter with an `l` indicates that the argument list is long.

- **f** Floating point - `m` is the total width and `n` is the number of digits after the decimal point. If `n` is omitted
it defaults to 6. If the precision is zero, the decimal point will be omitted unless the alternate format is specified.

Example

printf("Total = %4d%%", 23)
yields ’Total = 23%’

printf("Size is %lx", size)
where size is a long, prints size
as hexadecimal.

printf("Name = %.8s", "a1234567890")
yields ’Name = a1234567’

printf("xx%*d", 3, 4)
yields ’xx 4’

/* vprintf example */

#include <stdio.h>

int error (char * s, ...)
{
    va_list ap;
    va_start(ap, s);
    printf("Error: ");
    vprintf(s, ap);
    putchar(‘\n’);
    va_end(ap);
}

void main (void)
{
    int i;
    i = 3;
    error("testing 1 2 %d", i);
}

See Also

sprintf()
Library Functions

Return Value
The `printf()` routine returns the number of characters written to `stdout`. NB The return value is a char, NOT an int.

Note
Certain features of printf are only available when linking in alternative libraries. Printing floating point numbers requires that the float to be printed be no larger than the largest possible long integer. In order to use long or float formats, the appropriate supplemental library must be included. See the description on the PICC18 -L library scan option for more details.
Synopsis

```c
#include <stdlib.h>

int rand (void)
```

Description

The `rand()` function is a pseudo-random number generator. It returns an integer in the range 0 to 32767, which changes in a pseudo-random fashion on each call. The algorithm will produce a deterministic sequence if started from the same point. The starting point is set using the `srand()` call. The example shows use of the `time()` function to generate a different starting point for the sequence each time.

Example

```c
#include <stdlib.h>
#include <stdio.h>
#include <time.h>

void main (void)
{
    time_t toc;
    int i;

    time(&toc);
    srand((int)toc);
    for(i = 0 ; i != 10 ; i++)
    {
        printf("%d\t", rand());
        putchar(\n');
    }
}
```

See Also

`srand()`

Note

The example will require the user to provide the `time()` routine as one cannot be supplied with the compiler. See `time()` for more detail.
**SIN**

**Synopsis**

```
#include <math.h>

double sin (double f)
```

**Description**

This function returns the sine function of its argument.

**Example**

```c
#include <math.h>
#include <stdio.h>

#define C 3.141592/180.0

void main (void)
{
    double i;

    for(i = 0 ; i <= 180.0 ; i += 10)
        printf("sin(%3.0f) = %f, cos = %f\n", i, sin(i*C), cos(i*C));
}
```

**See Also**

cos(), tan(), asin(), acos(), atan(), atan2()

**Return Value**

Sine value of \( f \).
**SPRINTF**

Synopsis

```c
#include <stdio.h>

unsigned char sprintf (char *buf, const char * fmt, ...)
```

Description

The `sprintf()` function operates in a similar fashion to `printf()`, except that instead of placing the converted output on the `stdout stream`, the characters are placed in the buffer at `buf`. The resultant string will be null terminated, and the number of characters in the buffer will be returned.

See Also

`printf()`

Return Value

The `sprintf()` routine returns the number of characters placed into the buffer.

NB: The return value is a `char` not an `int`. 
**SQRT**

**Synopsis**

```
#include <math.h>

double sqrt (double f)
```

**Description**
The function `sqrt()` implements a square root routine using Newton’s approximation.

**Example**
```
#include <math.h>
#include <stdio.h>

void main (void)
{
    double i;

    for(i = 0 ; i <= 20.0 ; i += 1.0)
        printf("square root of %.1f = %f\n", i, sqrt(i));
}
```

**See Also**
exp()

**Return Value**
Returns the value of the square root.

**Note**
*A domain error occurs if the argument is negative.*
### SRAND

**Synopsis**

```c
#include <stdlib.h>

void srand (unsigned int seed)
```

**Description**

The `srand()` function initializes the random number generator accessed by `rand()` with the given `seed`. This provides a mechanism for varying the starting point of the pseudo-random sequence yielded by `rand()`. On the z80, a good place to get a truly random seed is from the refresh register. Otherwise timing a response from the console will do, or just using the system time.

