

**Homework Set #1**  
**Properties of Entropy, Mutual Information and Divergence**

**1. Entropy of functions of a random variable.**

Let  $X$  be a discrete random variable. Show that the entropy of a function of  $X$  is less than or equal to the entropy of  $X$  by justifying the following steps:

$$\begin{aligned} H(X, g(X)) &\stackrel{(a)}{=} H(X) + H(g(X)|X) \\ &\stackrel{(b)}{=} H(X). \\ H(X, g(X)) &\stackrel{(c)}{=} H(g(X)) + H(X|g(X)) \\ &\stackrel{(d)}{\geq} H(g(X)). \end{aligned}$$

Thus  $H(g(X)) \leq H(X)$ .

**2. Example of joint entropy.**

Let  $p(x, y)$  be given by

	Y	
X	0	1
0	$\frac{1}{3}$	$\frac{1}{3}$
1	0	$\frac{1}{3}$

Find

- (a)  $H(X), H(Y)$ .
- (b)  $H(X|Y), H(Y|X)$ .
- (c)  $H(X, Y)$ .
- (d)  $H(Y) - H(Y|X)$ .
- (e)  $I(X; Y)$ .

**3. Bytes.**

The entropy,  $H_a(X) = -\sum p(x) \log_a p(x)$  is expressed in bits if the logarithm is to the base 2 and in bytes if the logarithm is to the base 256. What is the relationship of  $H_2(X)$  to  $H_{256}(X)$ ?

4. **Two looks.**

Here is a statement about pairwise independence and joint independence. Let  $X, Y_1$ , and  $Y_2$  be binary random variables. If  $I(X; Y_1) = 0$  and  $I(X; Y_2) = 0$ , does it follow that  $I(X; Y_1, Y_2) = 0$ ?

- (a) Yes or no?
- (b) Prove or provide a counterexample.
- (c) If  $I(X; Y_1) = 0$  and  $I(X; Y_2) = 0$  in the above problem, does it follow that  $I(Y_1; Y_2) = 0$ ? In other words, if  $Y_1$  is independent of  $X$ , and if  $Y_2$  is independent of  $X$ , is it true that  $Y_1$  and  $Y_2$  are independent?

5. **A measure of correlation.**

Let  $X_1$  and  $X_2$  be *identically distributed*, but not necessarily independent. Let

$$\rho = 1 - \frac{H(X_2|X_1)}{H(X_1)}.$$

- (a) Show  $\rho = \frac{I(X_1; X_2)}{H(X_1)}$ . (There is no typo in the definition of  $\rho$ )
- (b) Show  $0 \leq \rho \leq 1$ .
- (c) When is  $\rho = 0$ ?
- (d) When is  $\rho = 1$ ?

6. **The value of a question.**

Let  $X \sim p(x)$ ,  $x = 1, 2, \dots, m$ .

We are given a set  $S \subseteq \{1, 2, \dots, m\}$ . We ask whether  $X \in S$  and receive the answer

$$Y = \begin{cases} 1, & \text{if } X \in S \\ 0, & \text{if } X \notin S. \end{cases}$$

Suppose  $\Pr\{X \in S\} = \alpha$ .

- (a) Find the decrease in uncertainty  $H(X) - H(X|Y)$ .
- (b) Is the set  $S$  with a given probability  $\alpha$  is as good as any other  $S' \neq S$  with  $\Pr\{X \in S'\} = \alpha$ ?

7. **Relative entropy is not symmetric**

Let the random variable  $X$  have three possible outcomes  $\{a, b, c\}$ . Consider two distributions on this random variable

Symbol	$p(x)$	$q(x)$
a	1/2	1/3
b	1/4	1/3
c	1/4	1/3

Calculate  $H(p)$ ,  $H(q)$ ,  $D(p \parallel q)$  and  $D(q \parallel p)$ .

Verify that in this case  $D(p \parallel q) \neq D(q \parallel p)$ .

8. **“True or False” questions**

Copy each relation and write **true** or **false**. Then, if it's true, prove it. If it is false give a counterexample or prove that the opposite is true.

(a)  $H(X) \geq H(X|Y)$

(b)  $H(X) + H(Y) \leq H(X, Y)$

(c) Let  $X, Y$  be two independent random variables. Then

$$H(X - Y) \geq H(X).$$

(d) Let  $X, Y, Z$  be three random variables that satisfies  $H(X, Y) = H(X) + H(Y)$  and  $H(Y, Z) = H(Z) + H(Y)$ . Then the following holds

$$H(X, Y, Z) = H(X) + H(Y) + H(Z).$$

(e) For any  $X, Y, Z$  and the deterministic function  $f, g$   $I(X; Y|Z) = I(X, f(X, Y); Y, g(Y, Z)|Z)$ .

