

## Homework 2

### Capacity of the semi-deterministic relay channel

#### 1) Semi-Deterministic Relay Channel.

Consider a relay channel as introduced in the class that is given by the conditional probability  $P(y, y_1|x, x_1)$ . We have learned a coding scheme called partial decode and forward. This scheme turned to be optimal for a semi-deterministic relay channel, where  $Y_1$  is a function of  $(X, X_1)$ , i.e.,  $y_1 = f(x, x_1)$ . Hence, the joint distribution of the semi-deterministic relay channel is of the form  $P(y, y_1|x, x_1) = P(y|x, x_1)\mathbf{1}_{y_1=f(x, x_1)}$  where  $\mathbf{1}_{y_1=f(x, x_1)}$  is 1 if  $y_1 = f(x, x_1)$  and zero otherwise.

a) Show that any rate satisfying

$$R \leq \min(H(Y_1|X_1) + I(X; Y|X_1, Y_1), I(X, X_1; Y)) \quad (1)$$

for some  $P(x, x_1)$  is achievable.

b) Show that a rate that is achievable must satisfy (1) for some joint distribution  $P(x, x_1)$ .

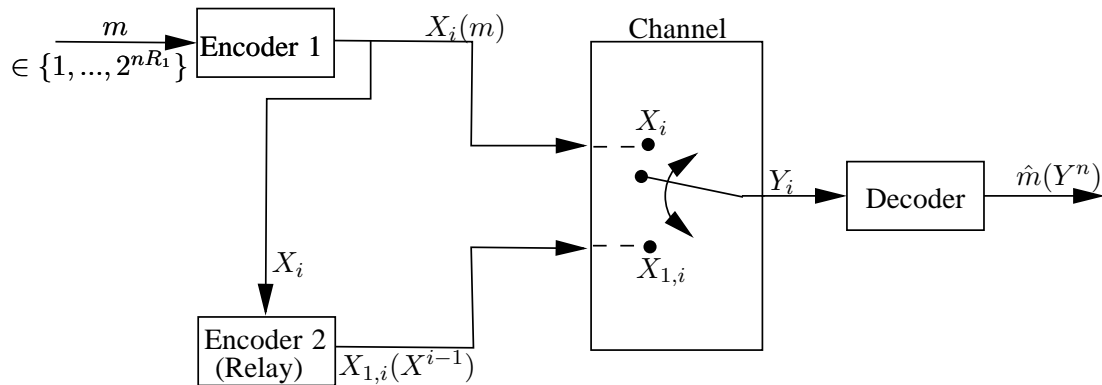


Fig. 1. Semi deterministic relay channel:  $Y_1 = X$  and  $Y$  is chosen with equal probability to be either  $X_1$  or  $X$ .

c) Consider the example in Fig. 1. The relay observe  $X$  and encode it with a delay as in the regular relay setting. The input to the channel is  $(X_i(m), X_{1,i}(X^{i-1}))$ . The output channel at time  $i$  i.e.,  $Y_i$ , is randomly chosen with equal probability to be either  $X_i$  or  $X_{1,i}$ . Find the capacity of the example in Fig. 1. If analytical solution is not possible, you may use the computer to obtain a numerical solution.

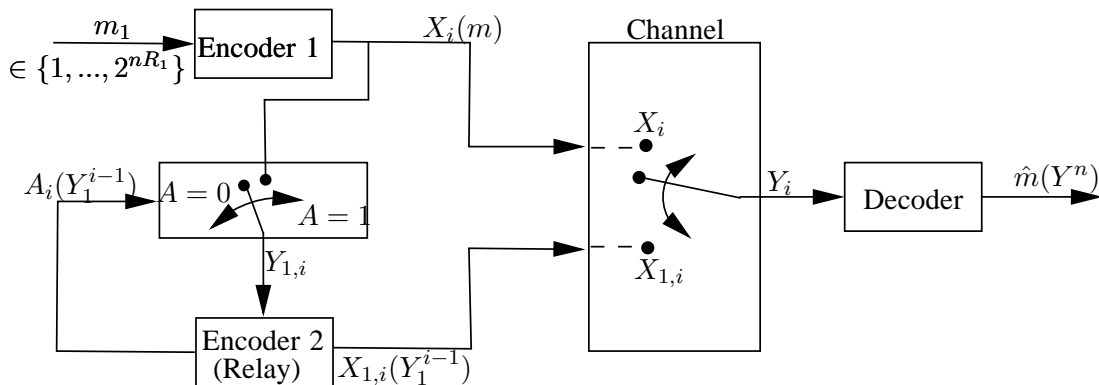


Fig. 2. Semi deterministic Relay with action. If  $A_i = 1$  the relay observe  $X_i$  and if  $A_i = 0$  the relay does not observe  $X_i$ . There exists a constraint on the actions such that  $\frac{1}{n} \sum_{i=1}^n E[A_i] \leq \Gamma$ .

d) Now, consider the case where the relay is taking an action  $A_i$ . The role of the action is to decide if the relay observe or not observe  $X_i$ . Namely, if  $A_i = 1$  than  $Y_{1,i} = X_i$  and if  $A_i = 0$  than  $Y_{1,i}$  is a

constant. Similar to  $X_{1,i}$  which may depend on  $Y_1^{i-1}$  the action  $A_i$  may also depend on  $Y_1^{i-1}$ . Fig. 2 depicts the setting. There exists a constraint on the relay action that  $\frac{1}{n} \sum_{i=1}^n E[A_i] \leq \Gamma$ .

Define the code of the setting and find the capacity region as a function of  $\Gamma$ . (Hint: Note that  $A_i$  has a similar role as  $X_{1,i}$ ; hence can you formulate the setting of Fig. 2 as being a regular formulation of semi-deterministic relay channel?)

- e) Find the capacity region (numerically or analytically) where  $\Gamma = 0$  and explain the result.
- f) Draw using a computer the capacity region as a function of  $\Gamma$ .

Good Luck!!!