**Example**

```c
#include <stdlib.h>
#include <stdio.h>
#include <time.h>

void main (void)
{
    time_t toc;
    int i;
    
    time(&toc);
    srand((int)toc);
    for(i = 0 ; i != 10 ; i++)
        printf("%d\t", rand());
    putchar('\\n');
}
```

**See Also**

`rand()`
STRCAT

Synopsis

#include <string.h>

char * strcat (char * s1, const char * s2)

Description

This function appends (catenates) string s2 to the end of string s1. The result will be null terminated. The argument s1 must point to a character array big enough to hold the resultant string.

Example

#include <string.h>
#include <stdio.h>

void
main (void)
{
    char buffer[256];
    char * s1, * s2;

    strcpy(buffer, "Start of line");
    s1 = buffer;
    s2 = "... end of line";
    strcat(s1, s2);
    printf("Length = %d
", strlen(buffer));
    printf("string = \"%s\"\n", buffer);
}

See Also

strcpy(), strcmp(), strncat(), strlen()

Return Value

The value of s1 is returned.
STRCHR, STRICHR

Synopsis

#include <string.h>

char * strchr (const char * s, int c)
char * strichr (const char * s, int c)

Description

The strchr() function searches the string s for an occurrence of the character c. If one is found, a pointer to that character is returned, otherwise NULL is returned.

The strichr() function is the case-insensitive version of this function.

Example

#include <strings.h>
#include <stdio.h>

#include <stdio.h>

#define void
#define main (void)
{
    static char temp[] = "Here it is...";
    char c = ‘s’;

    if(strchr(temp, c))
        printf("Character %c was found in string\n", c);
    else
        printf("No character was found in string");
}

See Also

strrchr(), strlen(), strcmp()

Return Value

A pointer to the first match found, or NULL if the character does not exist in the string.

Note

Although the function takes an integer argument for the character, only the lower 8 bits of the value are used.
**STRCMP, STRICMP**

**Synopsis**

```c
#include <string.h>

int strcmp (const char * s1, const char * s2)
int stricmp (const char * s1, const char * s2)
```

**Description**

The `strcmp()` function compares its two, null terminated, string arguments and returns a signed integer to indicate whether `s1` is less than, equal to or greater than `s2`. The comparison is done with the standard collating sequence, which is that of the ASCII character set.

The `stricmp()` function is the case-insensitive version of this function.

**Example**

```c
#include <string.h>
#include <stdio.h>

void main (void)
{
    int i;
    if((i = strcmp("ABC", "ABc")) < 0)
        printf("ABC is less than ABc
");
    else if(i > 0)
        printf("ABC is greater than ABc
");
    else
        printf("ABC is equal to ABc
");
}
```

**See Also**

`strlen()`, `strncmp()`, `strcpy()`, `strcat()`

**Return Value**

A signed integer less than, equal to or greater than zero.

**Note**

Other C implementations may use a different collating sequence; the return value is negative, zero or positive, i.e. do not test explicitly for negative one (-1) or one (1).
**STRCPY**

Synopsis

```c
#include <string.h>

char * strcpy (char * s1, const char * s2)
```

Description

This function copies a null terminated string `s2` to a character array pointed to by `s1`. The destination array must be large enough to hold the entire string, including the null terminator.

Example

```c
#include <string.h>
#include <stdio.h>

void
main (void)
{
    char buffer[256];
    char * s1, * s2;

    strcpy(buffer, "Start of line");
    s1 = buffer;
    s2 = "... end of line";
    strcat(s1, s2);
    printf("Length = %d\n", strlen(buffer));
    printf("string = \\
"%s"
", buffer);
}
```

See Also

`strncpy()`, `strlen()`, `strcat()`, `strlen()`

Return Value

The destination buffer pointer `s1` is returned.
**STRCSPN**

Synopsis

```c
#include <string.h>

size_t strcspn (const char * s1, const char * s2)
```

Description

The `strcspn()` function returns the length of the initial segment of the string pointed to by `s1` which consists of characters NOT from the string pointed to by `s2`.