9. **Entropy of 3 pairwise independent random variables**

Let  $W, X, Y$  be 3 random variables distributed each Bernoulli (0.5) that are pairwise independent, i.e.,  $I(W; X) = I(X; Y) = I(W; Y) = 0$ .

(a) What is the **maximum** possible value of  $H(W, X, Y)$ ?

(b) What is the condition under which this **maximum** is achieved?

(c) What is the **minimum** possible value of  $H(W, X, Y)$ ?

(d) Give a specific example achieving this **minimum**.

10. **Joint Entropy** Consider  $n$  different discrete random variables, named  $X_1, X_2, \dots, X_n$ . Each random variable separately has an entropy  $H(X_i)$ , for  $1 \leq i \leq n$ .

(a) What is the upper bound on the joint entropy  $H(X_1, X_2, \dots, X_n)$  of all these random variables  $X_1, X_2, \dots, X_n$  given that  $H(X_i)$ , for  $1 \leq i \leq n$  are fixed?

(b) Under what conditions will this upper bound be reached?

(c) What is the lower bound on the joint entropy  $H(X_1, X_2, \dots, X_n)$  of all these random variables?

(d) Under what condition will this upper bound be reached?

11. **More question of True or False**

Let  $X, Y, Z$  be discrete random variable. Copy each relation and write **true** or **false**. If it's true, prove it. If it is false give a counterexample or prove that the opposite is true.

For instance:

- $H(X) \geq H(X|Y)$  is **true**. Proof: In the class we showed that  $I(X;Y) > 0$ , hence  $H(X) - H(X|Y) > 0$ .
- $H(X) + H(Y) \leq H(X, Y)$  is **false**. Actually the opposite is true, i.e.,  $H(X) + H(Y) \geq H(X, Y)$  since  $I(X;Y) = H(X) + H(Y) - H(X, Y) \geq 0$ .

(a) If  $H(X|Y) = H(X)$  then  $X$  and  $Y$  are independent.

(b) For any two probability mass functions (pmf)  $P, Q$ ,

$$D\left(\frac{P+Q}{2} \parallel Q\right) \leq \frac{1}{2}D(P \parallel Q),$$

where  $D(\parallel)$  is a divergence between two pmfs.

(c) Let  $X$  and  $Y$  be two independent random variables. Then

$$H(X + Y) \geq H(X).$$

- (d)  $I(X;Y) - I(X;Y|Z) \leq H(Z)$
- (e) If  $f(x, y)$  is a convex function in the pair  $(x, y)$ , then for a fixed  $y$ ,  $f(x, y)$  is convex in  $x$ , and for a fixed  $x$ ,  $f(x, y)$  is convex in  $y$ .
- (f) If for a fixed  $y$  the function  $f(x, y)$  is a convex function in  $x$ , and for a fixed  $x$ ,  $f(x, y)$  is convex function in  $y$ , then  $f(x, y)$  is convex in the pair  $(x, y)$ . (Examples of such functions are  $f(x, y) = f_1(x) + f_2(y)$  or  $f(x, y) = f_1(x)f_2(y)$  where  $f_1(x)$  and  $f_2(y)$  are convex.)
- (g) Let  $X, Y, Z, W$  satisfy the Markov chain  $X - Y - Z$  and  $Y - Z - W$ . Does the Markov  $X - Y - Z - W$  hold? (The Markov  $X - Y - Z - W$  means that  $P(x|y, z, w) = P(x|y)$  and  $P(x, y|z, w) = P(x, y|z)$ .)
- (h)  $H(X|Z)$  is concave in  $P_{X|Z}$  for fixed  $P_Z$ .

12. **Random questions.**

One wishes to identify a random object  $X \sim p(x)$ . A question  $Q \sim r(q)$  is asked at random according to  $r(q)$ . This results in a deterministic answer  $A = A(x, q) \in \{a_1, a_2, \dots\}$ . Suppose the object  $X$  and the question  $Q$  are independent. Then  $I(X; Q, A)$  is the uncertainty in  $X$  removed by the question-answer  $(Q, A)$ .

- (a) Show  $I(X; Q, A) = H(A|Q)$ . Interpret.
- (b) Now suppose that two i.i.d. questions  $Q_1, Q_2 \sim r(q)$  are asked, eliciting answers  $A_1$  and  $A_2$ . Show that two questions are less valuable than twice the value of a single question in the sense that  $I(X; Q_1, A_1, Q_2, A_2) \leq 2I(X; Q_1, A_1)$ .

