## Homework Set #5 Polar codes

- 1. Polarization and the idea of polar codes (28 Points): The question is about polarization effect in memoryless channels that can lead to simple coding schemes that achieve the capacity which are called polar codes.
  - (a) Consider the channel in Fig. 1 where two parallel binary erasure channels can be used at once (the input is  $X = (X_1, X_2)$ ). The inputs alphabets are binary, so that  $Y_1$  and

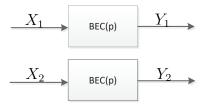


Figure 1: Two parallel binary erasure channels

 $Y_2$  are the outputs of a BEC(p) with inputs  $X_1$  and  $X_2$ , respectively. Compute the capacity of this channel, namely,

$$\max_{p(x_1, x_2)} I(X_1, X_2; Y_1, Y_2).$$
(1)

What is the input distribution  $p(x_1, x_2)$  that achieves the capacity?

(b) Consider the system in Fig. 2, where addition is modulo 2:

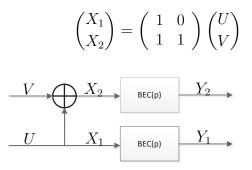


Figure 2: Two parallel binary erasure channels with modified inputs

Compute the capacity of the new channel, i.e.  $\max_{p(u,v)} I(U, V; Y_1, Y_2)$ .

What is the p(u, v) that achieves the capacity?

Next, the channel is decomposed into two parallel channels as appears in Fig. 3. The input of Channel 1 is U and its output is  $(Y_1, Y_2, V)$ . The input of Channel 2 is V and its output is  $(Y_1, Y_2)$ .

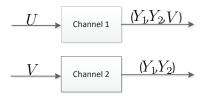


Figure 3: Two new channels

- (c) Compute the expressions  $I(U; Y_1, Y_2, V)$  and  $I(V; Y_1, Y_2)$  with respect to the p(u, v) that achieves the maximum in (b). What is the sum of the expressions you computed?
- (d) Compare the mutual information of Channels 1 and 2 with the capacity of a binary erasure channel (that is, write  $\langle , \rangle$  or = with simple proof).

\*For large n, repeating this decomposition n times, ends up in nc clean channels and in n(1-c) totally noisy channels. This is the main idea of polar codes, which achieves capacity.