Example

```c
#include <stdio.h>
#include <string.h>

void main (void)
{
    static char set[] = "xyz";

    printf("%d\n", strcspn( "abcdevwxyz", set));
    printf("%d\n", strcspn( "xxxbcadefs", set));
    printf("%d\n", strcspn( "1234567890", set));
}
```

See Also

`strspn()`

Return Value

Returns the length of the segment.
**STRLen**

Synopsis

```c
#include <string.h>

size_t strlen (const char * s)
```

Description

The `strlen()` function returns the number of characters in the string `s`, not including the null terminator.

Example

```c
#include <string.h>
#include <stdio.h>

void main (void)
{
    char buffer[256];
    char * s1, * s2;

    strcpy(buffer, "Start of line");
    s1 = buffer;
    s2 = "... end of line";
    strcat(s1, s2);
    printf("Length = %d\n", strlen(buffer));
    printf("string = \"%s\"\n", buffer);
}
```

Return Value

The number of characters preceding the null terminator.
**STRNCAT**

Synopsis

```c
#include <string.h>
char * strncat (char * s1, const char * s2, size_t n)
```

Description

This function appends (catenates) string `s2` to the end of string `s1`. At most `n` characters will be copied, and the result will be null terminated. `s1` must point to a character array big enough to hold the resultant string.

Example

```c
#include <string.h>
#include <stdio.h>

void main (void)
{
    char buffer[256];
    char * s1, * s2;

    strcpy(buffer, "Start of line");
    s1 = buffer;
    s2 = "... end of line";
    strncat(s1, s2, 5);
    printf("Length = %d\n", strlen(buffer));
    printf("string = \"\%s\"\n", buffer);
}
```

See Also

`strcpy()`, `strcmp()`, `strcat()`, `strlen()`

Return Value

The value of `s1` is returned.
**STRNCMP, STRNICMP**

**Synopsis**

```
#include <string.h>

int strncmp (const char * s1, const char * s2, size_t n)
int strnicmp (const char * s1, const char * s2, size_t n)
```

**Description**

The `strcmp()` function compares its two, null terminated, string arguments, up to a maximum of `n` characters, and returns a signed integer to indicate whether `s1` is less than, equal to or greater than `s2`. The comparison is done with the standard collating sequence, which is that of the ASCII character set.

The `stricmp()` function is the case-insensitive version of this function.

**Example**

```
#include <stdio.h>
#include <string.h>

void main (void)
{
    int i;

    i = strcmp("abcxyz", "abcxyz");
    if(i == 0)
        printf("Both strings are equal\n");
    else if(i > 0)
        printf("String 2 less than string 1\n");
    else
        printf("String 2 is greater than string 1\n");
}
```

**See Also**

`strlen()`, `strcmp()`, `strcpy()`, `strcat()`

**Return Value**

A signed integer less than, equal to or greater than zero.

**Note**

*Other C implementations may use a different collating sequence; the return value is negative, zero or positive, i.e. do not test explicitly for negative one (-1) or one (1).*
Synopsis

```
#include <string.h>

char * strncpy (char * s1, const char * s2, size_t n)
```

Description

This function copies a null terminated string *s2* to a character array pointed to by *s1*. At most *n* characters are copied. If string *s2* is longer than *n* then the destination string will not be null terminated. The destination array must be large enough to hold the entire string, including the null terminator.

