

Homework Set #2
Entropy, Mutual Information, Divergence and Jensen's Inequality

1. **The value of a question.**

Let $X \sim p(x)$, $x = 1, 2, \dots, m$.

We are given a set $S \subseteq \{1, 2, \dots, m\}$. We ask whether $X \in S$ and receive the answer

$$Y = \begin{cases} 1, & \text{if } X \in S \\ 0, & \text{if } X \notin S. \end{cases}$$

Suppose $\Pr\{X \in S\} = \alpha$.

- (a) Find the decrease in uncertainty $H(X) - H(X|Y)$.
- (b) Is the set S with a given probability α is as good as any other.

2. **Relative entropy is not symmetric**

Let the random variable X have three possible outcomes $\{a, b, c\}$. Consider two distributions on this random variable

Symbol	$p(x)$	$q(x)$
a	1/2	1/3
b	1/4	1/3
c	1/4	1/3

Calculate $H(p)$, $H(q)$, $D(p \parallel q)$ and $D(q \parallel p)$.

Verify that in this case $D(p \parallel q) \neq D(q \parallel p)$.

3. **True or False**

Let X, Y, Z be discrete random variable. Copy each relation and write **true** or **false**. If it's true, prove it. If it is false give a counterexample or prove that the opposite is true.

For instance:

- $H(X) \geq H(X|Y)$ is **true**. Proof: In the class we showed that $I(X;Y) > 0$, hence $H(X) - H(X|Y) > 0$.
- $H(X) + H(Y) \leq H(X,Y)$ is **false**. Actually the opposite is true, i.e., $H(X) + H(Y) \geq H(X,Y)$ since $I(X;Y) = H(X) + H(Y) - H(X,Y) \geq 0$.

- (a) If $H(X|Y) = H(X)$ then X and Y are independent.
 (b) For any two probability mass functions (pmf) P, Q ,

$$D\left(\frac{P+Q}{2} \parallel Q\right) \leq \frac{1}{2}D(P \parallel Q),$$

where $D(\parallel)$ is a divergence between two pmfs.

- (c) Let X and Y be two independent random variables. Then

$$H(X+Y) \geq H(X).$$

- (d) $I(X;Y) - I(X;Y|Z) \leq H(Z)$
 (e) If $f(x,y)$ is a convex function in the pair (x,y) , then for a fixed y , $f(x,y)$ is convex in x , and for a fixed x , $f(x,y)$ is convex in y .
 (f) If for a fixed y the function $f(x,y)$ is a convex function in x , and for a fixed x , $f(x,y)$ is convex function in y , then $f(x,y)$ is convex in the pair (x,y) . (Examples of such functions are $f(x,y) = f_1(x) + f_2(y)$ or $f(x,y) = f_1(x)f_2(y)$ where $f_1(x)$ and $f_2(y)$ are convex.)

4. Random questions.

One wishes to identify a random object $X \sim p(x)$. A question $Q \sim r(q)$ is asked at random according to $r(q)$. This results in a deterministic answer $A = A(x,q) \in \{a_1, a_2, \dots\}$. Suppose the object X and the question Q are independent. Then $I(X;Q,A)$ is the uncertainty in X removed by the question-answer (Q,A) .

- (a) Show $I(X;Q,A) = H(A|Q)$. Interpret.
 (b) Now suppose that two i.i.d. questions $Q_1, Q_2 \sim r(q)$ are asked, eliciting answers A_1 and A_2 . Show that two questions are less valuable than twice the value of a single question in the sense that $I(X;Q_1, A_1, Q_2, A_2) \leq 2I(X;Q_1, A_1)$.

5. Entropy bounds.

Let $X \sim p(x)$, where x takes values in an alphabet \mathcal{X} of size m . The entropy $H(X)$ is given by

$$\begin{aligned} H(X) &\equiv -\sum_{x \in \mathcal{X}} p(x) \log p(x) \\ &= E_p \log \frac{1}{p(X)}. \end{aligned}$$

Use Jensen's inequality ($Ef(X) \leq f(EX)$, if f is concave) to show

- (a) $H(X) \leq \log E_p \frac{1}{p(X)}$
 $= \log m$.
- (b) $-H(X) \leq \log(\sum_{x \in \mathcal{X}} p^2(x))$, thus establishing a lower bound on $H(X)$.
- (c) Evaluate the upper and lower bounds on $H(X)$ when $p(x)$ is uniform.
- (d) Let X_1, X_2 be two independent drawings of X . Find $\Pr\{X_1 = X_2\}$ and show $\Pr\{X_1 = X_2\} \geq 2^{-H}$.

6. Bottleneck.

Suppose a (non-stationary) Markov chain starts in one of n states, necks down to $k < n$ states, and then fans back to $m > k$ states. Thus $X_1 \rightarrow X_2 \rightarrow X_3$, $X_1 \in \{1, 2, \dots, n\}$, $X_2 \in \{1, 2, \dots, k\}$, $X_3 \in \{1, 2, \dots, m\}$, and $p(x_1, x_2, x_3) = p(x_1)p(x_2|x_1)p(x_3|x_2)$.

- (a) Show that the dependence of X_1 and X_3 is limited by the bottleneck by proving that $I(X_1; X_3) \leq \log k$.
- (b) Evaluate $I(X_1; X_3)$ for $k = 1$, and conclude that no dependence can survive such a bottleneck.