13. **Entropy bounds.**

Let  $X \sim p(x)$ , where  $x$  takes values in an alphabet  $\mathcal{X}$  of size  $m$ . The entropy  $H(X)$  is given by

$$\begin{aligned} H(X) &\equiv - \sum_{x \in \mathcal{X}} p(x) \log p(x) \\ &= E_p \log \frac{1}{p(X)}. \end{aligned}$$

Use Jensen's inequality ( $Ef(X) \leq f(EX)$ , if  $f$  is concave) to show

- (a)  $H(X) \leq \log E_p \frac{1}{p(X)}$   
 $= \log m$ .

- (b)  $-H(X) \leq \log(\sum_{x \in \mathcal{X}} p^2(x))$ , thus establishing a lower bound on  $H(X)$ .
- (c) Evaluate the upper and lower bounds on  $H(X)$  when  $p(x)$  is uniform.
- (d) Let  $X_1, X_2$  be two independent drawings of  $X$ . Find  $\Pr\{X_1 = X_2\}$  and show  $\Pr\{X_1 = X_2\} \geq 2^{-H}$ .

14. **Bottleneck.**

Suppose a (non-stationary) Markov chain starts in one of  $n$  states, necks down to  $k < n$  states, and then fans back to  $m > k$  states. Thus  $X_1 \rightarrow X_2 \rightarrow X_3$ ,  $X_1 \in \{1, 2, \dots, n\}$ ,  $X_2 \in \{1, 2, \dots, k\}$ ,  $X_3 \in \{1, 2, \dots, m\}$ , and  $p(x_1, x_2, x_3) = p(x_1)p(x_2|x_1)p(x_3|x_2)$ .

- (a) Show that the dependence of  $X_1$  and  $X_3$  is limited by the bottleneck by proving that  $I(X_1; X_3) \leq \log k$ .
- (b) Evaluate  $I(X_1; X_3)$  for  $k = 1$ , and conclude that no dependence can survive such a bottleneck.

15. **Convexity of Halfspaces, hyperplanes and polyhedron**

Let  $\mathbf{x}$  be a real vector of finite dimension  $n$ , i.e.,  $x \in \mathbb{R}^n$ . A *halfspace* is the set of all  $x \in \mathbb{R}^n$  that satisfies  $\mathbf{a}^T \mathbf{x} \leq b$ , where  $a \neq 0$ . In other words a halfspace is the set

$$\{\mathbf{x} \in \mathbb{R}^n : \mathbf{a}^T \mathbf{x} \leq b\}.$$

A hyperplan is the set of the form

$$\{\mathbf{x} \in \mathbb{R}^n : \mathbf{a}^T \mathbf{x} = b\}.$$

- (a) Show that a halfspace and a hyperplan are convex sets.
- (b) show that for any two sets  $\mathcal{A}$  and  $\mathcal{B}$  that are convex the intersection  $\mathcal{A} \cap \mathcal{B}$  is also convex.
- (c) A *polyhedron* is an intersection of halfspaces and a hyperplans. Deduce that a polyhedron is a convex set.

- (d) A probability vector  $\mathbf{x}$  is such that each element is positive and it sums to 1. Is the set of all vector probabilities of dimension  $n$  (called the probability simplex) a halfspace, hyperplan or polyhedron?

**16. Some sets of probability distributions.**

Let  $X$  be a real-valued random variable with  $\Pr(X = a_i) = p_i, i = 1, \dots, n$ , where  $a_1 < a_2 < \dots < a_n$ . Let  $\mathbf{p}$  denote the vector  $p_1, p_2, \dots, p_n$ . Of course  $\mathbf{p} \in \mathbb{R}^n$  lies in the standard probability simplex. Which of the following conditions are convex in  $\mathbf{p}$ ? (That is, for which of the following conditions is the set of  $\mathbf{p} \in \mathbf{P}$  that satisfy the condition convex?)

- (a)  $\alpha \leq E[f(X)] \leq \beta$ , where  $E[f(X)]$  is the expected value of  $f(X)$ , i.e.  $E[f(x)] = \sum_{i=1}^n p_i f(a_i)$  (The function  $f : \mathbb{R} \mapsto \mathbb{R}$  is given.)  
 (b)  $\Pr(X > \alpha) \leq \beta$   
 (c)  $E[|X^3|] \leq \alpha E[|X|]$ .  
 (d)  $\text{var}(X) \leq \alpha$ , where  $\text{var}(X) = E(X - EX)^2$  is the variance of  $X$ .  
 (e)  $E[X^2] \leq \alpha$   
 (f)  $E[X^2] \geq \alpha$

**17. Perspective transformation preserve convexity** Let  $f(x), f : \mathbb{R} \rightarrow \mathbb{R}$ , be a convex function.

- (a) Show that the function

$$tf\left(\frac{x}{t}\right), \tag{1}$$

is a convex function in the pair  $(x, t)$  for  $t > 0$ . (The function  $tf\left(\frac{x}{t}\right)$  is called perspective transformation of  $f(x)$ .)