Example

```
#include <string.h>
#include <stdio.h>

void main (void)
{
    char buffer[256];
    char * s1, * s2;

    strncpy(buffer, "Start of line", 6);
    s1 = buffer;
    s2 = " ... end of line";
    strcat(s1, s2);
    printf("Length = %d\n", strlen(buffer));
    printf("string = \"%s\"\n", buffer);
}
```

See Also

`strcpy()`, `strcat()`, `strlen()`, `strcmp()`

Return Value

The destination buffer pointer *s1* is returned.
**STRPBRK**

Synopsis

```c
#include <string.h>

char * strpbrk (const char * s1, const char * s2)
```

Description

The `strpbrk()` function returns a pointer to the first occurrence in string `s1` of any character from string `s2`, or a null pointer if no character from `s2` exists in `s1`.

Example

```c
#include <stdio.h>
#include <string.h>

void main (void)
{
    char * str = "This is a string.";

    while(str != NULL) {
        printf( "%s\n", str );
        str = strpbrk( str+1, "aeiou" );
    }
}
```

Return Value

Pointer to the first matching character, or NULL if no character found.
**STRRCHR, STRRICHR**

**Synopsis**

```c
#include <string.h>

char * strrchr (char * s, int c)
char * strrichr (char * s, int c)
```

**Description**

The `strrchr()` function is similar to the `strchr()` function, but searches from the end of the string rather than the beginning, i.e. it locates the last occurrence of the character `c` in the null terminated string `s`. If successful it returns a pointer to that occurrence, otherwise it returns `NULL`.

The `strrichr()` function is the case-insensitive version of this function.

**Example**

```c
#include <stdio.h>
#include <string.h>

void
main (void)
{
    char * str = "This is a string."

    while(str != NULL) {
        printf( "%s\n", str );
        str = strrchr( str+1, 's');
    }
}
```

**See Also**

`strchr()`, `strlen()`, `strcmp()`, `strcpy()`, `strcpy()`, `strcat()`

**Return Value**

A pointer to the character, or `NULL` if none is found.
Synopsis

#include <string.h>

size_t strspn (const char * s1, const char * s2)

Description
The `strspn()` function returns the length of the initial segment of the string pointed to by `s1` which consists entirely of characters from the string pointed to by `s2`.

Example
#include <stdio.h>
#include <string.h>

void main (void)
{
    printf("%d\n", strspn("This is a string", "This"));
    printf("%d\n", strspn("This is a string", "this"));
}

See Also
strcspn()

Return Value
The length of the segment.
STRSTR, STRISTR

Synopsis

```c
#include <string.h>

char * strstr (const char * s1, const char * s2)
char * stristr (const char * s1, const char * s2)
```

Description

The `strstr()` function locates the first occurrence of the sequence of characters in the string pointed to by `s2` in the string pointed to by `s1`.

The `stristr()` routine is the case-insensitive version of this function.

Example

```c
#include <stdio.h>
#include <string.h>

void main (void)
{
    printf("%d\n", strstr("This is a string", "str"));
}
```

Return Value

Pointer to the located string or a null pointer if the string was not found.
STRTOK

Synopsis

#include <string.h>

char * strtok (char * s1, const char * s2)

Description

A number of calls to `strtok()` breaks the string `s1` (which consists of a sequence of zero or more text tokens separated by one or more characters from the separator string `s2`) into its separate tokens.

The first call must have the string `s1`. This call returns a pointer to the first character of the first token, or NULL if no tokens were found. The inter-token separator character is overwritten by a null character, which terminates the current token.

For subsequent calls to `strtok()`, `s1` should be set to a null pointer. These calls start searching from the end of the last token found, and again return a pointer to the first character of the next token, or NULL if no further tokens were found.

Example

#include <stdio.h>
#include <string.h>

void main (void)
{
    char * ptr;
    char * buf = "This is a string of words.";
    char * sep_tok = ".,?! ";

    ptr = strtok(buf, sep_tok);
    while(ptr != NULL) {
        printf("%s\n", ptr);
        ptr = strtok(NULL, sep_tok);
    }
}

Return Value

Returns a pointer to the first character of a token, or a null pointer if no token was found.

Note

The separator string `s2` may be different from call to call.
Library Functions

TAN

Synopsis

    #include <math.h>
    
    double tan (double f)

Description

The tan() function calculates the tangent of \( f \).

Example

    #include <math.h>
    #include <stdio.h>
    #define C 3.141592/180.0

    void
    main (void)
    {
        double i;
        for(i = 0 ; i <= 180.0 ; i += 10)
            printf("tan(%3.0f) = %f\n", i, tan(i*C));
    }

See Also

sin(), cos(), asin(), acos(), atan(), atan2()

Return Value

The tangent of \( f \).
**TIME**

Synopsis

```c
#include <time.h>

time_t time (time_t * t)
```

Description

This function is not provided as it is dependant on the target system supplying the current time. This function will be user implemented. When implemented, this function should return the current time in seconds since 00:00:00 on Jan 1, 1970. If the argument `t` is not equal to NULL, the same value is stored into the object pointed to by `t`.

Example

```c
#include <stdio.h>
#include <time.h>

void main (void)
{
    time_t clock;
    time(&clock);
    printf("%s", ctime(&clock));
}
```

See Also

cftime(), gmtime(), localtime(), asctime()

Return Value

This routine when implemented will return the current time in seconds since 00:00:00 on Jan 1, 1970.

Note

*The time() routine is not supplied, if required the user will have to implement this routine to the specifications outlined above.*
**TOLOWER, TOUPPER, TOASCII**

**Synopsis**

```c
#include <ctype.h>

char toupper (int c)
char tolower (int c)
char toascii (int c)
```

**Description**

The `toupper()` function converts its lower case alphabetic argument to upper case, the `tolower()` routine performs the reverse conversion and the `toascii()` macro returns a result that is guaranteed in the range 0-0177. The functions `toupper()` and `tolower()` return their arguments if it is not an alphabetic character.

**Example**

```c
#include <stdio.h>
#include <ctype.h>
#include <string.h>

void main (void)
{
    char * array1 = "aBcDE";
    int i;

    for(i=0; i < strlen(array1); ++i) {
        printf("%c", tolower(array1[i]));
    }
    printf("\n");
}
```

**See Also**

`islower()`, `isupper()`, `isascii()`, et. al.
VA_START, VA_ARG, VA_END

Synopsis

```c
#include <stdarg.h>

void va_start (va_list ap, parmN)

type va_arg  (ap, type)

void va_end  (va_list ap)
```

Description

These macros are provided to give access in a portable way to parameters to a function represented in a prototype by the ellipsis symbol (...), where type and number of arguments supplied to the function are not known at compile time.

The rightmost parameter to the function (shown as `parmN`) plays an important role in these macros, as it is the starting point for access to further parameters. In a function taking variable numbers of arguments, a variable of type `va_list` should be declared, then the macro `va_start()` invoked with that variable and the name of `parmN`. This will initialize the variable to allow subsequent calls of the macro `va_arg()` to access successive parameters.

Each call to `va_arg()` requires two arguments; the variable previously defined and a type name which is the type that the next parameter is expected to be. Note that any arguments thus accessed will have been widened by the default conventions to `int`, `unsigned int` or `double`. For example if a character argument has been passed, it should be accessed by `va_arg(ap, int)` since the `char` will have been widened to `int`.

An example is given below of a function taking one integer parameter, followed by a number of other parameters. In this example the function expects the subsequent parameters to be pointers to char, but note that the compiler is not aware of this, and it is the programmers responsibility to ensure that correct arguments are supplied.

Example

```c
#include <stdio.h>
#include <stdarg.h>

void pf (int a, ...)
{
    va_list ap;

    va_start(ap, a);
    while(a--)
        puts(va_arg(ap, char *));
}
```


```c
    va_end(ap);
}

void
main (void)
{
    pf(3, "Line 1", "line 2", "line 3");
}
```
XTOI

Synopsis

#include <stdlib.h>

unsigned xtoi (const char * s)

Description
The xtoi() function scans the character string passed to it, skipping leading blanks reading an optional sign, and converts an ASCII representation of a hexadecimal number to an integer.