- (b) Is the preservation true for concave functions too?  
 (c) Use this property to prove that  $D(P||Q)$  is a convex function in  $(P, Q)$ .

**18. Coin Tosses**

Consider the next joint distribution:  $X$  is the number of coin tosses until the first head appears and  $Y$  is the number of coin tosses until the second head appears. The probability for a head is  $q$ , and the tosses are independent.

- a. Compute the distribution of  $X$ ,  $p(x)$ , the distribution of  $Y$ ,  $p(y)$ , and the conditional distributions  $p(y|x)$  and  $p(x|y)$ .
- b. Compute  $H(X)$ ,  $H(Y|X)$ ,  $H(X, Y)$ . Each term should not include a series. Hint: Is  $H(Y|X) = H(Y - X|X)$ ?
- c. Compute  $H(Y)$ ,  $H(X|Y)$ , and  $I(X; Y)$ . If necessary, answers may include a series.

19. **Inequalities** Copy each relation to your notebook and write  $\leq$ ,  $\geq$  or  $=$ , prove it.

- (a) Let  $X$  be a discrete random variable. Compare  $\frac{1}{2^{H(X)}}$  vs.  $\max_x p(x)$ .
- (b) Let  $H_b(a)$  denote the binary entropy for  $a \in [0, 1]$  and  $H_{ter}$  is the ternary entropy i.e.  $H_{ter}(a, b, c) = -a \log a - b \log b - c \log c$ , where  $p_1, p_2, p_3 \in [0, 1]$ , and  $p_1 + p_2 + p_3 = 1$ . Compare  $H_{ter}(ab, a\bar{b}, \bar{a})$  vs  $H_b(a) + \bar{a}H_b(b)$ .

20. **True or False of a constrained inequality:**

Given are three discrete random variables  $X, Y, Z$  that satisfy  $H(Y|X, Z) = 0$ .

- (a) Copy the next relation to your notebook and write **true** or **false**.

$$I(X; Y) \geq H(Y) - H(Z)$$

- (b) What are the conditions for which the equality  $I(X; Y) = H(Y) - H(Z)$  holds.
- (c) Assume that the conditions for  $I(X; Y) = H(Y) - H(Z)$  are satisfied. Is it true that there exists a function such that  $Z = g(Y)$ ?

21. **True or False of:** Copy each relation to your notebook and write **true** or **false**. If true, prove the statement, and if not provide a counterexample.



(a) Let  $X - Y - Z - W$  be a Markov chain, then the following holds:

$$I(X; W) \leq I(Y; Z).$$

(b) For two probability distributions,  $p_{XY}$  and  $q_{XY}$ , that are defined on  $\mathcal{X} \times \mathcal{Y}$ , the following holds:

$$D(p_{XY} || q_{XY}) \geq D(p_X || q_X).$$

(c) If  $X$  and  $Y$  are dependent and also  $Y$  and  $Z$  are dependent, then  $X$  and  $Z$  are dependent.

22. **Cross entropy:**

Often in Machine learning, cross entropy is used to measure performance of a classifier model such as neural network. Cross entropy is defined for two PMFs  $P_X$  and  $Q_X$  as

$$H(P_X, Q_X) \triangleq - \sum_{x \in \mathcal{X}} P_X(x) \log Q_X(x).$$

In a shorter notation we write as

$$H(P, Q) \triangleq - \sum_{x \in \mathcal{X}} P(x) \log Q(x).$$

Copy each of the following relations to your notebook and write **true** or **false** and provide a proof/disproof.

- (a)  $0 \leq H(P, Q) \leq \log |\mathcal{X}|$  for all  $P, Q$ .
- (b)  $\min_Q H(P, Q) = H(P, P)$  for all  $P$ .
- (c)  $H(P, Q)$  is concave in the pair  $(P, Q)$ .
- (d)  $H(P, Q)$  is convex in the pair  $(P, Q)$ .

23. **Properties of mutual information:** A joint distribution is given by  $P(x, \theta, y) = P(x)P(\theta)P(y|x, \theta)$ . Answer the following three questions:

(a) **True/False:** Is it true that there is a Markov chain  $X - Y - \theta$ ? Prove or provide a counter example.

- (b) **Inequalities:** Fill (and prove) one of the relations  $\leq, =, \geq$  between the following expressions :

$$I(X; Y) \quad ??? \quad I(X; Y|\theta).$$

- (c) **Convex/Concave:** Determine whether the mutual information,  $I(X_1; X_2)$  is convex OR concave function of  $P(x_2|x_1)$  for a fixed  $P(x_1)$ . **Hint: You can use your answers from the previous questions.**