Example

#include <stdlib.h>
#include <stdio.h>

void main (void)
{
    char buf[80];
    int i;

    gets(buf);
    i = xtoi(buf);
    printf("Read %s: converted to %x\n", buf, i);
}

See Also

atoi()

Return Value

A signed integer. If no number is found in the string, zero will be returned.
__CONFIG

Synopsis

#include <pic18.h>
__CONFIG(n, data)

Description
This macro is used to program the configuration fuses that set the device into various modes of operation.

The macro accepts the number corresponding to the configuration register it is to program, then the 16-Bit value it is to update it with.

16-Bit masks have been defined to describe each programmable attribute available on each device. These attribute masks can be found tabulated in this manual in the Features and Runtime Environment section.

Multiple attributes can be selected by ANDing them together

Example

#include <pic18.h>

__CONFIG(1, RC & OSCEN)
__CONFIG(2, WDTPS16 & BORV45)
__CONFIG(4, DEBUGEN)

void
main (void)
{
}

__EEPROM_DATA

Synopsis

#include <pic18.h>

__EEPROM_DATA(a,b,c,d,e,f,g,h)

Description
This macro is used to store initial values into the device’s EEPROM registers at the time of programming.

The macro must be given blocks of 8 bytes to write each time it is called, and can be called repeatedly to store multiple blocks.

__EEPROM_DATA() will begin writing to EEPROM address zero, and will auto-increment the address written to by 8, each time it is used.

Example

#include <pic18.h>

__EEPROM_DATA(0x00,0x01,0x02,0x03,0x04,0x05,0x06,0x07)
__EEPROM_DATA(0x08,0x09,0x0A,0x0B,0x0C,0x0D,0x0E,0x0F)

void
main (void)
{
}
__IDLOC

Synopsis

#include <pic18.h>
__IDLOC(x)

Description
This macro places data into the device’s special locations outside of addressable memory reserved for ID. This would be useful for storage of serial numbers etc.
The macro will attempt to write 5 nibbles of data to the 5 locations reserved for ID purposes.

Example
#include <pic18.h>

__IDLOC(15F01);
/* will store 1, 5, F, 0 and 1 in the ID registers*/

void
main (void)
{
}
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<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-processor</td>
<td>Define the processor</td>
</tr>
<tr>
<td>-AAHEX</td>
<td>Generate an American Automation symbolic HEX file</td>
</tr>
<tr>
<td>-ASMLIST</td>
<td>Generate assembler .LST file for each compilation</td>
</tr>
<tr>
<td>-Address</td>
<td>Specify offset for ROM code</td>
</tr>
<tr>
<td>-A-option</td>
<td>Specify -option to be passed directly to the assembler</td>
</tr>
<tr>
<td>-BIN</td>
<td>Generate a Binary output file</td>
</tr>
<tr>
<td>-Bl</td>
<td>Selects large memory model</td>
</tr>
<tr>
<td>-Bs</td>
<td>Selects small memory model</td>
</tr>
<tr>
<td>-C</td>
<td>Compile to object files only</td>
</tr>
<tr>
<td>-CKfile</td>
<td>Make OBJTOHEX use a checksum file</td>
</tr>
<tr>
<td>-CP16</td>
<td>Use 16-bit wide pointers to program space</td>
</tr>
<tr>
<td>-CP24</td>
<td>Use 24-bit wide pointers to program space</td>
</tr>
<tr>
<td>-CRfile</td>
<td>Generate cross-reference listing</td>
</tr>
<tr>
<td>-D24</td>
<td>Use truncated 24-bit floating point format for doubles</td>
</tr>
<tr>
<td>-D32</td>
<td>Use IEEE754 32-bit floating point format for doubles</td>
</tr>
<tr>
<td>-Dmacro</td>
<td>Define pre-processor macro</td>
</tr>
<tr>
<td>-E</td>
<td>Use “editor” format for compiler errors</td>
</tr>
<tr>
<td>-E+file</td>
<td>Redirect compiler errors to a file</td>
</tr>
<tr>
<td>-FAKELOCAL</td>
<td>Produce MPLAB-specific debug information</td>
</tr>
<tr>
<td>-FDOUBLE</td>
<td>Use Fast 32-bit floating point libraries</td>
</tr>
<tr>
<td>-Gfile</td>
<td>Generate enhanced source level symbol table</td>
</tr>
<tr>
<td>-HELP</td>
<td>Print summary of options</td>
</tr>
<tr>
<td>-ICD</td>
<td>Compile code for MPLAB-ICD</td>
</tr>
<tr>
<td>-ipath</td>
<td>Specify a directory pathname for include files</td>
</tr>
<tr>
<td>-INTEL</td>
<td>Generate an Intel HEX format output file (default)</td>
</tr>
<tr>
<td>-library</td>
<td>Specify a library to be scanned by the linker</td>
</tr>
<tr>
<td>-L-option</td>
<td>Specify -option to be passed directly to the linker</td>
</tr>
<tr>
<td>-Mfile</td>
<td>Request generation of a MAP file</td>
</tr>
<tr>
<td>-MOT</td>
<td>Generate a Motorola S1/S9 HEX format output file</td>
</tr>
<tr>
<td>-MPLAB</td>
<td>Specify compilation and debugging under MPLAB IDE</td>
</tr>
<tr>
<td>-Nsize</td>
<td>Specify identifier length</td>
</tr>
<tr>
<td>-NODEL</td>
<td>Do not remove temporary/intermediate files.</td>
</tr>
</tbody>
</table>
## PICC18 Options

<table>
<thead>
<tr>
<th>Option</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>-NOERRATA</td>
<td>Disable errata-fix modifications of the output code</td>
</tr>
<tr>
<td>-NORT</td>
<td>Do not link standard runtime module</td>
</tr>
<tr>
<td>-O</td>
<td>Enable post-pass optimization</td>
</tr>
<tr>
<td>-Ofile</td>
<td>Specify output filename</td>
</tr>
<tr>
<td>-P</td>
<td>Preprocess assembler files</td>
</tr>
<tr>
<td>-PRE</td>
<td>Produce preprocessed source files only</td>
</tr>
<tr>
<td>-PROTO</td>
<td>Generate function prototype information</td>
</tr>
<tr>
<td>-PSECTMAP</td>
<td>Display complete memory segment usage after linking</td>
</tr>
<tr>
<td>-q</td>
<td>Specify quiet mode</td>
</tr>
<tr>
<td>-RESRAM&lt;ranges&gt;</td>
<td>Reserve the specified RAM address ranges.</td>
</tr>
<tr>
<td>-RESROM&lt;ranges&gt;</td>
<td>Reserve the specified ROM address ranges.</td>
</tr>
<tr>
<td>-ROM&lt;ranges&gt;</td>
<td>Specify program space memory for ROM-less devices</td>
</tr>
<tr>
<td>-S</td>
<td>Compile to assembler source files only</td>
</tr>
<tr>
<td>-SIGNED_CHAR</td>
<td>Make the default char signed.</td>
</tr>
<tr>
<td>-STRICT</td>
<td>Enable strict ANSI keyword conformance</td>
</tr>
<tr>
<td>-TEK</td>
<td>Generate a Tektronix HEX format output file</td>
</tr>
<tr>
<td>-Usymbol</td>
<td>Undefined a predefined pre-processor symbol</td>
</tr>
<tr>
<td>-UBROF</td>
<td>Generate an UBROF format output file</td>
</tr>
<tr>
<td>-V</td>
<td>Verbose: display compiler pass command lines</td>
</tr>
<tr>
<td>-Wlevel</td>
<td>Set compiler warning level</td>
</tr>
<tr>
<td>-Wlevel!</td>
<td>Set compiler warning level and stop on warnings</td>
</tr>
<tr>
<td>-X</td>
<td>Eliminate local symbols from symbol table</td>
</tr>
<tr>
<td>-Zg[level]</td>
<td>Enable global optimization in the code generator</td>
</tr>
</tbody>
</table>